

# Equivariant dendroidal sets

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We extend the Cisinski–Moerdijk–Weiss theory of  $\infty$ -operads to the equivariant setting to obtain a notion of  $G$ - $\infty$ -operads that encode “equivariant operads with norm maps” up to homotopy. At the root of this work is the identification of a suitable category of  $G$ -trees together with a notion of  $G$ -inner horns capable of encoding the compositions of norm maps.

Additionally, we follow Blumberg and Hill by constructing suitable variants associated to each of the indexing systems featured in their work.

55U10, 55U35, 55U40; 18G30

## 1 Introduction

Operads encode a variety of algebraic structures, such as monoids, commutative monoids or (depending on the ambient category) Lie algebras,  $E_n$ -algebras, etc. Indeed, all such instances can be regarded as categories of algebras for some (fixed) suitable operad. Informally, an operad  $\mathcal{O}$  consists of “sets/spaces of  $n$ -ary operations”  $\mathcal{O}(n)$ ,  $n \geq 0$ , each of which carries a  $\Sigma_n$ -action encoding “reordering the inputs of the operations”, and a suitable notion of “composition of operations”.

From the homotopy theory point of view, one of the most important classes of operads is certainly that of the  $E_\infty$ -operads, which are “up to homotopy” replacements of the commutative operad  $\text{Com}$ . More concretely, while algebras for  $\text{Com}$  are the usual commutative monoids, the algebras for an  $E_\infty$ -operad are “up to homotopy commutative monoids”, where associativity and commutativity are only enforced up to homotopy. Further,  $E_\infty$ -operads  $\mathcal{O}$  are characterized by the property that each space  $\mathcal{O}(n)$  is a contractible space with a free  $\Sigma_n$ -action.

This work lies at the intersection of operad theory and equivariant homotopy theory. Briefly, in  $G$ -equivariant homotopy theory a map of  $G$ -spaces  $X \rightarrow Y$  is considered a  $G$ -weak equivalence only if all the induced fix point maps  $X^H \rightarrow Y^H$ ,  $H \leq G$  are weak equivalences. Therefore, it is no surprise that the characterization of  $G$ - $E_\infty$ -operads, ie  $G$ -equivariant operads whose algebras are “ $G$ -equivariant up to homotopy

commutative monoids”, would need to be modified. Indeed, a naive first guess might be that a  $G$ -operad  $\mathcal{O}$  should be called  $G$ - $E_\infty$  if (i) each space  $\mathcal{O}(n)$  has a free  $\Sigma_n$ -action and (ii)  $\mathcal{O}(n)$  is  $G$ -contractible. Accepting this tentative characterization for the moment, such a  $G$ -operad is easily produced: simply taking a (nonequivariant)  $E_\infty$ -operad and giving it a trivial  $G$ -action yields such an example. However, it has long been known that such “ $G$ -trivial  $E_\infty$ -operads” are not the correct replacement for the commutative operad in the equivariant setting; see Costenoble and Waner [9]. To see why, we consider the much-studied example of  $R$  a (strictly) commutative  $G$ -ring spectrum. For a finite  $G$ -set  $T$  with  $n$  elements it is possible to equip  $R^{\wedge T} \simeq R^{\wedge n}$  with a mixed  $G$ -action combining the actions on  $R$  and  $T$ . One often writes  $N^T R$  for  $R^{\wedge T}$  together with this action and calls it the Hill–Hopkins–Ravenel *norm*. Multiplication then induces *norm maps*

$$(1.1) \quad N^T R \rightarrow R$$

satisfying equivariance and associativity conditions. The flaw of “ $G$ -trivial  $E_\infty$ -operads” is then that they lack all norm maps (1.1) with  $T$  a nontrivial  $G$ -set (or, after restriction to  $H \leq G$ , with  $T$  a nontrivial  $H$ -set).

In understanding this issue, note first that though  $\mathcal{O}(n)$  has a  $(G \times \Sigma_n)$ -action when  $\mathcal{O}$  is a  $G$ -operad, conditions (i) and (ii) above actually fail to determine a *unique*  $(G \times \Sigma_n)$ -homotopy type. Indeed, (i) implies that  $\mathcal{O}(n)^\Gamma = \emptyset$  whenever  $\Gamma \cap \Sigma_n \neq *$  while (ii) implies that  $\mathcal{O}(n)^\Gamma \sim *$  if  $\Gamma \leq G$ , but these conditions leave out many subgroups  $\Gamma \leq G \times \Sigma_n$ . Indeed, there are identifications

$$(1.2) \quad \Gamma \text{ such that } \Gamma \cap \Sigma_n = * \Leftrightarrow \text{graph of } G \geq H \rightarrow \Sigma_n \Leftrightarrow H\text{-action on } \{1, \dots, n\}$$

and (ii) covers only those  $\Gamma$  encoding trivial  $H$ -actions. The correct characterization of  $G$ - $E_\infty$ -operads is then that (i)  $\mathcal{O}(n)$  is  $\Sigma_n$ -free; (ii')  $\mathcal{O}(n)$  is graph-contractible; ie  $\mathcal{O}(n)^\Gamma \sim *$  for any  $\Gamma \leq G \times \Sigma_n$  such that  $\Gamma \cap \Sigma_n = *$ .

A key observation of Blumberg and Hill in [2] is that the reason a “ $G$ -trivial  $E_\infty$ -operad” induces only the norm maps for trivial sets is that it satisfies (ii') only for those  $\Gamma$  encoding trivial sets. Indeed, their work takes this observation much further. Motivated by the study of equivariant spectra in incomplete universes, they define a whole lattice of types of  $G$ -operads, which they dub  $N_\infty$ -operads, and which satisfy (ii') only for  $\Gamma$  encoding  $H$ -sets within certain special families. Further, they call such special families *indexing systems*. “ $G$ -trivial  $E_\infty$ ” and  $G$ - $E_\infty$  are then the

minimum and maximum types of  $N_\infty$ -operads, with the remaining types interpolating in between.

The motivation for this paper (and the larger project it belongs to) is the observation that the closure conditions for the  $H$ -sets in an indexing system identified in [2, Definition 3.22] admit a nice diagrammatic interpretation (discussed in Section 9) and that this suggests the possibility of encoding equivariant operads with norm maps (ie with operations in  $\mathcal{O}(n)^\Gamma$ ) via suitable diagrammatic models.

Indeed, it is well known that composition of operations in an operad can be encoded using tree diagrams, and work of Moerdijk and Weiss in [17], with follow-up work of Cisinski and Moerdijk in [5], builds a category  $\Omega$  of trees and a model structure on the presheaf category  $\mathbf{dSet} = \mathbf{Set}^{\Omega^{\text{op}}}$  which is shown in the follow-up papers [6] and [7] to be Quillen equivalent to the category of colored simplicial operads.

The role of this paper is to provide the equivariant analogue of the work in [17] and [5]. We first identify a (nonobvious) category  $\Omega_G$  of  $G$ -trees (introduced in Section 4.3 and formally defined in Section 5.3) capable of encoding norm maps and their compositions, and then adapt the proofs in [17] and [5] to prove the existence of a model structure on  $\mathbf{dSet}^G$  whose fibrant objects, which we call  $G$ - $\infty$ -operads, are “up to homotopy  $G$ -operads with norm maps”. We note that our results are not formal: indeed, while our proofs closely follow those in [17] and [5] the presence of equivariance often requires significant modifications. Moreover, we note that alternative so called “genuine” model structures on  $\mathbf{dSet}^G$  built by formal methods (say, by mimicking the definition of the genuine model structure on  $\mathbf{Top}^G$ ) would instead only model  $G$ -operads *without nontrivial norm maps*.

## Acknowledgements

The current work owes much to Peter Bonventre: the category  $\Omega_G$  of  $G$ -trees, which is at the root of this paper, is a joint discovery with him; the author first heard of the notion of broad posets (due to Weiss) from him; and lastly, this work has also greatly benefited from extensive joint conversations.

I would also like to thank the referee for helpful suggestions.

## 2 Main results

Our main result follows. It is the equivariant analogue of [5, Theorem 2.4].

**Theorem 2.1** *There exists a model structure on  $\text{dSet}^G$  such that*

- *the cofibrations are the  $G$ -normal monomorphisms;*
- *the fibrant objects are the  $G$ - $\infty$ -operads;*
- *the weak equivalences are the smallest class that contains the  $G$ -inner anodyne extensions and the trivial fibrations and is closed under “2-out-of-3”.*

[Theorem 2.1](#) will be proven as the combination of [Proposition 8.8](#), [Theorem 8.22](#) and [Corollary 8.23](#).

Further, letting  $\mathcal{F}$  denote an indexing system (see [[2](#), Definition 3.22], and as reinterpreted in [Definition 9.5](#)), we also prove the following more general result.

**Theorem 2.2** *For  $\mathcal{F}$  a weak indexing system there exists a model structure on  $\text{dSet}^G$  such that*

- *the cofibrations are the  $\mathcal{F}$ -normal monomorphisms;*
- *the fibrant objects are the  $\mathcal{F}$ - $\infty$ -operads;*
- *the weak equivalences are the smallest class that contains the  $\mathcal{F}$ -inner anodyne extensions and the  $\mathcal{F}$ -trivial fibrations and is closed under “2-out-of-3”.*

[Theorem 2.2](#) is proven at the end of [Section 9](#).

**Remark 2.3** In the special case where  $\mathcal{F}$  is the indexing system containing only the trivial  $H$ -sets, the model structure given by [Theorem 2.2](#) coincides with the “formal genuine model structure” as built using [[18](#), Proposition 2.6] (for the collection of all subgroups  $H \leq G$ ). However, this is not the case for either [Theorem 2.1](#) or indeed the vast majority of instances of [Theorem 2.2](#), which are not formal consequences of the existence of the model structure on  $\text{dSet}$ .

### 3 Outline

After reviewing the familiar types of trees found elsewhere in the literature (eg [[16](#); [17](#); [5](#)], among others), [Section 4](#) provides an introductory look at the new equivariant trees that motivate this paper, focusing on examples. Most notably, each  $G$ -tree can be represented by two *distinctly* shaped tree diagrams, called the *expanded* and *orbital* representations, each capturing different key features.

[Section 5](#) lays the necessary framework for our work. Specifically, [Section 5.1](#) recalls Weiss’s algebraic *broad poset* model [[19](#)] for the category  $\Omega$  of trees, which we prefer

since planar representations of  $G$ -trees can easily get prohibitively large. Section 5.2 discusses forests, which play an auxiliary role. Section 5.3 formally introduces the category  $\Omega_G$  of  $G$ -trees. Lastly, Section 5.4 introduces all the necessary presheaf categories, most notably the category  $\text{dSet}^G$  featured in Theorems 2.1 and 2.2.

Section 6 discusses the notions of  $G$ -normal monomorphism and  $G$ - $\infty$ -operad needed to state Theorem 2.1. The former of these is straightforward, but the latter requires the key (and more subtle) notion of  $G$ -inner horn (Definition 6.11).

Section 7 is the technical heart of the paper, extending the key technical results [17, Proposition 9.2] and [5, Theorems 5.2 and 4.2] concerning tensor products of dendroidal sets and the dendroidal join to the equivariant setting (Theorems 7.1, 7.2 and 7.4).

Section 8 then finishes the proof of Theorem 2.1 by combining the results of Section 7 with the arguments in the proof of the original nonequivariant result [5, Theorem 2.4].

Finally, Section 9 proves Theorem 2.2 by straightforward generalizations of our arguments to the framework of general indexing systems.

## 4 An introduction to equivariant trees

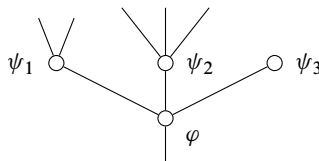
### 4.1 Planar trees

Operads are a tool for studying various types of algebraic structures that possess operations of several arities. More concretely, an operad  $\mathcal{O}$  consists of a sequence of sets (or, more importantly to us, spaces/simplicial sets)  $\mathcal{O}(n)$ ,  $n \geq 0$ , which behave as sets (spaces/simplicial sets) of  $n$ -ary operations. That is, one should have *composition product* maps (see [15, Definition 1.1] or [10, Definition 1.4])

$$(4.1) \quad \begin{aligned} \mathcal{O}(k) \times \mathcal{O}(n_1) \times \cdots \times \mathcal{O}(n_k) &\xrightarrow{\circ} \mathcal{O}(n_1 + \cdots + n_k), \\ (\varphi, \psi_1, \dots, \psi_k) &\longmapsto \varphi(\psi_1, \dots, \psi_k), \end{aligned}$$

and an identity  $\text{id} \in \mathcal{O}(1)$  satisfying suitable associativity and unital conditions.

A powerful tool for visualizing operadic compositions and their compatibilities is given by tree diagrams. For instance, the tree

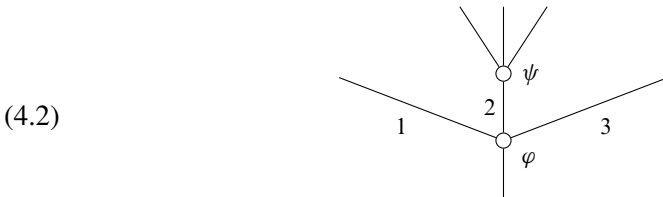


encodes the composition of operations  $\varphi \in \mathcal{O}(3)$ ,  $\psi_1 \in \mathcal{O}(2)$ ,  $\psi_2 \in \mathcal{O}(3)$  and  $\psi_3 \in \mathcal{O}(0)$ , and one has  $\varphi(\psi_1, \psi_2, \psi_3) \in \mathcal{O}(5)$ , with arity of the composite given by counting *leaves* (ie the edges at the top of the tree, not capped by a circle).

Alternatively, given the presence of the identity  $\text{id} \in \mathcal{O}(1)$ , one can instead define operads using so called *partial composition products* [14, Definition 1.16]

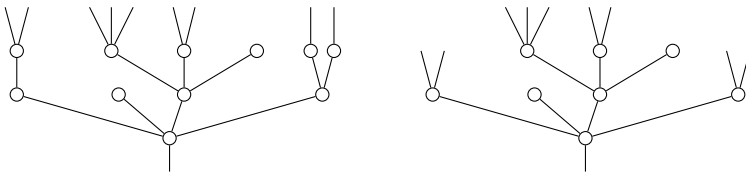
$$\begin{aligned} \mathcal{O}(k) \times \mathcal{O}(n) &\xrightarrow{\circ_i} \mathcal{O}(n + k - 1), \\ (\varphi, \psi) &\longmapsto \varphi(\text{id}, \dots, \text{id}, \psi, \text{id}, \dots, \text{id}), \end{aligned}$$

which are also readily visualized using trees. For example,



encodes the partial composition  $\varphi \circ_2 \psi = \varphi(\text{id}, \psi, \text{id})$ .

Heuristically, trees encoding iterations of  $\circ$  are naturally tiered whereas trees encoding iterations of  $\circ_i$  operations are not, as exemplified by the following:



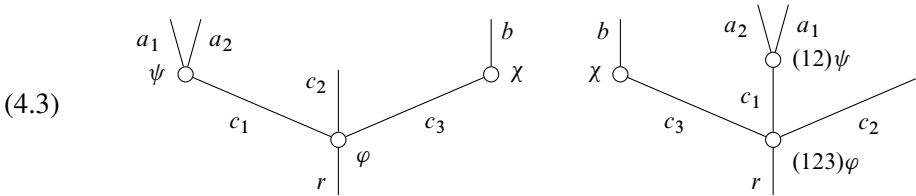
In the leftmost tree, encoding an iterated composition of  $\circ$ , all leaves appear at the same height and the operations, encoded by *nodes* (ie the circles) are naturally divided into levels. On the other hand, this fails for the rightmost tree, which encodes iterated compositions of  $\circ_i$ . Indeed, while the definition  $\varphi \circ_i \psi = \varphi(\text{id}, \dots, \text{id}, \psi, \text{id}, \dots, \text{id})$  would allow us to convert the rightmost tree into a tiered tree by inserting nodes labeled by *id*, there are multiple ways to do so (indeed, the leftmost tree represents one such possibility).

In practice, trees of the second type seem to be the most convenient and we will henceforth work only with such trees.

### 4.2 Symmetric trees

The (planar) tree notation just described is suitable for working with so-called “non- $\Sigma$  operads”. In many applications, however, operads possess an additional piece of structure: each set (space/simplicial set)  $\mathcal{O}(n)$  has a left action of the symmetric group  $\Sigma_n$ . Heuristically, the role of this action is to “change the order of the inputs of an operation”: thinking of  $\varphi \in \mathcal{O}(n)$  as an operation  $x_1, \dots, x_n \mapsto \varphi(x_1, \dots, x_n)$  and letting  $\sigma \in \Sigma_n$ , then  $\sigma\varphi \in \mathcal{O}(n)$  would correspond to the operation  $x_1, \dots, x_n \mapsto \varphi(x_{\sigma(1)}, \dots, x_{\sigma(n)})$ .

When representing compositions on a (symmetric) operad, it thus becomes convenient to think of the edges above a node, which represent the inputs for the operation labeling the node, as not having a fixed order. One immediate drawback of this perspective, however, is that drawing a planar representation of such a tree on paper necessarily requires choosing an (arbitrary) order for the input edges of every node. Therefore, it is possible for different planar representations to encode the exact same information. For example, the pictures



display two planar representations of the same tree that encode *the same* composition data. To explain why, we first point out that  $a_1, a_2, b, c_1, c_2, c_3$  are simply the *names* of the edges of the tree (needed so as to distinguish different representations of the tree in the plane), by contrast with  $\psi, \chi, (123)\varphi$ , etc, which are operations in  $\mathcal{O}$ . Next, for a finite set  $S$  of size  $n$ , denote  $(\text{Iso}(\{1, \dots, n\}, S) \times \mathcal{O}(n))_{\Sigma_n}$  by  $\mathcal{O}(S)$ , where the orbits  $(-)\Sigma_n$  are defined using the diagonal action, acting on the  $\text{Iso}(\{1, \dots, n\}, S)$  component by precomposition with the inverse. Note that though  $\mathcal{O}(S)$  is of course isomorphic to  $\mathcal{O}(n)$ , the isomorphism is not canonical, depending on the choice of an isomorphism  $\{1, \dots, n\} \xrightarrow{\cong} S$ . With this convention, both trees in (4.3) represent the same instance of a composition

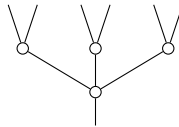
$$\mathcal{O}(\{c_1, c_2, c_3\}) \times \mathcal{O}(\{a_1, a_2\}) \times \mathcal{O}(\{b\}) \xrightarrow{(-)\circ_{c_1}(-)\circ_{c_3}(-)} \mathcal{O}(\{a_1, a_2, c_2, b\}).$$

The reason for the differing labels on the nodes of the trees is then that different planar representations correspond to different choices of isomorphisms  $\{1, \dots, n\} \xrightarrow{\cong} S$ .

For example, the leftmost tree uses the identification  $\{1, 2, 3\} \simeq \{c_1, c_2, c_3\}$  while the rightmost tree uses  $\{1, 2, 3\} \simeq \{c_3, c_1, c_2\}$ , so that, for example, the classes  $[(\{1, 2, 3\} \xrightarrow{\simeq} \{c_1, c_2, c_3\}, \varphi)]$  and  $[(\{1, 2, 3\} \xrightarrow{\simeq} \{c_3, c_1, c_2\}, (123)\varphi)]$  are in fact the same element of  $\mathcal{O}(c_1, c_2, c_3)$ .

For reasons to become apparent when we discuss how to encode compositions in equivariant operads using equivariant trees, the fact that the labels in (4.3) change depending on the planar representation is rather inconvenient. Since the source of the problem is the use of noncanonical isomorphisms  $\mathcal{O}(S) \simeq \mathcal{O}(|S|)$ , an easy solution is to use labels in  $\mathcal{O}(S)$  instead. Thus, denoting by  $[\psi] \in \mathcal{O}(\{a_1, a_2\})$ ,  $[\chi] \in \mathcal{O}(\{b\})$ ,  $[\varphi] \in \mathcal{O}(\{c_1, c_2, c_3\})$  the classes of the operations in the leftmost tree in (4.3), the same information could then be instead encoded by labeling both trees with the common labels  $[\varphi]$ ,  $[\psi]$ ,  $[\chi]$ . In what follows, we refer to labels of the form  $\varphi \in \mathcal{O}(n)$  as *coordinate-dependent* labels and to  $[\varphi] \in \mathcal{O}(S)$  labels as *coordinate-free* labels.

Additionally, an important feature of symmetric trees not present in planar trees is that they have nontrivial “automorphism groups”. For example, the tree



has automorphism group isomorphic to the wreath product  $\Sigma_3 \wr \Sigma_2 = \Sigma_3 \times (\Sigma_2)^{\times 3}$ .

**Remark 4.4** An alternative and more rigorous perspective on (4.3) is provided by [16, Section 3], where it is explained that any tree  $T$  has an associated *colored* operad  $\Omega(T)$ . Briefly, colored operads generalize operads much in the way that categories generalize monoids: each colored operad  $\mathcal{O}$  has a collection of objects, morphism sets  $\mathcal{O}(\underline{b}; a)$  for an ordered tuple of source objects  $\underline{b} = b_1, \dots, b_n$  and target object  $a$ , units  $\text{id}_a \in \mathcal{O}(a; a)$ , compositions

$$\mathcal{O}(b_1, \dots, b_n; a) \times \mathcal{O}(\underline{c}_1; b_1) \times \dots \times \mathcal{O}(\underline{c}_n; b_n) \xrightarrow{\circ} \mathcal{O}(\underline{c}_1, \dots, \underline{c}_n; a),$$

and isomorphisms  $\mathcal{O}(b_1, \dots, b_n; a) \xrightarrow{\sigma_*} \mathcal{O}(b_{\sigma^{-1}(1)}, \dots, b_{\sigma^{-1}(n)}; a)$  for each  $\sigma \in \Sigma_n$  satisfying natural associativity, unitality and symmetry conditions. Note that regular operads are then “colored operads with a single color”.

$\Omega(T)$  is then the colored operad with objects the set of edges of  $T$  and freely generated by morphisms associated to each node. Explicitly, for the tree in (4.3) these generators are (unique) morphisms  $a_1 a_2 \rightarrow c_1$ ,  $b \rightarrow c_3$  and  $c_1 c_2 c_3 \rightarrow d$  if using the leftmost



planar representation, or generators  $a_2 a_1 \rightarrow c_1$ ,  $b \rightarrow c_3$  and  $c_3 c_1 c_2 \rightarrow d$  if using the rightmost planar representation (that the two descriptions coincide follows from the symmetry isomorphisms).

Equation (4.3) is then a diagrammatic representation of a morphism  $\Omega(T) \rightarrow \mathcal{O}$ , between the colored operad  $\Omega(T)$  and the regular operad  $\mathcal{O}$ , with the node labels being the image of the associated generators of  $\Omega(T)$ .

### 4.3 Equivariant trees

Throughout the following we fix a finite group  $G$  and the term operad will now refer to an operad  $\mathcal{O}$  together with a left  $G$ -action compatible with all the structure. Notably,  $\mathcal{O}(n)$  now has a  $(G \times \Sigma_n)$ -action, so that from an equivariant homotopy theory perspective it is natural to consider fixed points  $\mathcal{O}(n)^\Gamma$  for  $\Gamma \leq G \times \Sigma_n$ . On the other hand, since operad theory often focuses on  $\Sigma$ -cofibrant operads (ie such that  $\mathcal{O}(n)$  is  $\Sigma_n$ -free), it is natural to focus attention on  $\Gamma$  such that  $\Gamma \cap \Sigma_n = *$ , since for such  $\mathcal{O}$  we have that  $\mathcal{O}(n)^\Gamma = \emptyset$  otherwise. The key identifications

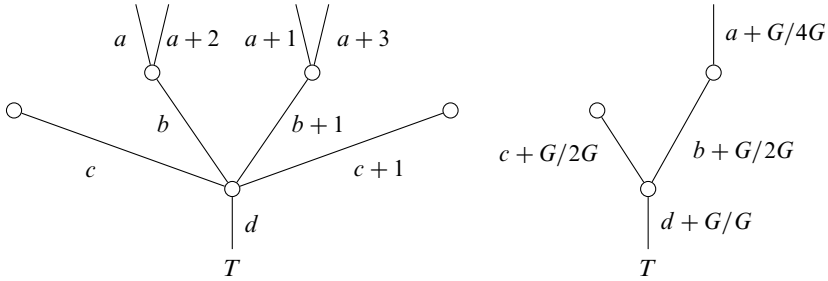
$\Gamma$  such that  $\Gamma \cap \Sigma_n = * \leftrightarrow$  graph of  $G \geq H \rightarrow \Sigma_n \leftrightarrow$  action of  $H \leq G$  on  $\{1, \dots, n\}$

then hint at a deep connection between  $G$ -operads and  $H$ -sets that is at the core of Blumberg and Hill's work in [2]. Briefly, to each  $\Sigma$ -cofibrant  $G$ -operad  $\mathcal{O}$  they associate the family of those  $\Gamma$  such that  $\mathcal{O}(n)^\Gamma \neq \emptyset$  and, in turn, the family of the corresponding  $H$ -sets,  $H \leq G$  [2, Definition 4.5]. They then show that such families satisfy a number of novel closure conditions [2, Lemmas 4.10, 4.11, 4.12, 4.15], and dub such a family an *indexing system* [2, Definition 3.22]. Moreover, analyzing their proofs one sees that the key idea is that carefully chosen fixed-point conditions on the source of the composition (4.1) induce fixed-point conditions on its target (for an explicit example, see (4.11) below).

The discovery of equivariant trees was the result of an attempt to encode the closure conditions of Blumberg and Hill diagrammatically, and we provide more details on how that works in Section 9. For now, however, we focus on examples.

As a first guess, one might attempt to define  $G$ -trees simply as symmetric trees together with a  $G$ -action (using the automorphisms mentioned in the previous section). As it turns out, such "trees with a  $G$ -action" are only a part of what is required, though we will choose such trees as our first examples.

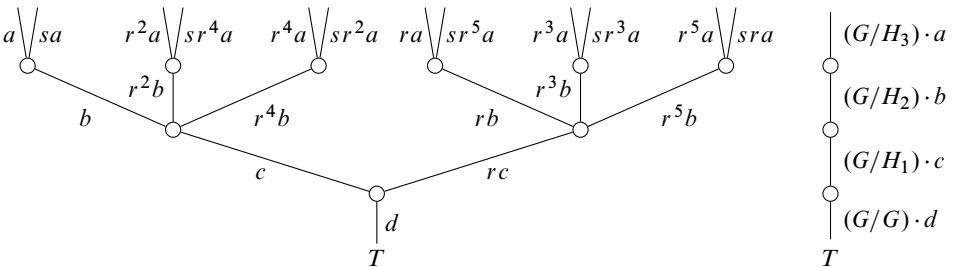
**Example 4.5** Let  $G = \mathbb{Z}/4$ . The following are two *equivalent* representations of a symmetric tree  $T$  with a  $G$ -action:



The leftmost representation, which we call the *expanded representation*, is simply a planar representation of the corresponding equivariant tree, together with a naming convention for the edges that reflects the  $G$ -action. More concretely,  $1 \in G$  acts on the tree by sending  $a$  to  $a + 1$ ,  $a + 1$  to  $a + 2$ ,  $b$  to  $b + 1$ , etc (note that implicitly  $b + 2 = b$ ,  $c + 2 = c$ ,  $d + 1 = d$ ).

The rightmost representation, which we call the *orbital representation*, is obtained from the expanded representation by “identifying edges that lie in the same  $G$ -orbit” and then labeling the corresponding “edge orbit” by the  $G$ -set of the edges corresponding to it.

**Example 4.6** Let  $G = D_6 = \{e, r, r^2, r^3, r^4, r^5, s, sr, sr^2, sr^3, sr^4, sr^5\}$  denote the hexagonal dihedral group with generators  $r$  and  $s$  such that  $r^6 = e$ ,  $s^2 = e$ ,  $srs = r^5$ . Letting  $H_1 \geq H_2 \geq H_3$  denote the subgroups  $H_1 = \langle r^2, s \rangle$ ,  $H_2 = \langle s \rangle$ ,  $H_3 = \{e\}$ , one has the following representations of a tree  $T$  with  $G$ -action:



We note that it is implicit in the orbital representation that, for example, the assignment  $b \mapsto c$  defines a  $G$ -set map  $(G/H_2) \cdot b \rightarrow (G/H_1) \cdot c$  (ie  $H_1 \geq H_2$ ).

We can now ask what the analogue of the node labels in (4.2) are for such  $G$ -trees. For example, for  $G = \mathbb{Z}/4$  consider the label  $\varphi$  in the leftmost expanded representation of the following equivariant *corolla* (ie tree with a single node):

(4.7)

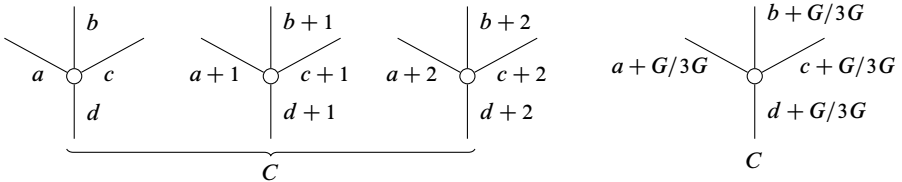
An immediate answer is provided by Remark 4.4: indeed, the corolla  $C$  with a  $G$ -action will generate a colored operad  $\Omega(C)$  with a  $G$ -action, so that (4.7) should encode a  $G$ -equivariant map  $\Omega(C) \rightarrow \mathcal{O}$ . Unpacking this observation, one should have  $\varphi \in \mathcal{O}(4)$ , but an additional equivariance condition is to be expected. To make this explicit, note first that there are left actions of both  $G$  and  $\Sigma_4$  on the set of all morphisms of  $\Omega(C)$  and that these actions commute, assembling to an action of  $G \times \Sigma_4$ . As concrete examples,  $1 \in \mathbb{Z}/4$  sends the morphism  $cb(b+1)(c+1) \rightarrow d$  to  $(c+1)(b+1)bc \rightarrow d$  while  $(124) \in \Sigma_4$  sends  $cb(b+1)(c+1) \rightarrow d$  to  $(c+1)c(b+1)b \rightarrow d$ . Further, one can readily check that the  $(G \times \Sigma_4)$ -isotropy of the morphism  $cb(b+1)(c+1) \rightarrow d$  is precisely the subgroup  $\Gamma_{\{c,b,b+1,c+1\}}$  given by the graph of the homomorphism  $G \rightarrow \Sigma_4$  encoding the  $G$ -set  $\{c, b, b+1, c+1\}$ . And since  $\varphi$  is the image of that morphism, we get the sought condition  $\varphi \in \mathcal{O}(4)^{\Gamma_{\{c,b,b+1,c+1\}}}$ .

We now turn our attention to the orbital representation of  $C$  on the right side of (4.7), which is often preferable both for conceptual reasons and compactness. Writing  $S = \{b, c, b+1, c+1\}$ , our node label is now the coordinate-free label  $[\varphi] \in \mathcal{O}(S)$  (indeed, a label in  $\mathcal{O}(4)$  can not be used since the orbital notation provides no ordering of  $S$ ). Further,  $\mathcal{O}(S)$  now has two commuting  $G$ -actions, one induced by the structural  $G$ -action on  $\mathcal{O}$  and one induced by the  $G$ -action on  $S$ . Referring to the combined diagonal  $G$ -action as the canonical  $G$ -action, the equivariance condition is now straightforward: it is simply  $[\varphi] \in \mathcal{O}(S)^G$ .

It seems helpful to make the equivalence between the previous two paragraphs explicit. The homomorphism  $\rho: G \rightarrow \Sigma_S$  encoding  $S$  induces a semidirect product  $G \ltimes \Sigma_S$  with multiplication  $(g, \sigma)(\bar{g}, \bar{\sigma}) = (g\bar{g}, \sigma\rho(g)\bar{\sigma}\rho(g)^{-1})$ . The semidirect product  $G \ltimes \Sigma_S$  naturally acts on  $S$  with the action of  $(g, \sigma)$  given by  $\sigma\rho(g)$  and on  $\mathcal{O}(S)$  with action given by  $\sigma g$  (where  $\sigma \in \Sigma_S$  acts via symmetries and  $g \in G$  acts via the canonical action). The isomorphism  $\tau: \{c, b, b+1, c+1\} \xrightarrow{\cong} \{1, 2, 3, 4\}$  then induces an isomorphism  $\tau_*: G \ltimes \Sigma_S \xrightarrow{\cong} G \ltimes \Sigma_4$  via  $\tau_*(g, \sigma) = (g, \tau\sigma\rho(g)\tau^{-1})$ . That the previous two paragraphs are equivalent is then simply the observation that  $\tau_*$  sends the subgroup  $G \leq G \ltimes \Sigma_S$  to  $\Gamma_{\{c,b,b+1,c+1\}} \leq G \ltimes \Sigma_4$ , ie that the graph subgroup encodes the canonical action when in coordinate-dependent notation.

Now that we know that corollas with  $G$ -actions encode operations fixed by graphs of full homomorphisms  $G \rightarrow \Sigma_n$ , we turn to the question of how to encode operations fixed by graphs of *partial* homomorphisms  $G \geq H \rightarrow \Sigma_n$ . A natural first guess might be that this role is played by corollas with an  $H$ -action. However, due to the lack of full  $G$ -actions this would not quite provide the necessary maps for the category  $\Omega_G$  of  $G$ -trees that we introduce in Section 5.3. The solution is both simple and surprising: one simply “induces a tree with an  $H$ -action into a  $G$ -object”. We start with an example where  $H = *$ .

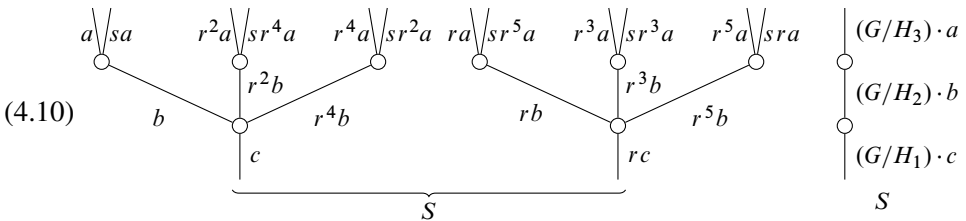
**Example 4.8** Let  $G = \mathbb{Z}/3$ . The equivariant corolla  $C$  with orbital representation given on the right



has expanded representation given by the union of the three (nonequivariant) corollas on the left. For clarity, we stress that we refer to the three trees on the left together as a forming a single  $\mathbb{Z}/3$ -tree. The legitimacy of this nomenclature is born out of the role such  $G$ -trees play in the theory, though for now we simply point out that at least the orbital representation is an “honest” tree (for further discussion, see Section 5.3).

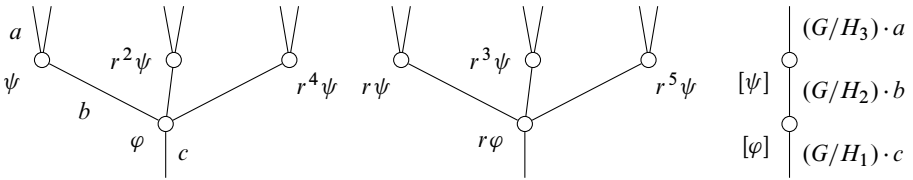
A map  $\Omega(C) \rightarrow \mathcal{O}$  is then determined by the image of the morphism  $abc \rightarrow d$ , and hence by an arbitrary operation  $[\varphi] \in \mathcal{O}(\{a, b, c\}) \simeq \mathcal{O}(3)$ , which determines operations  $[\varphi] + 1 \in \mathcal{O}(\{a + 1, b + 1, c + 1\})$  and  $[\varphi] + 2 \in \mathcal{O}(\{a + 2, b + 2, c + 2\})$ .

**Example 4.9** Keeping  $G = D_6$  and  $H_1, H_2, H_3$  as in Example 4.6, removing the *root* orbit (ie bottom orbit) from the  $G$ -tree  $T$  therein yields the  $D_6$ -tree



We end this introduction by illustrating the kind of compositions that  $G$ -trees encode.

Taking the  $D_6$ -tree  $S$  from [Example 4.9](#), a map  $\Omega(S) \rightarrow \mathcal{O}$  leads to node labels



where  $[\psi] \in \mathcal{O}((H_2/H_3) \cdot a)^{H_2}$  and  $[\phi] \in \mathcal{O}((H_1/H_2) \cdot b)^{H_1}$ . We note that in particular  $r[\phi] \in \mathcal{O}(r(H_1/H_2)r^{-1} \cdot rb)^{rH_1r^{-1}} = \mathcal{O}((H_1^r/H_2^r) \cdot rb)^{H_1^r}$  and likewise for  $r\psi, r^2\psi, \dots, r^5\psi$ , so that we are adopting the convention that labels in the orbital notation are chosen according to the edge orbit generators  $a, b, c$ .

Further unpacking the map  $\Omega(S) \rightarrow \mathcal{O}$ , we see that  $S$  encodes the fact that the composition product

$$\mathcal{O}(H_1/H_2) \times \prod_{[h] \in H_1/H_2} \mathcal{O}(H_2^h/H_3^h) \rightarrow \mathcal{O}(H_1/H_3)$$

restricts to

$$(4.11) \quad \mathcal{O}(H_1/H_2)^{H_1} \times \left( \prod_{[h] \in H_1/H_2} \mathcal{O}(H_2^h/H_3^h)^{H_2^h} \right)^{H_1} \rightarrow \mathcal{O}(H_1/H_3)^{H_1}$$

or, using that  $(\prod_{[h] \in H_1/H_2} \mathcal{O}(H_2^h/H_3^h)^{H_2^h})^{H_1} \simeq \mathcal{O}(H_2/H_3)^{H_2}$ , simply

$$\mathcal{O}(H_1/H_2)^{H_1} \times \mathcal{O}(H_2/H_3)^{H_2} \rightarrow \mathcal{O}(H_1/H_3)^{H_1}.$$

## 5 Categories of trees and forests

In this section we introduce the several categories of trees, forests and presheaves we will be working with. We will make heavy use of the broad poset framework introduced by Weiss in [\[19\]](#), which provides an algebraically flavored model for the category  $\Omega$  of trees (see [\[16; 17; 5; 11\]](#), among others). We will find this particularly convenient since, when using tree diagrams as in [Section 4.3](#), representative examples of equivariant trees are typically quite large.

### 5.1 Broad posets and the category of trees

We start by recalling the key notions in [\[19\]](#) and establishing some basic results.

Given a set  $T$  we denote by  $T^+$  the free abelian monoid generated by  $T$ . Elements of  $T^+$  will be written in tuple notation, such as  $\underline{e} = e_1e_3e_1e_2 = e_1e_1e_2e_3 \in T^+$

for  $e_1, e_2, e_3 \in T$ . We will also write  $e_i \in \underline{e}$  whenever  $e_i$  is a “letter” appearing in  $\underline{e}$ ,  $\underline{f} \subset \underline{e}$  if  $\underline{f}\underline{g} = \underline{e}$  for some  $\underline{g} \in T^+$ , and denote the “empty tuple” of  $T^+$  by  $\epsilon$ .

**Definition 5.1** A (commutative) broad poset structure [19, Definition 3.2] on  $T$  is a relation  $\leq$  on  $T^+$ ,  $T$  (ie a subset of  $T^+ \times T$ ) satisfying the following conditions:

- **Reflexivity**  $e \leq e$  (for  $e \in T$ ).
- **Antisymmetry** If  $e \leq f$  and  $f \leq e$ , then  $e = f$  (for  $e, f \in T$ ).
- **Broad transitivity** If  $f_1 f_2 \cdots f_n = \underline{f} \leq e$  and  $\underline{g}_i \leq f_i$ , then  $\underline{g}_1 \cdots \underline{g}_n \leq e$  (for  $e, f_i \in T$  and  $\underline{f}, \underline{g}_i \in T^+$ ).

**Remark 5.2** Omitting antisymmetry yields the notion of a *prebroad poset*.

Since the main examples of broad posets are induced by constructions involving trees, we will refer to the elements of a broad poset as its *edges*.

**Definition 5.3** A broad poset  $P$  is called *simple* if for any broad relation  $e_1 \cdots e_n \leq e$  one has  $e_i = e_j$  only if  $i = j$ .

**Notation 5.4** A broad poset structure  $\leq$  on  $T$  naturally induces the following preorder relations on  $T$  and  $T^+$ :

- For  $f, e \in T$  we say that  $f$  is a *descendant* of  $e$ , written  $f \leq_d e$ , if there exists a broad relation  $\underline{f} \leq e$  such that  $f \in \underline{f}$ .
- For  $\underline{f}, \underline{e} \in T^+$ , we write  $\underline{f} \leq \underline{e}$  if it is possible to write  $\underline{f} = \underline{f}_1 \cdots \underline{f}_k$  and  $\underline{e} = e_1 \cdots e_k$  such that  $\underline{f}_i \leq e_i$  for  $i = 1, \dots, k$ .

**Remark 5.5** Generally, these preorders can be fairly counterintuitive. For example, it is possible to have  $ab \leq a$ , or even both  $aa \leq a$  and  $a \leq aa$  simultaneously. The case of simple broad posets, however, is much simpler.

**Proposition 5.6** Let  $T$  be a simple broad poset. Then  $\leq_d$  (resp.  $\leq$ ) is an order relation on  $T$  (resp. on  $T^+$ ). Further, if  $f_1 \cdots f_k \leq e$  then the  $f_i$  are  $\leq_d$ -incomparable (in particular,  $e \underline{f} \leq e$  only if  $\underline{f} = \epsilon$ ).

**Proof** The “further” part is immediate: if two  $f_i$  were  $\leq_d$ -comparable then broad transitivity would produce a nonsimple broad relation.

To see that  $\leq_d$  satisfies antisymmetry, note that if  $e' \underline{f} \leq e$  and  $e \underline{g} \leq e'$ , then  $e \underline{g} \underline{f} \leq e$ , so it must be that  $\underline{g} = \underline{f} = \epsilon$  and the antisymmetry of  $\leq$  on  $T$  implies  $e = e'$ .

Finally, we show antisymmetry of  $\leq$  on  $T^+$  by induction on the size of the tuple  $\underline{e}$  in a pair of relations  $\underline{f} \leq \underline{e}$  and  $\underline{e} \leq \underline{f}$ . The  $\underline{e} = \epsilon$  case is immediate. Otherwise let  $e \in \underline{e}$  be  $\leq_d$ -maximal and choose  $\underline{eg} \subset \underline{e}$  and  $f \in \underline{f}$  such that  $\underline{eg} \leq f$  and choose  $\underline{hf} \subset \underline{f}$  and  $e' \in \underline{e}$  such that  $\underline{hf} \leq e'$ . Then  $e \leq_d f \leq_d e'$  and by  $\leq_d$ -maximality of  $e$  it must hold that  $e = f = e'$  and hence, by the “further” claim, also that  $\underline{g} = \epsilon = \underline{h}$ . And since this must hold regardless of how  $f, e', \underline{g}, \underline{h}$  are chosen, one concludes that, writing  $\underline{e} = \underline{ee'}$  and  $\underline{f} = \underline{ef'}$ , it must in fact also hold that  $\underline{e'} \leq \underline{f'}$  and  $\underline{f'} \leq \underline{e'}$ ; thus the induction hypothesis applies.  $\square$

**Definition 5.7** An edge  $e \in T$  is called

- a *leaf* if there are no  $\underline{f} \in T^+$  such that  $\underline{f} < e$  (ie  $\underline{f} \leq e$  and  $\underline{f} \neq e$ );
- a *node* if there is a nonempty maximum  $\underline{f} \neq \epsilon$  such that  $\underline{f} < e$ ;
- a *stump* if  $\underline{f} = \epsilon$  is the maximum (in fact, only)  $\underline{f}$  such that  $\underline{f} < e$ .

Further, in either the node or stump case the maximum such  $\underline{f}$  is denoted by  $e^\uparrow$ .

**Remark 5.8** While it is customary to regard stumps simply as a type of node, we will find it convenient, in light of Proposition 7.15 and Lemma 7.22, to separate the two cases.

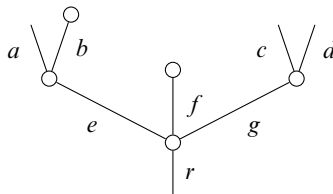
The following definition is the key purpose of [19].

**Definition 5.9** A *dendroidally ordered set* is a finite simple broad poset  $T$  satisfying the following additional conditions:

- **Nodal condition** Each edge  $e \in T$  is either a leaf, a node or a stump.
- **Root condition** There is a maximum  $r_T \in T$  for  $\leq_d$ , called the *root* of  $T$ .

Weiss proves in [19] that the category of dendroidally ordered sets (together with the obvious notion of monotonous function) is equivalent to the category  $\Omega$  of trees (see [16; 17; 5; 11] etc). As such, we will henceforth refer to dendroidally ordered sets simply as *trees* and use them as our model for  $\Omega$ .

**Example 5.10** The tree diagram



represents a broad poset structure on  $\{a, b, c, d, e, f, g, r\}$ . The nodes represent generating broad relations  $\epsilon \leq b$ ,  $ab \leq e$ ,  $\epsilon \leq f$ ,  $cd \leq g$  and  $efg \leq r$  with the other broad relations, such as  $afcd \leq r$ , obtained by “composition” (ie using broad transitivity). Note that, alternatively, one can also write  $b^\uparrow = \epsilon$ ,  $e^\uparrow = ab$ ,  $f^\uparrow = \epsilon$ ,  $g^\uparrow = cd$  and  $r^\uparrow = efg$  to denote the generating broad relations.

We will make use of the following basic results.

**Proposition 5.11** *Let  $T$  be a tree and  $A$  any broad poset. A set map  $\varphi: T \rightarrow A$  is a broad poset map if and only if  $\varphi(e^\uparrow) \leq \varphi(e)$  for each node/stump  $e \in T$ .*

*In particular, the broad relations of  $T$  are generated by the  $e^\uparrow \leq e$  relations.*

**Proof** Since for any nonidentity relation  $\underline{f} < e$  one has  $\underline{f} \leq e^\uparrow < e$ , one can write  $e^\uparrow = e_1 \cdots e_k$  and  $\underline{f} = \underline{f}_1 \cdots \underline{f}_k$  so that  $\underline{f}_i \leq e_i$  ( $k = 0$  is allowed, in which case  $e^\uparrow = \underline{f} = \epsilon$ ), so the result follows by  $\leq_d$  induction on  $e$ .

The “in particular” claim follows from the identity map  $\text{id}_T: T \rightarrow \tilde{T}$ , where  $\tilde{T}$  has the broad poset structure generated by the  $e^\uparrow \leq e$ . □

**Lemma 5.12** *Let  $T$  be a tree. For any  $e \in T$  there exists a minimum  $e^\lambda \in T^+$  such that  $e^\lambda \leq e$ . In fact,  $e^\lambda = l_1 \cdots l_k$  consists of those leaves  $l_i$  such that  $l_i \leq_d e$ .*

*Further, a broad relation  $\underline{f} = f_1 \cdots f_n \leq e$  holds if and only if  $f_i \leq_d e$ , the  $f_i$  are  $\leq_d$ -incomparable and  $f_1^\lambda \cdots f_n^\lambda = e^\lambda$ .*

**Proof** The proof is by  $\leq_d$  induction on  $e$ . The leaf case is obvious. Otherwise, let  $\underline{f} \leq e^\uparrow < e$  be any nonidentity relation and write  $e^\uparrow = e_1 \cdots e_k$  and  $\underline{f} = \underline{f}_1 \cdots \underline{f}_k$  so that  $\underline{f}_i \leq e_i$  ( $k = 0$  is allowed). By induction,  $e_i^\lambda \leq \underline{f}_i \leq e_i$  where  $e_i^\lambda$  consists of the leaves  $l$  such that  $l \leq_d e_i$  and hence indeed we have that  $e^\lambda = e_1^\lambda \cdots e_k^\lambda \leq \underline{f}$ .

Only the “if” half of the “further” statement needs proof. We use the same induction argument: incomparability yields  $e \notin \underline{f}$  provided  $\underline{f} \neq e$  and, writing  $e^\uparrow = e_1 \cdots e_k$  and  $\underline{f} = \underline{f}_1 \cdots \underline{f}_k$  so that  $s \in \underline{f}_i$  if and only if  $s \leq_d e_i$ , the induction hypothesis applies. □

**Example 5.13** In Example 5.10,  $e^\lambda = a$ ,  $g^\lambda = cd$ ,  $f^\lambda = \epsilon$  and  $r^\lambda = acd$ .

We now discuss a key operation on trees: grafting. We recall that  $\eta$  denotes the tree with a single edge, also denoted by  $\eta$ , and only the identity relation  $\eta \leq \eta$ .



**Definition 5.14** Let  $T, U \in \Omega$  be trees and let  $v$  denote both a leaf  $v: \eta \rightarrow T$  and the root  $v: \eta \rightarrow U$ . The *grafting*  $T \amalg_v U$  is the pushout (of prebroad posets):

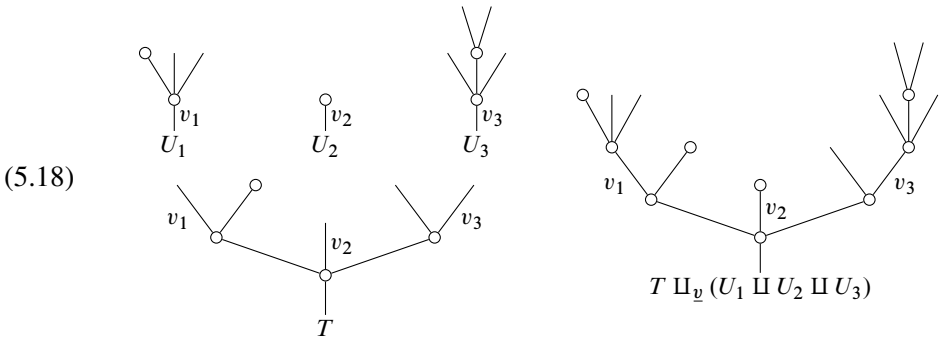
$$\begin{array}{ccc} \eta & \xrightarrow{v} & U \\ v \downarrow & & \downarrow \\ T & \longrightarrow & T \amalg_v U \end{array}$$

**Proposition 5.15**  $T \amalg_v U$  is a tree.

**Proof** The underlying set of  $T \amalg_v U$  is the underlying coproduct and the broad relations are easily seen to come in three types: (i)  $\underline{t} \leq t$ , a relation in  $T$ ; (ii)  $\underline{u} \leq u$ , a relation in  $U$ ; (iii)  $\underline{t}\underline{u} \leq t$  whenever both  $\underline{t}v \leq t$  in  $T$  and  $\underline{u} \leq v$  in  $U$ . The antisymmetry, simple, nodal and root conditions are straightforward.  $\square$

**Remark 5.16** More generally, letting  $\underline{v} = v_1 \cdots v_n$  denote both a tuple of leaves of  $T$  and the roots  $v_i \in U_i$  we similarly define a grafted tree  $T \amalg_{\underline{v}} (U_1 \amalg \cdots \amalg U_n)$ . Explicitly, its broad relations have the form (i)  $\underline{t} \leq t$ , a relation in  $T$ ; (ii)  $\underline{u} \leq u$ , a relation in some  $U_i$ ; (iii)  $\underline{t}u_{k_1} \cdots u_{k_p} \leq t$  whenever  $\{k_i\} \subset \{1, \dots, n\}$  and both  $\underline{t}v_{k_1} \cdots v_{k_p} \leq t$  in  $T$  and  $\underline{u}_{k_i} \leq v_{k_i}$  in  $U_i$ .

**Example 5.17** A tree representation of the grafting procedure follows:



We will also find it useful to be able to reverse the grafting procedure.

For any tree  $T$  and edge  $e \in T$  we will let  $T^{\leq e}$  denote the prebroad poset with set  $\{f \in T \mid f \leq_d e\}$  and the generating relations  $f^\uparrow \leq f$  for  $f \leq_d e$ .

Similarly, for a  $\leq_d$ -incomparable tuple  $\underline{e} = e_1 \cdots e_n$  we will let  $T_{\not\leq \underline{e}}$  denote the prebroad poset with set  $\{f \in T \mid f \not\leq_d e_i \text{ for all } i\}$  and the generating relations  $f^\uparrow \leq f$  for  $f$  satisfying  $f \not\leq_d e_i$  and  $f \neq e_i$  for all  $i$ . Note that by incomparability we have that  $e_i \in T_{\not\leq \underline{e}}$ .

**Proposition 5.19**  $T^{\leq e}$  and  $T_{\neq e}$  are trees. Hence, for any  $\leq_d$ -incomparable tuple  $\underline{e} = e_1 \cdots e_n$ , one has

$$(5.20) \quad T \simeq T_{\neq \underline{e}} \amalg_{\underline{e}} (T^{\leq e_1} \amalg \cdots \amalg T^{\leq e_n}).$$

**Proof** In both cases antisymmetry and simplicity are inherited from  $T$ . Further, since any string  $f_1 \leq_d f_2 \leq_d \cdots \leq_d f_n$  in  $T$  (note that we can assume the  $\leq_d$  are induced by generating relations) where  $f_1, f_n \in T^{\leq e}$  (resp.  $f_1, f_n \in T_{\neq e}$ ) is a string in  $T^{\leq e}$  (resp.  $T_{\neq e}$ ), the nodal and root conditions also follow.

Equation (5.20) follows from the “in particular” claim in Proposition 5.11. □

**Example 5.21** Denoting by  $V$  the rightmost tree in (5.18), one has  $U_i = V^{\leq v_i}$  and  $T = V_{\neq \underline{v}}$ . We note here a pictorial mnemonic for our index/exponent notation:  $V^{\leq v_i}$  denotes an “upper subtree” of  $V$  while  $V_{\neq \underline{v}}$  denotes a “lower subtree”.

**Remark 5.22** The leaves of  $T_{\neq \underline{e}}$  consist of the edges  $e_i \in \underline{e}$  together with the leaves of  $T$  not in  $\underline{e}^\lambda$ , ie those leaves  $\leq_d$ -incomparable with the  $e_i \in \underline{e}$ .

**Corollary 5.23** Let  $U \xrightarrow{\varphi} T$  be a map in  $\Omega$  and let  $r$  be the root and  $\underline{l}$  be the leaf tuple of  $U$ . Then  $\varphi$  naturally factors as  $U \rightarrow T_{\neq \varphi(\underline{l})}^{\leq \varphi(r)} \hookrightarrow T$ .

**Proof** By the “in particular” claim in Proposition 5.11 it suffices to check that any relation  $\underline{e} \leq^T e$  in  $T$  is a relation in  $T_{\neq \underline{e}}^{\leq e}$ . Since the root of  $T_{\neq \underline{e}}^{\leq e}$  is  $e$  and the leaf tuple is  $\underline{e}$ , this follows by Lemma 5.12. □

**Lemma 5.24** Let  $T$  be a tree and  $\underline{e}$  a tuple of  $\leq_d$ -incomparable edges of  $T$ . Then, letting  $r$  be the root of  $T$ , there exists a broad relation of the form  $\underline{e} \underline{f} \leq r$ .

**Proof** The proof is by induction on the sum of the distances (measured in  $\leq_d$  inequality chains) from the edges in  $\underline{e}$  to the root. Clearly  $r \in \underline{e}$  only if  $r = \underline{e}$ . Otherwise one can write  $r^\uparrow = r_1 \cdots r_k$  and  $\underline{e} = \underline{e}_1 \cdots \underline{e}_k$  so that  $s \in \underline{e}_i$  if and only if  $s \leq_d r_i$ . The induction hypothesis applies to  $\underline{e}_i$  and the subtrees  $T^{\leq r_i}$ . □

Since the converse of the previous lemma holds by Proposition 5.6, we obtain the following.

**Corollary 5.25** If  $e, f \in T$  are  $\leq_d$ -incomparable and  $e' \leq_d e$ , then  $e', f$  are  $\leq_d$ -incomparable.

### 5.2 Categories of forests

We will also need to discuss forests, ie formal coproducts of trees. We start by generalizing Definition 5.9.

**Definition 5.26** A forestally ordered set is a finite simple broad poset  $F$  satisfying the following additional conditions:

- **Nodal condition** Each edge  $e \in F$  is either a leaf, a node or a stump.
- **Root tuple condition** For each edge  $e$  there is a unique  $\leq_d$ -maximal element  $r \in T$  such that  $e \leq_d r$ . Further, we denote by  $\underline{r}_F$  the tuple of  $\leq_d$ -maximal elements of  $F$  and refer to it as the *root tuple*.

The last condition guarantees that any forestally ordered set decomposes as a disjoint union of dendroidally ordered sets, ie trees. We shall hence refer to these simply as *forests* and denote by  $\Phi$  the category formed by them.

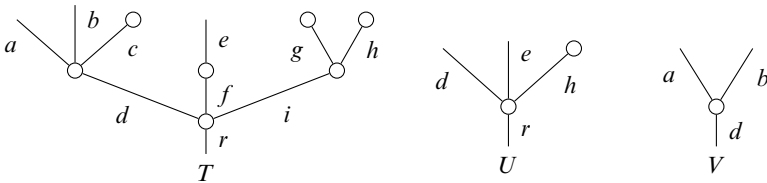
**Definition 5.27** A map of forests  $F \xrightarrow{\varphi} F'$  is called

- *wide* if  $\varphi(\underline{r}_F) \leq \underline{r}_{F'}$ ;
- *independent* if there exists a tuple  $\underline{e} \in F^+$  such that  $\varphi(\underline{r}_F)\underline{e} \leq \underline{r}_{F'}$ .

The subcategory of forests and independent (resp. wide) maps is denoted by  $\Phi_w$  (resp.  $\Phi_i$ ). Note that there are inclusions  $\Phi_w \hookrightarrow \Phi_i \hookrightarrow \Phi$ .

**Remark 5.28** The category  $\Phi_i$  nearly coincides with the category of forests discussed in [11, Section 3.1], the only difference being that here we include the empty forest  $\emptyset$ .

**Example 5.29** Consider the following tree  $T$  and subtrees (all labels are on edges):



Further denoting by  $\eta_a, \eta_b, \dots$  the single-edge subtrees corresponding to each edge, some examples of wide morphisms are given by the inclusions

$$\begin{aligned}
 &V \hookrightarrow T, & V \amalg \eta_f \hookrightarrow T, & V \amalg \eta_e \hookrightarrow T, \\
 &V \amalg \eta_f \amalg \eta_i \hookrightarrow T, & V \amalg \eta_f \amalg \eta_g \hookrightarrow T, & \eta_a \amalg \eta_b \amalg \eta_f \hookrightarrow T,
 \end{aligned}$$

some examples of nonwide independent maps are given by

$$V \sqcup \eta_g \hookrightarrow T, \quad \eta_f \sqcup \eta_i \hookrightarrow T, \quad \eta_g \sqcup \eta_h \hookrightarrow T, \quad \emptyset \hookrightarrow T,$$

and examples of nonindependent maps are given by

$$V \sqcup \eta_r \rightarrow T, \quad V \sqcup \eta_f \sqcup \eta_i \sqcup \eta_g \rightarrow T, \quad \eta_i \sqcup \eta_i \rightarrow T.$$

The following is a useful forestal strengthening of [Proposition 5.6](#), which follows by combining that result with [Lemma 5.24](#).

**Proposition 5.30** *Let  $F$  be a forest. If  $\underline{f} \leq \underline{e} = e_1 \cdots e_k$  with  $\leq_d$ -incomparable  $e_i$ , then the decomposition  $\underline{f} = \underline{f}_1 \cdots \underline{f}_k$  with  $\underline{f}_i \leq e_i$  is unique. In fact, for  $s \in \underline{f}$  one has  $s \in \underline{f}_i$  if and only if  $s \leq_d e_i$ .*

**Proof** If  $s \leq_d e_i$  and  $s \leq_d e_j$  with  $i \neq j$ , the unique root  $r$  such that  $s \leq_d r$  must also satisfy  $e_i \leq_d r$  and  $e_j \leq_d r$  and [Lemma 5.24](#) can be used to produce a nonsimple relation. Thus such  $e_i$  is unique and the result follows. □

We now turn to describing the degeneracy-face decomposition of maps in the broad poset setting, which we obtain as [Proposition 5.37](#) below.

**Definition 5.31** A map  $F \xrightarrow{\varphi} F'$  in  $\Phi_i$  is called

- a *face map* if the underlying set map is injective;
- a *degeneracy* if the underlying set map is surjective and for each leaf  $l \in F$  we have that  $\epsilon \not\leq \varphi(l)$ .

**Lemma 5.32** *For any map  $F \xrightarrow{\varphi} F'$  in  $\Phi_i$ , we have that  $e, \bar{e}$  are  $\leq_d$ -comparable if and only if  $\varphi(e), \varphi(\bar{e})$  are  $\leq_d$ -comparable.*

*Further, if  $\varphi(e) = \varphi(\bar{e})$ , then  $\bar{e}\underline{f} \leq e$  can hold only if  $\underline{f} = \epsilon$  and thus either  $e \leq \bar{e}$  or  $\bar{e} \leq e$ .*

**Proof** If  $e, \bar{e}$  are not  $\leq_d$ -comparable, [Lemma 5.24](#) ensures that there exists  $\underline{f}$  such that  $e\bar{e}\underline{f} \leq r_F$ . Thus, by definition of  $\Phi_i$ , there exists  $\underline{g}$  such that  $\varphi(e)\varphi(\bar{e})\varphi(\underline{f})\underline{g} \leq \varphi(r_{F'})\underline{g} \leq r_{F'}$ , hence  $\varphi(e), \varphi(\bar{e})$  are  $\leq_d$ -incomparable.

The “further” claim is immediate. □

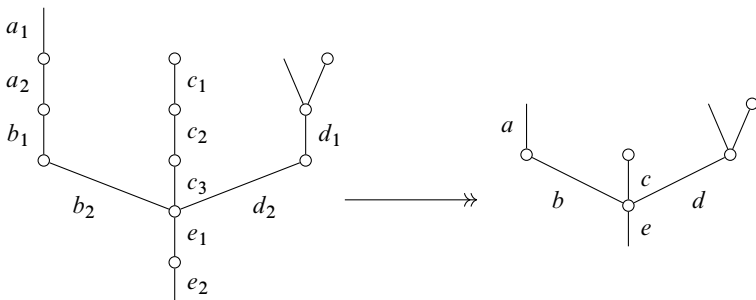
**Lemma 5.33** *If  $F \xrightarrow{\varphi} F'$  is both a face map and a degeneracy then  $\varphi$  is an isomorphism.*

**Proof** Bijectivity allows us to assume that the underlying sets are the same and it thus follows from Lemma 5.32 that the relations  $\leq_d^F$  and  $\leq_d^{F'}$  coincide. Hence, both forest structures have the same roots (ie  $\leq_d$ -maximal edges) and one needs only show that broad relations coincide in each of the constituent trees. But noting that a leaf  $l$  is precisely a  $\leq_d$ -minimal edge such that  $\epsilon \not\leq l$ , the definition of degeneracy implies that  $F$  and  $F'$  have the same leaves and the result follows by the criterion in the “further” part of Lemma 5.12.  $\square$

**Lemma 5.34** Let  $F \xrightarrow{\varphi} F'$  be a map in  $\Phi_i$  and  $\underline{f} \leq e$  be a broad relation in  $F$  such that  $\varphi(\underline{f}) \neq \varphi(e)$ . Then for any  $\underline{f}'$  and  $e'$  such that  $\varphi(\underline{f}') = \varphi(\underline{f})$  and  $\varphi(e') = \varphi(e)$ , we also have that  $\underline{f}' \leq e'$ .

**Proof** Writing  $\underline{f} = f_1 \cdots f_n$  and  $\underline{f}' = f'_1 \cdots f'_n$ , the condition  $\varphi(\underline{f}) \neq \varphi(e)$  ensures  $f_i <_d e$ , and thus Lemma 5.32 implies that  $f'_i <_d e'$  also. Using the characterization in Lemma 5.12, the desired relation  $\underline{f}' \leq e'$  will hold provided we show that  $a^\lambda = \bar{a}^\lambda$  whenever  $\varphi(a) = \varphi(\bar{a})$ . But since leaves are  $\leq_d$ -minimal, this too follows from Lemma 5.32.  $\square$

**Remark 5.35** It follows from the “further” part in Lemma 5.32 together with anti-symmetry that the preimage of any edge by a map  $\varphi$  always consists of a linearly ordered subset of edges. As such, degeneracies are necessarily maps that “collapse linear sections of a tree”, and indeed that was their original description in [19]. A typical degeneracy (sending edges labeled  $a_i, b_i, c_i, d_i$  to respective edges  $a, b, c, d$ ) is pictured below:



**Lemma 5.36** Let  $\varphi: F \rightarrow F'$  be a map in  $\Phi_i$  and let  $U \subset F$  be any sub-broad poset consisting of exactly one edge in each preimage of  $\varphi$  and the broad relations of  $F$  between them. Then  $U$  is a forest.

**Proof** Simplicity of  $U$  is inherited from  $F$ . Given  $u \in U$ , let  $\bar{u} \in F$  be the  $\leq_d$ -minimal edge such that  $\varphi(\bar{u}) = \varphi(u)$ . Lemma 5.34 implies that  $u$  is a leaf in  $U$  if  $\bar{u}$  is a leaf in  $F$ . Otherwise,  $\bar{u}^\uparrow < \bar{u}$  (in  $F$ ) and again by Lemma 5.34 the unique tuple  $\underline{v}$  of  $U$  such that  $\varphi(\underline{v}) = \varphi(\bar{u}^\uparrow)$  provides the desired node tuple  $u^\uparrow = \underline{v}$  for  $U$ . Lastly, by Lemma 5.34,  $s$  will be a root of  $U$  if and only if  $\varphi(s) = \varphi(r)$  for  $r$  a root of  $F$ .  $\square$

We now prove the following factorization result, which in the  $\Omega$  case first appeared as [16, Lemma 3.1] and in the  $\Phi_i$  case was proven in [11, Lemma 3.1.3].

**Proposition 5.37** *Each map  $\varphi$  of  $\Phi_i$  has a factorization  $\varphi = \varphi^+ \circ \varphi^-$  as a degeneracy followed by a face. Further, this decomposition is unique up to unique isomorphism.*

*Finally, the decomposition restricts to the subcategories  $\Omega$  and  $\Phi_w$ .*

**Proof** Given a map  $F \xrightarrow{\varphi} F'$  and picking any  $U$  as in Lemma 5.36, the isomorphism  $U \simeq \varphi(F)$  allows us to equip  $\varphi(F)$  with a forest structure. Moreover, by Lemma 5.34 the broad relations of  $\varphi(F)$  are exactly the image of those in  $F$ , and thus independent of the chosen  $U$ . The existence of a factorization follows.

The uniqueness of the  $F \xrightarrow{\varphi^-} G \xrightarrow{\varphi^+} F'$  factorization follows since by the description of the broad structure on  $\varphi(F)$  there is clearly a broad poset map  $\varphi(F) \rightarrow G$ , which by Lemma 5.33 is an isomorphism.

Further, the decomposition restricts to  $\Omega$  since the image by a degeneracy map of a tree is necessarily a tree and restricts to  $\Phi_w$  since any edge surjective map  $F \xrightarrow{\varphi^-} G$  must map roots to roots and hence  $\varphi^-(r_F) = r_G$  must hold.  $\square$

**Corollary 5.38** *If  $F \xrightarrow{\rho} F'$  is a degeneracy in any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$  then the broad relations in  $F'$  are precisely the image of the broad relations of  $F$ . Further, any section  $F' \xrightarrow{s} F$  of the underlying set map is a section in  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$ .*

The following will be needed in Section 5.4 when discussing dendroidal boundaries.

**Corollary 5.39** *In any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$ , pairs of degeneracies with common domain have absolute pushouts.<sup>1</sup>*

<sup>1</sup>Recall that an absolute colimit can be described as either a colimit that is preserved by the Yoneda embedding or (equivalently) a colimit that is preserved by any functor.

**Proof** We will throughout write  $\Xi$  for any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$  and  $\Xi[F] \in \text{Set}^{\Xi^{\text{op}}}$  for the presheaf represented by  $F \in \Xi$ . Given a diagram of degeneracies  $E \xleftarrow{\bar{\rho}} F \xrightarrow{\rho} E_0$  we first inductively extend it to a diagram of degeneracies as on the left below

$$(5.40) \quad \begin{array}{ccccccc} F & \xrightarrow{\rho} & E_0 & \xrightarrow{\rho} & E_1 & \xrightarrow{\rho} & \cdots & \xrightarrow{\rho} & E_n \\ \downarrow \bar{\rho} & & & & & & & & \downarrow \bar{\rho} \\ \bar{E} & \nearrow & & & & & & & E_{i+1} \end{array} \quad \begin{array}{ccc} E_i & \xrightarrow{s_i} & F \\ \downarrow \rho & & \downarrow \bar{\rho} \\ E_{i+1} & \xrightarrow{\iota} & \bar{E} \end{array}$$

as follows: assuming  $E_0 \rightarrow \cdots \rightarrow E_i$  have been built, Corollary 5.38 implies that a (necessarily unique) compatible degeneracy  $\bar{E} \rightarrow E_i$  exists if and only if it exists at the level of the underlying sets; otherwise, there must exist edges  $e_1, e_2 \in E_i$  and lifts  $f_1, f_2 \in F$  such that  $\rho^{o_i+1}(f_k) = e_k$  and  $\bar{\rho}(f_1) = \bar{\rho}(f_2)$ , and choosing  $s_i$  to be a section of  $\rho^{o_i+1}$  such that  $s_i(e_k) = f_k$ , one defines  $\rho: E_i \rightarrow E_{i+1}$  via the degeneracy-face factorization of  $\bar{\rho}s_i$ , as in the right square in (5.40). The condition on  $s_i$  then guarantees that  $\rho: E_i \twoheadrightarrow E_{i+1}$  is never an isomorphism, so that this procedure always terminates.

It remains to show that  $\Xi[\bar{E}] \amalg_{\Xi[F]} \Xi[E_0] \rightarrow \Xi[E_n]$  is an isomorphism. Surjectivity is immediate from the existence of a section  $s: E_n \rightarrow E_0$  of  $\rho^n$ . For injectivity, the existence of a section  $\bar{s}: \bar{E} \rightarrow F$  of  $\bar{\rho}$  implies that one can represent any element of  $\Xi[\bar{E}] \amalg_{\Xi[F]} \Xi[E_0]$  by an element in  $\Xi[E_0]$ , ie as an equivalence class  $[\varphi]$  for a map  $\varphi: T \rightarrow E_0$ . It now suffices to show by induction on  $i$  that if  $\rho^{o_i}\varphi = \rho^{o_i}\psi$  then  $[\varphi] = [\psi]$  or, equivalently, that  $[\varphi] = [\rho s_i \rho^{o_i}\varphi]$ . The case  $i = 0$  is trivial. Otherwise, we claim that

$$[\varphi] = [\rho s_{i-1} \rho^{o_{i-1}} \varphi] = [\rho s_{i-1} \rho^{o_i} s_i \rho^{o_i} \varphi] = [\rho s_i \rho^{o_i} \varphi].$$

The leftmost and rightmost identities follow from the induction hypothesis for  $i - 1$  and the functions  $\varphi$  and  $\rho s_i \rho^{o_i} \varphi$ . For the inner identity, one considers the lifts  $s_{i-1} \rho^{o_{i-1}} \varphi: T \rightarrow F$  and  $s_{i-1} \rho^{o_i} s_i \rho^{o_i} \varphi: T \rightarrow F$  and notes that postcomposition with  $\bar{\rho}$  yields  $\bar{\rho} s_{i-1} \rho^{o_{i-1}} \varphi = \iota \rho^{o_i} \varphi$  and  $\bar{\rho} s_{i-1} \rho^{o_i} s_i \rho^{o_i} \varphi = \iota \rho^{o_{i+1}} s_i \rho^{o_i} \varphi = \iota \rho^{o_i} \varphi$ . Injectivity now follows, finishing the proof.  $\square$

We now recall the usual [19; 17; 5; 11] description of faces as composites of maximal ‘‘codimension 1’’ faces. We first discuss some terminology.

Firstly, we will regard a face  $F'$  of  $F$  as a subset of  $F$  together with a subset of the broad relations of  $F$ , and write  $F' \hookrightarrow F$  to indicate this. Further, if the broad relations

between edges of  $F'$  in the broad posets  $F'$  and  $F$  coincide, then we will call  $F'$  a *full face* of  $F$  and write  $F' \subset F$  instead.

Secondly, an edge  $e \in F$  is called *external* if  $e$  is either a leaf or a root and *internal* otherwise.

Finally, we denote a generating broad relation of  $F$  by  $v_e = (e^\uparrow \leq e)$  and refer to it as the *vertex at  $e$* .

**Notation 5.41** The maximal faces of  $F$  in  $\Omega$ ,  $\Phi_i$ ,  $\Phi_w$  have the following types:

- The *inner face* (valid for any of  $\Omega$ ,  $\Phi_i$ ,  $\Phi_w$ ) associated to an inner edge  $e$  is the full face  $F - e \subset F$  obtained by removing  $e$ .
- The *leaf vertex outer face* (valid for any of  $\Omega$ ,  $\Phi_i$ ,  $\Phi_w$ ) associated to a vertex  $v_e$  such that  $e^\uparrow$  consists of leaves is the full face  $F_{\neq e} \subset F$  obtained by removing the leaves in  $e^\uparrow$ .
- The *stump outer face* (valid for any of  $\Omega$ ,  $\Phi_i$ ,  $\Phi_w$ ) associated to a stump vertex  $v_e = (\epsilon = e^\uparrow \leq e)$  is the face  $F_{\neq e} \hookrightarrow F$  with the same edges as  $F$  but removing  $\epsilon \leq e$  as a generating broad relation (note that this also removes some composite relations).
- The *root vertex outer face* (valid only for  $\Omega$ ) for an edge  $e \in r^\uparrow$  such that the edges of  $r^\uparrow$  other than  $e$  are leaves is the full face  $T^{\leq e} \subset T$  consisting of those edges  $\bar{e}$  such that  $\bar{e} \leq_d e$ .
- The *root face* (valid only for  $\Phi_i$ ,  $\Phi_w$ ) associated to a root  $r_i \in \underline{r}_F$  which is *not* also a leaf is the full face  $F - r_i \subset F$  obtained by removing  $r_i$ .
- The *stick component face* (valid only for  $\Phi_i$ ) associated to a stick  $\eta \in F$  (ie an edge that is simultaneously a root and a leaf) is the full face  $F - \eta \subset F$  obtained by removing  $\eta$ .

**Remark 5.42** The implicit claim that an inner face  $F - e$  is itself a forest can easily be checked using [Lemma 5.12](#), which shows that  $f^{\uparrow, F-e}$  can be defined to consist of the  $\leq_d$ -maximal edges  $f_i \neq e$  such that  $f_i \leq_d f$ .

Similarly, [Lemma 5.12](#) shows that a broad relation  $f_1 \cdots f_n \leq f$  in  $F$  holds in a stump outer face  $F_{\neq e}$  if the condition  $e \leq_d f$  implies that  $e \leq_d f_i$  for some  $i$ .

**Example 5.43** Consider the trees in [Example 5.29](#). One can write

$$U = (((((T_{\neq c})_{\neq d}) - f) - i) - g),$$

where the intermediate steps (from the inside out) are a stump face, a leaf vertex face



and three inner faces. Further, both  $\eta_a = V^{\leq a}$  and  $\eta_b = V^{\leq b}$  are root vertex outer faces of  $V$  when viewing  $V$  as a tree and  $\eta_a \amalg \eta_b = V - d$  is a root face of  $V$  when viewing  $V$  as a forest.

### 5.3 The category of equivariant trees

Let  $G$  be a finite group. We will denote by  $\Phi^G$  the category of  $G$ -forests, ie forests equipped with a  $G$ -action.

**Definition 5.44** The category of  $G$ -trees, denoted by  $\Omega_G$ , is the full subcategory  $\Omega_G \subset \Phi^G$  of  $G$ -forests  $F$  such that the root tuple  $\underline{r}_F$  consists of a single  $G$ -orbit.

**Remark 5.45** The relationship between  $\Phi^G$  and  $\Omega_G$  is similar to the relationship between the category  $\text{Fin}^G$  of finite  $G$ -sets and the orbital category  $\text{O}_G$  consisting of the orbital  $G$ -sets  $G/H$ .

Examples of equivariant trees can be found throughout Section 4.3. The author is aware that the fact that  $G$ -trees often “look like forests” is likely counterintuitive at first (indeed, that was a major hurdle in the development of the theory presented in this paper). However, the following two facts may assuage such concerns: (i) similarly to how a nonequivariant tree is a forest that can not be decomposed as a coproduct of forests, so too a  $G$ -tree can not be equivariantly decomposed as a coproduct of  $G$ -forests; (ii) the orbital representation of a  $G$ -tree (see Section 4.3) always *does* “look like a tree”.

**Remark 5.46** Note that  $\Omega_G$ , the category of  $G$ -trees, is rather different from  $\Omega^G$ , the category of trees with a  $G$ -action. In fact, each  $G$ -tree is (noncanonically) isomorphic to a forest of the form  $G \cdot_H T$  for some  $H \leq G$  and  $T \in \Omega^H$ . More precisely, one has the following elementary proposition.

**Proposition 5.47**  $\Omega_G$  is equivalent to the Grothendieck construction for the functor

$$\begin{aligned} \text{O}_G^{\text{op}} &\rightarrow \text{Cat}, \\ G/H &\mapsto \Omega^{G/H}, \end{aligned}$$

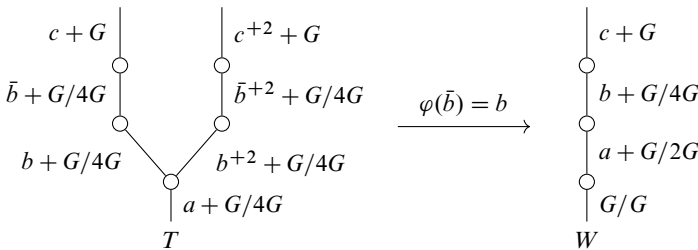
where  $G/H$  denotes the groupoid with objects the cosets  $gH$  and arrows  $gH \xrightarrow{\bar{g}} \bar{g}gH$ . (Here  $\text{O}_G$  is the orbit category; see Remark 5.45.)

**Remark 5.48** There is a natural inclusion  $G \times \Omega \hookrightarrow \Omega_G$  given by regarding each object  $(*, T) \in G \times \Omega$  as the  $G$ -tree given by the  $G$ -free forest  $G \cdot T$ .

**Remark 5.49** While maps in  $\Omega$  can be built out of two types of maps, faces and degeneracies (Proposition 5.37), in  $\Omega_G$  we need a third type of map: quotients.

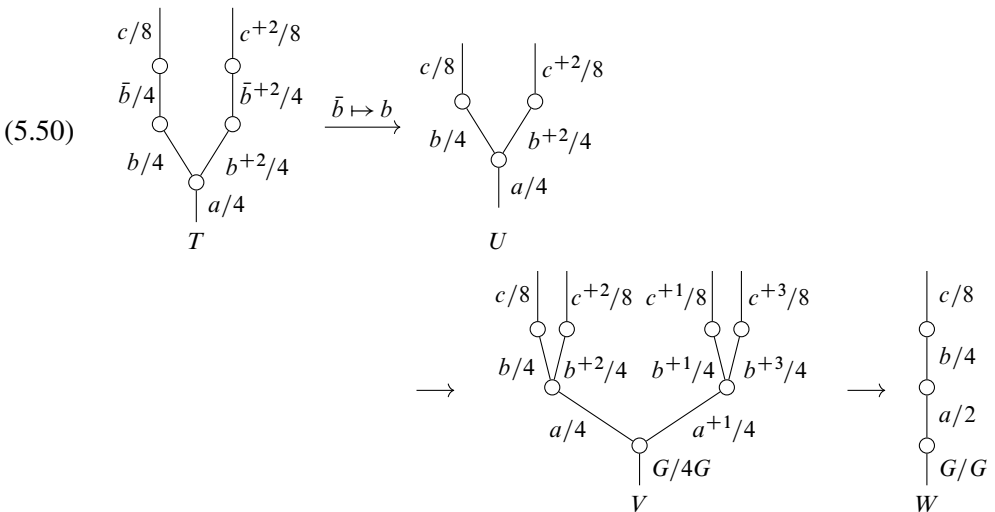
To see this, note that by Proposition 5.47, each  $G$ -tree  $T$  sits (up to equivalence) inside one of the subcategories  $\Omega^{\underline{G/H}}$ , and that  $\Omega^H$  is equivalent to the latter. Since it is immediate by (the proof of) Proposition 5.37 that the degeneracy-face decomposition extends to  $\Omega^H$ , Proposition 5.47 implies that any map in  $\Omega_G$  factors as a degeneracy followed by a face (both inside one of the fibers  $\Omega^{\underline{G/H}}$ ) followed by a cartesian map. We will prefer to refer to cartesian maps as quotients.

For a representative example, let  $G = \mathbb{Z}/8$  and consider the map (represented in orbital notation; see Section 4.3)

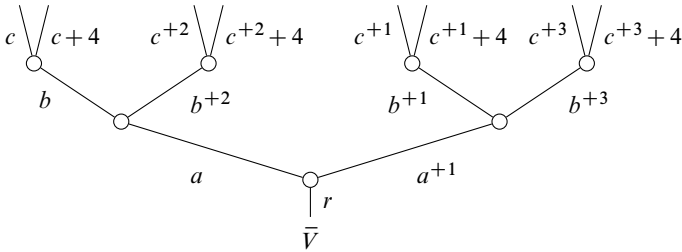


where we follow these conventions: (i) edges in different trees with the same label are mapped to each other; (ii) if an edge is labeled  $e^{+i}$ , we assume that its orbit is disjoint from that of  $e$  and that, if no edge labeled  $e^{+i}$  appears in the target tree, then  $\varphi(e^{+i}) = \varphi(e) + i$ .

This map can be factored as (where, for brevity, we write  $a + G/4G$  as  $a/4$ , etc):



It is perhaps worthwhile to unpack the last map in (5.50), which is an example of a quotient. The  $G$ -tree labeled  $V$  can be written as  $V \simeq G \cdot_{4G} \bar{V}$ , where  $\bar{V} \in \Omega^{4G}$  is the tree with a  $4G$ -action pictured below:



In words,  $V$  consists (nonequivariantly) of four trees identical to  $\bar{V}$  which are interchanged by the action of elements of  $G$  other than 0 and 4. The  $G$ -tree  $W$ , on the other hand, consists of a single (nonequivariant) tree, also shaped like  $\bar{V}$ , and can be thought of as the quotient of  $V$  obtained by gluing the four trees so that the edges  $e^{+i} + j$  and  $e + i + j$  are identified.

**Remark 5.51** One particularly convenient property of  $\Omega$  is that  $\Omega^{\text{op}}$  is a generalized Reedy category, in the sense of [1]. In fact,  $\Omega$  is a *dualizable* generalized Reedy category, so that both  $\Omega$  and  $\Omega^{\text{op}}$  are generalized Reedy.

Unfortunately, this is not the case for  $\Omega_G$ : while indeed  $\Omega_G$  itself can be shown to be generalized Reedy, the opposite category  $\Omega_G^{\text{op}}$  is not. The problem is readily apparent in the factorization in (5.50). Indeed, for the Reedy factorizations to hold (see [1, Definition 1.1(iii)]), quotient maps would need to be considered the same type of maps as face maps, ie degree-raising maps of  $\Omega_G$ . However, the quotient map in (5.50) fails [1, Definition 1.1(iv)'], since there is an automorphism of  $V$  (given by  $e^{+i} \mapsto e^{+i+1} - 1$ , where  $e^{+0}$  is interpreted as  $e$ , and yet undefined  $e^{+i}$  labels are interpreted by regarding  $i \in \mathbb{Z}/k$  as needed) compatible with the quotient map to  $W$ .

### 5.4 Presheaf categories

We now establish some key terminology and notation concerning the presheaf categories we will use. Recall that the category of *dendroidal sets* is the presheaf category  $\text{dSet} = \text{Set}^{\Omega^{\text{op}}}$ .

**Definition 5.52** The category of  $G$ -equivariant dendroidal sets is the category

$$\text{dSet}^G = \text{Set}^{\Omega^{\text{op}} \times G}.$$

The category of *genuine  $G$ -equivariant dendroidal sets* is the category

$$\text{dSet}_G = \text{Set}^{\Omega_G^{\text{op}}}.$$

Twisting the inclusion in [Remark 5.48](#) by the inverse map  $G^{\text{op}} \xrightarrow{(-)^{-1}} G$  yields an inclusion  $u: \Omega^{\text{op}} \times G \hookrightarrow \Omega_G$ .

**Proposition 5.53** *The adjunction*

$$(5.54) \quad u^*: \text{dSet}_G \rightleftarrows \text{dSet}^G : u_*$$

*identifies  $\text{dSet}^G$  as a reflexive subcategory of  $\text{dSet}_G$ .*

**Remark 5.55** Since [Theorem 2.1](#) concerns  $\text{dSet}^G$ , that category will be our main focus throughout the present paper, although  $\text{dSet}_G$  also plays a role in its proof (see [Section 8.2](#)).

Nonetheless,  $\text{dSet}_G$  is arguably the most interesting category. Indeed, the adjunction [\(5.54\)](#) bears many similarities to the adjunction  $\text{sSet}^{\Omega_G^{\text{op}}} \rightleftarrows \text{sSet}^{G^{\text{op}}}$  and, as will be shown in upcoming work, the full structure of the “homotopy operad” of a  $G$ - $\infty$ -operad is described as an object in  $\text{dSet}_G$  rather than in  $\text{dSet}^G$  (more precisely, the claim is that the homotopy operad of a  $G$ - $\infty$ -operad forms a “colored genuine equivariant operad”; (single colored) genuine equivariant operads have been recently formalized by the author and Peter Bonventre in [\[3\]](#)). We note that this is similar to how  $\pi_n$  of a  $G$ -space forms a  $G$ -coefficient system rather than just a  $G$ -set. We conjecture that a model structure on  $\text{dSet}_G$  making [\(5.54\)](#) into a Quillen equivalence exists, and that too is the subject of current work. The presence of extra technical difficulties when dealing with  $\Omega_G$  (see [Remark 5.51](#)), however, make it preferable to address the  $\text{dSet}^G$  case first.

**Notation 5.56** Recall the usual notation

$$(5.57) \quad \Omega \xrightarrow{T \mapsto \Omega[T]} \text{dSet}$$

for the Yoneda embedding.

One can naturally extend this notation to the category  $\Phi$  of forests: given  $F = \coprod_i T_i$ , set  $\Omega[F] = \coprod_i \Omega[T_i]$ . Passing to the  $G$ -equivariant object categories and using the inclusion  $\Omega_G \hookrightarrow \Phi^G$  we will slightly abuse notation and write

$$(5.58) \quad \Omega_G \xrightarrow{T \mapsto \Omega[T]} \text{dSet}^G.$$

More explicitly, if  $T \simeq G \cdot_H T_e$  for some  $T_e \in \Omega^H$ , then  $\Omega[T] \simeq G \cdot_H \Omega[T_e]$ , where  $\Omega[T_e]$  is just the Yoneda embedding of [\(5.57\)](#) together with the resulting  $H$ -action.

**Remark 5.59** While (5.58) defines “representable functors” for each  $T \in \Omega_G$ , given a presheaf  $X \in \text{dSet}^G$  the evaluations  $X(U)$  are defined only for  $U \in \Omega$ , ie for  $U$  a nonequivariant tree.

This is in contrast with  $\text{dSet}_G$ , where both representables and evaluations are defined in terms of  $\Omega_G$ . We note that to reconcile this observation with the inclusion  $u_*$  of (5.54) the nonequivariant tree  $U \in \Omega$  should be reinterpreted as the free  $G$ -tree  $G \cdot U \in \Omega_G$  (see Remark 5.48).

We end this section by discussing a category of “forestial sets” which, while secondary for our purposes, will greatly streamline our discussion of the dendroidal join in Section 7.4.

**Definition 5.60** The category of *wide forestial sets* is the category

$$\text{fSet}_w = \text{Set}^{\Phi_w^{\text{op}}}.$$

**Remark 5.61** The category  $\text{fSet}_i = \text{Set}^{\Phi_i^{\text{op}}}$  of what we might call “independent forestial sets” was one of the main objects of study in [11], where they are called simply “forest sets”.

Mimicking (5.57) by writing  $\Phi_i[F] \in \text{fSet}_i$  for the representable functor of  $F \in \Phi_i$ , it is shown in [11] that one can define a formal boundary  $\partial\Phi_i[F]$  possessing the usual properties one might expect.

We will find it desirable to be able to use the analogous construction for the representable  $\Phi_w[F] \in \text{fSet}_w$ , but this does not quite follow from the result in [11], since while  $\Phi_i[F] \in \text{fSet}_i$  can be forgotten to a presheaf  $u^*\Phi_i[F] \in \text{fSet}_w$ , one typically has a proper inclusion  $\Phi_w[F] \hookrightarrow u^*\Phi_i[F]$ .

We thus instead mimic the discussion in [1], making use of the key technical results established in Section 5.2.

Letting  $\Xi$  denote any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$  and setting

$$|F| = \#\{\text{edges of } F\} + \#\{\text{stumps of } F\},$$

Lemma 5.33 and Proposition 5.37 then say that  $\Xi$  is a dualizable generalized Reedy category (see [1, Definition 1.1]). As in [1, Section 6], call an element  $x: \Xi[F] \rightarrow X$  of a presheaf  $X \in \text{Set}^{\Xi^{\text{op}}}$  *degenerate* if it factors through a noninvertible degeneracy operator and *nondegenerate* otherwise. Corollary 5.39 then allows us to adapt the proof of [1, Proposition 6.9] to obtain the following.

**Proposition 5.62** *Let  $X \in \text{Set}^{\Xi^{\text{op}}}$  for  $\Xi$  any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$ . Then any element  $x: \Xi[F] \rightarrow X$  has a factorization, unique up to unique isomorphism,*

$$\Xi[F] \xrightarrow{\rho_x} \Xi[G] \xrightarrow{\bar{x}} X$$

as a degeneracy operator  $\rho_x$  followed by a nondegenerate element  $\bar{x}$ .

Defining skeleta as in [1, Section 6] the proof of [1, Corollary 6.8] yields the following.

**Corollary 5.63** *Let  $\Xi$  be any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$ . The counit  $\text{sk}_n X \rightarrow X$  for  $X \in \text{Set}^{\Xi^{\text{op}}}$  is a monomorphism whose image consists of those elements of  $X$  that factor through some  $\Xi[F] \rightarrow X$  for  $|F| \leq n$ .*

**Definition 5.64** Let  $\Xi$  be any of  $\Omega$ ,  $\Phi_i$  or  $\Phi_w$ . The formal boundary

$$\partial \Xi[F] \hookrightarrow \Xi[F]$$

is the subobject formed by those maps that factor through a noninvertible map in  $\Xi^+$ , ie through a noninvertible face map.

By combining the Reedy axioms [1, Definition 1.1] with Corollary 5.63 one has

$$\partial \Xi[F] \simeq \text{sk}_{|F|-1} \Xi[F].$$

## 6 Normal monomorphisms and anodyne extensions

### 6.1 Equivariant normal monomorphisms

Recalling that the cofibrations in  $\text{dSet}$  are not the full class of monomorphisms [5, Proposition 1.5], but rather the subclass of so called *normal monomorphisms*, one should expect a similar phenomenon to take place in  $\text{dSet}^G$ .

We start by noting that for  $X \in \text{dSet}^G$  and  $U \in \Omega$ , the set  $X_U = X(U)$  is acted on by the group  $G \times \Sigma_U$ , where  $\Sigma_U$  denotes the automorphism group of  $U$ .

**Definition 6.1** A subgroup  $N \leq G \times \Sigma_U$  is called a *G-graph subgroup* if  $N \cap \Sigma_U = *$ .

It is straightforward to check that a *G-graph subgroup*  $N$  can equivalently be described by a partial homomorphism  $G \geq H \xrightarrow{\rho} \Sigma_U$ . Further, since such a  $\rho$  allows us to write  $U \in \Omega^H$ , one has that each *G-graph subgroup*  $N$  has an associated *G-tree*  $G \cdot_H U$ . More precisely, one has the following result.

**Proposition 6.2** *The functor  $\Omega[-]: \Omega_G \rightarrow \text{dSet}^G$  induces an equivalence between  $\Omega_G$  and the full subcategory of quotients of the form  $(G \cdot \Omega[U])/N$  for  $U \in \Omega$  and  $N \leq G \times \Sigma_U$  a  $G$ -graph subgroup.*

Recalling the discussion following (5.57) one can, for a forest  $F \simeq \coprod_i T_i$  in  $\Phi$ , define  $\partial\Omega[F] = \coprod_i \partial\Omega[T_i]$ . Carrying this discussion through to  $G$ -objects leads to the following definition.

**Definition 6.3** *The boundary inclusions of  $\text{dSet}^G$  are the maps of the form*

$$(6.4) \quad \partial\Omega[T] \hookrightarrow \Omega[T]$$

for  $T \in \Omega_G$ .

More explicitly, if  $T \simeq G \cdot_H T_e$  for some  $T_e \in \Omega^H$ , then the (nonequivariant) presheaf  $\Omega[T_e]$  inherits an  $H$ -action and (6.4) is isomorphic to the map

$$G \cdot_H (\partial\Omega[T_e] \hookrightarrow \Omega[T_e])$$

or, letting  $N \leq G \times \Sigma_{T_e}$  denote the  $G$ -graph subgroup associated to  $T_e$ ,

$$(G \cdot (\partial\Omega[T_e] \hookrightarrow \Omega[T_e]))/N.$$

The following is an immediate generalization of [1, Proposition 7.2].

**Proposition 6.5** *Let  $\phi: X \rightarrow Y$  be a map in  $\text{dSet}^G$ . Then the following are equivalent:*

- (i) *For each tree  $U \in \Omega$ , the relative latching map*

$$l_U(\phi): X_U \coprod_{L_U X} L_U Y \rightarrow Y_U$$

*is a  $\Sigma_U$ -free extension.*

- (ii)  *$\phi$  is a monomorphism and, for each  $U \in \Omega$  and nondegenerate  $y \in Y_U - \phi(X)_U$ , the isotropy group  $\{g \in G \times \Sigma_U \mid gy = y\}$  is a  $G$ -graph subgroup.*
- (iii) *For each  $n \geq 0$ , the relative  $n$ -skeleton*

$$\text{sk}_n(\phi) = X \coprod_{\text{sk}_n X} \text{sk}_n Y$$

*is obtained from the relative  $(n-1)$ -skeleton by attaching boundary inclusions.*

**Definition 6.6** *A monomorphism that satisfies any of the equivalent conditions in Proposition 6.5 will be called a  $G$ -normal monomorphism.*

**Remark 6.7** By Proposition 6.5(i) a monomorphism in  $\text{dSet}^G$  is  $G$ -normal if and only if it is normal as a map in  $\text{dSet}$  [1, Proposition 7.2]. In view of this, we often drop  $G$  from the terminology. Moreover, all monomorphisms over a  $G$ -normal dendroidal set are hence  $G$ -normal monomorphisms [5, Corollaries 1.7 and 1.8].

### 6.2 Equivariant anodyne extensions

The key to the preceding section was the observation that if  $T \in \Omega^H$  then the usual boundary  $\partial\Omega[T]$  inherits an  $H$ -action. However, such is not the case for inner horns: if  $e \in T$  is an inner edge, then  $\Lambda^e[T]$  (see [17, Section 5]) will inherit an  $H$ -action if and only if  $e$  is an  $H$ -fixed edge.

Therefore, to define  $G$ -inner horns, one must treat all inner edges in an *inner edge orbit* in an uniform way. To do so, we first recall the notion of generalized inner horns (see [17, Lemma 5.1]).

**Definition 6.8** Let  $E \subset E^i(T)$  be a subset of the inner edges of  $T \in \Omega$ . We define

$$\Lambda^E[T] \hookrightarrow \partial\Omega[T] \hookrightarrow \Omega[T]$$

to be the subpresheaf formed by the union of those faces *other than* the inner faces of the form  $T - E'$  for  $E' \subset E$ .

More generally, given a forest  $F = \coprod_i T_i$  and  $E = \coprod_i E_i$  with  $E_i \subset E^i(T_i)$  we set

$$(6.9) \quad \Lambda^E[F] = \coprod_i \Lambda^{E_i}[T_i].$$

**Remark 6.10** The reader of [11] may note that (6.9) clashes with [11, Section 3.6]. This is because in [11] the presheaf being defined lives in  $\text{fSet}_i$  rather than in  $\text{dSet}$ .

**Definition 6.11** The *generating  $G$ -inner horn inclusions* are the maps in  $\text{dSet}^G$  of the form

$$\Lambda^{Ge}[T] \rightarrow \Omega[T],$$

where  $T \in \Omega_G$  is a  $G$ -tree and  $Ge$  is the  $G$ -orbit of an inner edge  $e$ .

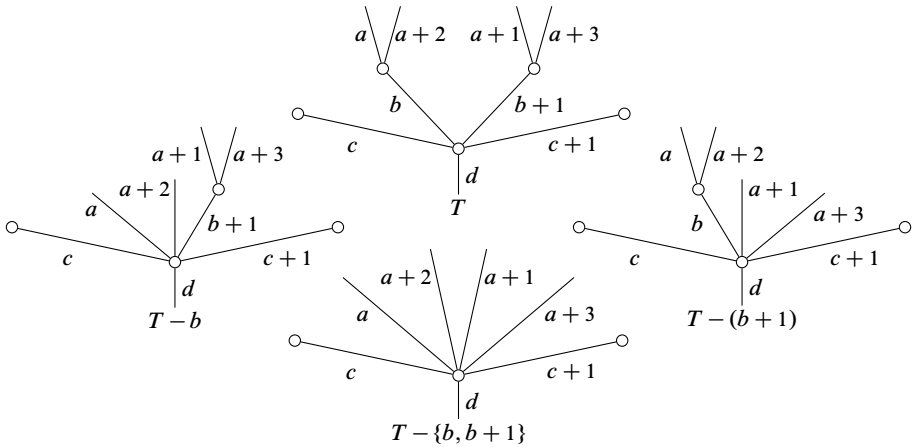
**Definition 6.12** A  $G$ -dendroidal set  $X$  is called a  $G$ - $\infty$ -operad if  $X$  has the right lifting property with respect to all generating  $G$ -inner horn inclusions:

$$\begin{array}{ccc} \Lambda^{Ge}[T] & \longrightarrow & X \\ \downarrow & \nearrow \text{dashed} & \\ \Omega[T] & & \end{array}$$



Further,  $A \rightarrow B$  is called a  $G$ -inner anodyne extension if it is in the saturation of the generating  $G$ -inner horn inclusions under pushouts, transfinite compositions and retracts.

**Example 6.13** If one considers the  $G = \mathbb{Z}/4$ -tree  $T$  in Example 4.5, one possible inner orbit edge is  $Gb = \{b, b + 1\}$ . The following are the (inner) faces of  $T$  not included in  $\Lambda^{Ge}[T]$ :



We recall (see Remark 5.59) that since presheaves  $X \in \text{dSet}^G$  are only evaluated on nonequivariant trees  $U \in \Omega$ , the faces above are merely nonequivariant faces of the equivariant tree  $T$ : indeed,  $T - b$  and  $T - (b + 1)$  do not admit a full compatible  $G$ -action. Rather,  $G$  acts instead on the set of such faces and since  $T - b$  admits a compatible  $H=2G$ -action and one can think of the disjoint union  $(T - b) \amalg (T - (b + 1))$  as the  $G$ -tree  $G \cdot_H (T - b)$ .

**Remark 6.14** An eventual goal of the project this work belongs to is to show that there is a Quillen equivalence

$$W_1: \text{dSet}^G \rightleftarrows \text{sOp}^G : \text{hc}N_d,$$

generalizing [7, Theorem 8.15] (where  $\text{sOp}^G$  is the category of  $G$ -equivariant colored simplicial operads and  $\text{hc}N_d$  is the dendroidal homotopy coherent nerve).

While the proof of such a result is work in progress, and some of the best evidence in that direction is the subject of a current parallel write-up making Remark 5.55 precise, we provide here a first piece of evidence by generalizing [17, Theorem 7.1].

**Proposition 6.15** *Suppose that  $\mathcal{O} \in \text{sOp}^G$  is locally  $G$ -graph fibrant, ie that  $\mathcal{O}(\underline{a}; b)^\Gamma$  is fibrant whenever  $\Gamma \leq G \times \Sigma_{\underline{a}}$  stabilizes  $\underline{a}$  and  $b$  and satisfies  $\Gamma \cap \Sigma_{\underline{a}} = *$ .*

*Then  $\text{hcN}_d(\mathcal{O})$  is a  $G$ - $\infty$ -operad.*

**Proof** Since any  $G$ -tree has the form  $G \cdot_H T$  for some  $T \in \Omega^H$ , it suffices after unpacking adjunctions to solve all the lifting problems as on the left below:

$$(6.16) \quad \begin{array}{ccc} W_1 \Lambda^{He}[T] & \xrightarrow{f} & \mathcal{O} \\ \downarrow & \nearrow \text{dashed} & \uparrow \\ W_1 \Omega[T] & & \end{array} \quad \begin{array}{ccc} W_1 \Lambda^{He}[T](\underline{l}; r) & \longrightarrow & \mathcal{O}(f(\underline{l}); f(r)) \\ \downarrow & \nearrow \text{dashed} & \uparrow \\ W_1 \Omega[T](\underline{l}; r) & & \end{array}$$

Repeating the argument in the proof of [17, Theorem 7.1], it suffices to build this lift for the mapping spaces between the leaves  $\underline{l}$  and root  $r$  of  $T$ . That is, one needs only solve the rightmost lifting problem in (6.16), which needs only be equivariant with respect to the subgroup  $\Gamma \leq H \times \Sigma_{\underline{l}}$  encoding the  $H$ -set  $\underline{l}$  (crucially, note that operadic compatibility is automatic). Since the condition on  $\mathcal{O}$  guarantees that  $\mathcal{O}(f(\underline{l}); f(r))$  is genuinely  $\Gamma$ -fibrant and the vertical map on the right side of (6.16) is (generalizing the formula in the proof of [7, Proposition 4.5])

$$(\{1\} \rightarrow \Delta[1]) \square^{He} \square (\partial \Delta[1] \rightarrow \Delta[1]) \square^{E^i(T)-He},$$

which is a genuine  $\Gamma$ -trivial cofibration, the result follows. □

We will develop for  $G$ -inner horns most of the key results of [17] and [5], the proofs of which turn out to need only moderate modifications in order to generalize to the equivariant context.

The hardest of those results, concerning the tensor product, will be the subject of Section 7. To finish this section, we collect a couple of easier results, starting with the analogue of [17, Lemma 5.1].

**Proposition 6.17** *Let  $T \in \Omega_G$  be a  $G$ -equivariant tree and  $E$  a  $G$ -equivariant subset of the inner edges of  $T$ . Then the generalized  $G$ -horn inclusion*

$$\Lambda^E[T] \hookrightarrow \Omega[T]$$

*is  $G$ -inner anodyne.*

**Proof** Since  $E$  consists of a union of edge orbits, one immediately reduces to proving that maps of the form

$$\Lambda^E[T] \rightarrow \Lambda^{E-Ge}[T]$$

are  $G$ -inner anodyne. In the nonequivariant case [17, Lemma 5.1] such maps can be described as single pushouts, but here we require multiple pushouts, naturally indexed by an equivariant poset which we now describe.

Firstly, let  $T_e$  denote the (nonequivariant) tree component containing the edge  $e$  and set  $H \leq G$  to be its isotropy, resulting in a canonical identification  $G \cdot_H T_e \simeq T$ . Writing  $\text{Inn}_{He}(T_e)$  for the  $H$ -poset (under inclusion) of the inner faces of  $T_e$  collapsing only edges in  $He$ , it suffices to check that, for any  $H$ -equivariant convex<sup>2</sup> subsets  $B \subset B' \subset \text{Inn}_{He}(T_e)$ ,

$$(6.18) \quad \Lambda^E[T] \cup \bigcup_{g \in G, U \in B} \Omega[gU] \rightarrow \Lambda^E[T] \cup \bigcup_{g \in G, U \in B'} \Omega[gU]$$

is  $G$ -inner anodyne. Without loss of generality, we may assume  $B'$  is obtained from  $B$  by adding a single orbit  $HV$  and, setting  $\bar{H} \leq H$  to be the isotropy of  $V$  in  $\text{Inn}_{He}(T_e)$ , we claim that (6.18) is a pushout of

$$G \cdot_{\bar{H}} (\Lambda^{E_e - He}[V] \rightarrow \Omega[V]),$$

where  $E_e = E \cap T_e$  denotes the subset of inner edges of  $T_e$  that are in  $E$ . This claim is straightforward except for the following: one needs to note that the  $G$ -isotropy of any faces in  $\text{Inn}_{E_e - He}(V)$  (ie those faces missing from  $\Lambda^{E_e - He}[V]$ ) is indeed contained in  $\bar{H}$ , and this follows since  $\bar{H}$  can also be described as the subgroup of  $G$  sending the edge subset  $He \cap V = Ge \cap V$  to itself.

This concludes the proof by nested induction on the order of  $G$  and the number of  $G$ -orbits of  $E$ . □

The following is the equivariant analogue of [17, Lemma 5.2]. Note that edge orbits of a  $G$ -tree  $T$  are encoded by maps  $G/H \cdot \eta \xrightarrow{G/H \cdot e} T$  for some  $H \leq G$ .

**Proposition 6.19** *Suppose that  $T$  has a leaf orbit and  $U$  a root orbit both isomorphic to  $G/H$ . Write  $V = T \amalg_{G/H \cdot \eta} U$  for the grafted  $G$ -tree.*

Then

$$(6.20) \quad \Omega[T] \amalg_{\Omega[G/H \cdot \eta]} \Omega[U] \rightarrow \Omega[V]$$

is inner  $G$ -anodyne.

**Proof** Let  $\text{Out}(V)$  denote the  $G$ -poset of outer faces (which have the form  $V_{\neq e}^{\leq e}$ , and

<sup>2</sup>Recall that a subset  $B \subset \mathcal{P}$  of a poset  $\mathcal{P}$  is called *convex* if  $\bar{b} \leq b$  and  $b \in B$  implies  $\bar{b} \in B$ .

are hence not inner faces of any other face) of the grafted tree  $V$ , and  $\text{Out}_{\not\subseteq T,U}(V)$  the  $G$ -subposet of those outer faces contained in neither  $T$  nor  $U$ .

It suffices to show that for all  $G$ -equivariant convex subsets  $B \subset B'$  of  $\text{Out}_{\not\subseteq T,U}(V)$ ,

$$(6.21) \quad \Omega[T] \cup \Omega[U] \cup \bigcup_{R \in B} R \rightarrow \Omega[T] \cup \Omega[U] \cup \bigcup_{R \in B'} R$$

is  $G$ -inner anodyne (indeed, (6.21) recovers (6.20) when  $B = \emptyset$  and  $B' = \text{Out}_{\not\subseteq T,U}(V)$ ).

As before, we can assume  $B'$  is obtained from  $B$  by adding the orbit  $GS$  of a single outer face  $S$ . Letting  $H \leq G$  denote the isotropy of  $S$ , one has that (6.21) is then the pushout (note that the  $G$ -isotropy of an inner face of the outer face  $S$  is at most  $H$ ) of

$$G \cdot_H (\Lambda^{E^i(S)}[S] \rightarrow \Omega[S]),$$

finishing the proof. □

**Remark 6.22** A key difference between the proofs of Propositions 6.17 and 6.19 versus their nonequivariant analogues is the need to check that the isotropies are correct when attaching equivariant horns.

## 7 Tensor products

Our goal in this section is to prove equivariant analogues of [17, Proposition 9.2], [5, Theorem 5.2] and [5, Theorem 4.2], which are the key technical results in their respective papers. These results concern the interaction of anodyne extensions with the tensor product and the join constructions, which we recall in Section 7.1 and Section 7.4. We now present our versions of the results, starting with the analogue of [17, Proposition 9.2].

**Theorem 7.1** *Let  $S, T \in \Omega_G$  be  $G$ -trees and let  $G\xi$  be an inner orbit edge of  $T$ . Then the map*

$$\partial\Omega[S] \otimes \Omega[T] \underset{\partial\Omega[S] \otimes \Lambda^{G\xi}[T]}{\amalg} \Omega[S] \otimes \Lambda^{G\xi}[T] \rightarrow \Omega[S] \otimes \Omega[T]$$

is a  $G$ -inner anodyne extension if either

- (i) both  $S$  and  $T$  are open  $G$ -trees (ie have no stumps);
- (ii) at least one of  $S, T$  is a linear  $G$ -tree (ie isomorphic to  $G \cdot_H [n]$  for  $[n] \in \Delta$ ).

The proof of Theorem 7.1 will be the subject of Section 7.3. More specifically, the result will follow from Proposition 7.44 when  $B = \emptyset$  and  $B' = |\mathbb{E}_{G\xi}(S \otimes T)$ .

The following is the equivariant analogue of [5, Theorem 5.2].

**Theorem 7.2** *Let  $S \in \Omega^G$  be a tree with a  $G$ -action such that  $S \neq \eta$  and denote by  $A \rightarrow \Omega[S] \otimes \Omega[1]$  the pushout product map*

$$\partial\Omega[S] \otimes \Omega[1] \coprod_{\partial\Omega[S] \otimes \{1\}} \Omega[S] \otimes \{1\} \rightarrow \Omega[S] \otimes \Omega[1].$$

*Then there is a factorization  $A \rightarrow B \rightarrow \Omega[S] \otimes \Omega[1]$  such that:*

- (a)  $A \rightarrow B$  is a  $G$ -inner anodyne extension.
- (b) There is a pushout (the join  $S \star \eta$  is introduced in Definition 7.46):

$$(7.3) \quad \begin{array}{ccc} \Lambda^\eta[S \star \eta] & \longrightarrow & B \\ \downarrow & & \downarrow \\ \Omega[S \star \eta] & \longrightarrow & \Omega[S] \otimes \Omega[1] \end{array}$$

- (c) Letting  $\eta \xrightarrow{r} S$  denote the root edge, the composite

$$\Omega[1] \simeq \Omega[\eta \star \eta] \xrightarrow{r \star \text{id}} \Omega[S \star \eta] \rightarrow \Omega[S] \otimes \Omega[1]$$

*coincides with the composite*

$$\Omega[1] \simeq \eta \otimes \Omega[1] \xrightarrow{r \otimes \text{id}} \Omega[S] \otimes \Omega[1].$$

Theorem 7.2 will be proven at the end of Section 7.3 as a direct consequence of the arguments used in the proof of Proposition 7.44. We note that  $\Lambda^\eta[S \star \eta]$  is an outer horn, lacking only the outer face  $(S \star \eta) - \eta = (S \star \eta)^{\leq r} \simeq S$ .

The following is the equivariant analogue of [5, Theorem 4.2]. Note that we use the notation  $i: \Delta \rightarrow \Omega$  for the inclusion and  $i^*: \text{dSet} \rightarrow \text{sSet}$  for the restriction.

**Theorem 7.4** *Let  $S \in \Omega^G$  be a tree with a  $G$ -action. Assume further that  $S$  has at least two vertices and unary root vertex  $G/G \cdot \Omega[1] \xrightarrow{v_r} S$ . Then a lift exists in any commutative diagram*

$$(7.5) \quad \begin{array}{ccc} \Lambda^r[S] & \xrightarrow{f} & X \\ \downarrow & \nearrow & \downarrow \\ \Omega[S] & \longrightarrow & Y \end{array}$$

*such that  $X \rightarrow Y$  is a  $G$ -inner fibration between  $G$ - $\infty$ -operads and  $f(v_r)$  is an equivalence in the  $\infty$ -category  $i^*(X^G)$ .*

We will prove [Theorem 7.4](#) at the end of [Section 7.4](#).

### 7.1 Tensor product

To keep the proofs of [Theorems 7.1](#) and [7.2](#) compact we will prefer to use broad poset language throughout. We start by defining tensor products in this framework.

Given  $\underline{s} \in S^+$  and  $\underline{t} \in T^+$ , we will let  $\underline{s} \times \underline{t} \in (S \times T)^+$  denote the obvious tuple whose elements  $(s, t) \in \underline{s} \times \underline{t}$  are those pairs with  $s \in \underline{s}$  and  $t \in \underline{t}$ .

**Definition 7.6** Given prebroad posets  $S$  and  $T$  (see [Remark 5.2](#)), their *tensor product*  $S \otimes T$  is the prebroad poset whose underlying set is  $S \times T$  and whose relations are generated by relations of the form  $\underline{s} \times t \leq (s, t)$  (resp.  $s \times \underline{t} \leq (s, t)$ ) for  $s \in S, t \in T$  and  $\underline{s} \leq s$  (resp.  $\underline{t} \leq t$ ) a broad relation in  $S$  (resp.  $T$ ).

**Proposition 7.7** *If  $S$  and  $T$  are simple broad posets then so is  $S \otimes T$ . Further, for any nonidentity broad relation  $(s_1, t_1) \cdots (s_n, t_n) \leq (s, t)$  in  $S \otimes T$  one has:*

- (i)  $s_i \leq_d s$  and  $t_j \leq_d t$  for all  $0 \leq i, j \leq n$ .
- (ii) For  $i \neq j$  either there exists  $\underline{s}$  such that  $s_i s_j \underline{s} \leq s$  or there exists  $\underline{t}$  such that  $t_i t_j \underline{t} \leq t$  (or both).
- (iii) If both of the pairs  $s_i, s_j$  and  $t_i, t_j$  are  $\leq_d$ -comparable then  $i = j$ .

**Proof** Note first that the “further” conditions suffice to check that  $S \otimes T$  is a simple broad poset, ie that they imply antisymmetry and simplicity. Indeed, antisymmetry follows from combining (i) with the fact that  $\leq_d^S$  and  $\leq_d^T$  are order relations while simplicity is a particular case of (iii).

Item (i) follows since the condition holds for generating relations and is preserved by transitivity. Similarly, (ii) holds for generating relations and is readily seen to be preserved by transitivity when applied to relations satisfying (i). Lastly, (iii) follows from (ii) and the  $\leq_d$ -incomparability result in [Proposition 5.6](#). □

**Remark 7.8** The main claim in [Proposition 7.7](#) fails for nonsimple broad posets. As an example, let  $S$  be the broad poset  $\{a, b\}$  with generating relations  $ab \leq a$  and  $ab \leq b$  (antisymmetry holds since  $\leq$  decreases the size of tuples) and  $T$  be  $\{c\}$  with generating relation  $\epsilon \leq c$ . Then the relations  $(a, c) \leq (b, c) \leq (a, c)$  hold in  $S \otimes T$ .

**Proposition 7.9** *Let  $S$  and  $T$  be trees. An edge  $(s, t) \in S \otimes T$  has one of five types:*

- (leaf) *It is a leaf if both  $s \in S$  and  $t \in T$  are leaves.*
- (stump) *It is a stump if  $s \in S$  is a leaf and  $t \in T$  is a stump or vice versa, or if both  $s \in S$  and  $t \in T$  are stumps.*
- (leaf node) *It is a node if  $s \in S$  is a node and  $t \in T$  is a leaf or vice versa. In fact  $(s, t)^\uparrow = s^\uparrow \times t$  or  $(s, t)^\uparrow = s \times t^\uparrow$ , accordingly.*
- (null node) *It is a node such that  $\epsilon \leq (s, t)$  if  $s \in S$  is a node and  $t \in T$  a stump or vice versa. In fact  $(s, t)^\uparrow = s^\uparrow \times t$  or  $(s, t)^\uparrow = s \times t^\uparrow$ , accordingly.*
- (fork) *If  $s \in S$  and  $t \in T$  are both nodes then there are exactly two maximal  $\underline{f}$  such that  $\underline{f} < (s, t)$ , namely  $s \times t^\uparrow$  and  $s^\uparrow \times t$ . We call such  $(s, t)$  a **fork**.*

**Proof** Only the fork case requires proof. In fact, it is tautological that only  $s \times t^\uparrow$  and  $s^\uparrow \times t$  can possibly be maximal, hence one needs only verify that neither  $s \times t^\uparrow \leq s^\uparrow \times t$  nor  $s^\uparrow \times t \leq s \times t^\uparrow$ . This follows since the  $S$  coordinates of the pairs in the tuple  $s^\uparrow \times t$  are  $<_d$  than those in  $s \times t^\uparrow$  and vice versa. □

In order to simplify notation, we will henceforth write  $e^{\uparrow S} = s^\uparrow \times t$ ,  $e^{\uparrow T} = s \times t^\uparrow$  and  $e^{\uparrow S, T} = s^\uparrow \times t^\uparrow$ .

**Proposition 7.10** *Let  $S$  and  $T$  be trees and consider the broad relation*

$$\underline{e} = (s_1, t_1)(s_2, t_2) \cdots (s_n, t_n) \leq (s, t) = e$$

in  $S \otimes T$ . Then:

- (i)  $\underline{e} \leq e^{\uparrow S}$  (resp.  $\underline{e} \leq e^{\uparrow T}$ ) if and only if  $s_i \neq s$  for all  $i$  (resp.  $t_i \neq t$  for all  $i$ ).
- (ii)  $\underline{e} \leq e^{\uparrow S, T}$  if and only if both  $s_i \neq s$  and  $t_i \neq t$  for all  $i$ .

**Proof** Only the “if” directions need proof, and the proof follows by upward  $\leq_d$  induction on  $s, t$ . The base cases of either  $s$  or  $t$  a leaf are obvious.

Otherwise, let  $\underline{e}$  satisfy the “if” condition in (i). Since it must be either that  $\underline{e} \leq e^{\uparrow S}$  or  $\underline{e} \leq e^{\uparrow T}$ , we can assume the latter case holds. Writing  $t^\uparrow = u_1 \cdots u_k$  and  $\underline{e} = \underline{e}_1 \cdots \underline{e}_k$  so that  $\underline{e}_i \leq (s, u_i)$  (note that possibly  $k = 0$ ), the induction hypothesis now yields  $\underline{e}_i \leq (s, u_i)^\uparrow S = s^\uparrow \times u_i$ , and hence

$$(7.11) \quad \underline{e} = \underline{e}_1 \cdots \underline{e}_k \leq (s^\uparrow \times u_1) \cdots (s^\uparrow \times u_k) = s^\uparrow \times t^\uparrow \leq s^\uparrow \times t = e^{\uparrow S}.$$

The proof of (ii) simply disregards the last inequality in (7.11). □

**Corollary 7.12**  $\underline{e} \leq e^{\uparrow S, T}$  if and only if both  $\underline{e} \leq e^{\uparrow S}$  and  $\underline{e} \leq e^{\uparrow T}$ .

**Lemma 7.13** Let  $S$  and  $T$  be trees. For any  $e = (s, t) \in S \otimes T$  there exists a minimum  $e^\lambda \in (S \times T)^+$  such that  $e^\lambda \leq e$ . In fact,  $e^\lambda = s^\lambda \times t^\lambda$ .

**Proof** The proof is by  $\leq_d$  induction on  $e$ . The case of  $s, t$  both leaves is obvious. Otherwise, for any nonidentity relation either  $\underline{f} \leq e^{\uparrow S} < e$  or  $\underline{f} \leq e^{\uparrow T} < e$ , and the analysis in the proof of Lemma 5.12 applies in either case to show that indeed  $s^\lambda \times t^\lambda \leq \underline{f}$ . □

## 7.2 Subtrees

**Definition 7.14** Let  $S$  and  $T$  be trees. A *subtree* of  $S \otimes T$  is a tree  $U$  together with a broad poset map  $U \hookrightarrow S \otimes T$  that is an underlying monomorphism. Further, a subtree is called *full* if the relations in  $U$  coincide with those in its image and in that case we instead write  $U \subset S \otimes T$ .

We have the following characterization.

**Proposition 7.15**  $U \xrightarrow{\varphi} S \otimes T$  is full if and only if for each leaf  $l \in U$  it is not the case that  $\epsilon \leq \varphi(l)$ .

**Proof** To simplify notation we will simply write  $u$  for both an edge  $u \in U$  and its image  $\varphi(u) \in S \otimes T$  and instead decorate the broad relations as  $\leq^U, \leq^{S \otimes T}$ , and similarly write  $u^{\lambda, U}, u^{\lambda, S \otimes T}$ , following Lemmas 5.12 and 7.13.

We need to show that a broad relation  $\underline{u} = u_1 \cdots u_n \leq^{S \otimes T} u$  can not fail the conditions in Lemma 5.12 with respect to  $U$ . Further, we note that, since  $S \otimes T$  has a  $\leq_d$ -maximal element  $(r_S, r_T)$ , the proof of Lemma 5.32 applies to show that the two  $\leq_d$ -relations on  $\varphi(U)$  coincide. Thus, only condition (iii) of Lemma 5.12 could possibly fail, and this would happen only if  $u_1^{\lambda, U} \cdots u_n^{\lambda, U}$  lacked some of the leaves in  $u^{\lambda, U}$ . But our hypothesis is that  $l^{\lambda, S \otimes T} \neq \epsilon$  for  $l$  any leaf of  $U$ , hence this is impossible. □

**Definition 7.16** Let  $S$  and  $T$  be trees. A subtree  $U \hookrightarrow S \otimes T$  is called

- *elementary* if all of its generating broad relations are of the form  $e^{\uparrow S} \leq e$  or  $e^{\uparrow T} \leq e$ ;
- *initial* if  $U$  contains the “double root”  $(r_S, r_T) \in S \otimes T$ .

Further, a *maximal elementary subtree* is a subtree that is not contained in any other. Note that, since one can graft new root vertices to  $U$ , Proposition 5.11 implies that maximal elementary subtrees are necessarily initial.



**Remark 7.17** Maximal elementary subtrees are called *percolation schemes* in [17, Example 9.4].

**Remark 7.18** As noted in Proposition 7.9, the relation  $\epsilon \leq e$  is a decomposable relation of  $S \otimes T$  whenever  $e$  is a null node. As a consequence, any elementary tree containing such a generating relation is in fact an inner face of a larger elementary tree.

**Lemma 7.19** *If  $U$  is a full face of  $S$ , then  $U \otimes T \subset S \otimes T$ ; ie  $U \otimes T$  contains all broad relations in its image.*

**Proof** Given a relation  $\underline{e} = (s_1, t_1) \cdots (s_n, t_n) \leq (s, t) = e$  with all  $s_i \in U$  (but not necessarily  $s \in U$ ), we first claim that there is a factorization  $\underline{e} \leq \underline{u} \times t \leq (s, t)$  with  $\underline{u}$  in  $U$  and  $\underline{u} \leq s$  in  $S$ . If  $s \in U$  one simply takes  $\underline{u} = s$ . Otherwise Proposition 7.10 yields  $\underline{e} \leq e^{\uparrow S} \leq e$  and the claim follows by  $\leq_d$  induction on  $e$ .

To check the desired claim that  $\underline{e} \leq e$  will be in  $U \otimes T$  if  $s_i, s \in U$ , we again argue by  $\leq_d$  induction on  $e$ , with the case  $\underline{e} \leq e^{\uparrow T} \leq e$  being immediate and the case  $\underline{e} \leq e^{\uparrow S} \leq e$  following by the result in the previous paragraph.  $\square$

**Definition 7.20** Let  $S, T$  be trees and  $A = \{a_i\}, B = \{b_j\}$  subsets of the sets of stumps of  $S, T$ , respectively, and let  $v_A = \{\epsilon \leq a_i\}, v_B = \{\epsilon \leq b_j\}$  denote the corresponding vertices.

We say that a subtree  $U \hookrightarrow S \otimes T$  *misses  $v_A$  and  $v_B$*  if one has a factorization  $U \hookrightarrow (S - v_A) \otimes (T - v_B) \hookrightarrow S \otimes T$ .

Further, if  $B = \emptyset$  (resp.  $A = \emptyset$ ) we say simply that “ $U$  misses  $v_A$ ” (resp. “ $U$  misses  $v_B$ ”).

**Remark 7.21** In [17] similar notions of “ $U$  missing an inner edge/leaf vertex” are also defined, but we note that (due to Lemma 7.19) those notions are far more straightforward. In fact, as explained in the errata to the follow-up paper [5], the earlier treatment overlooked some subtle properties of stumps.

For instance, note that Corollary 7.24 below implies that the notion “ $U$  misses  $v_A$  and  $v_B$ ” does not coincide with the notion “ $U$  misses  $v_A$  and  $U$  misses  $v_B$ ” whenever  $A$  and  $B$  are both nonempty.

**Lemma 7.22** *Let  $S$  and  $T$  be trees and  $A = \{a_i\}, B = \{b_j\}$  subsets of the stumps of  $S$  and  $T$ , respectively. Then a broad relation*

$$\underline{f} = f_1 f_2 \cdots f_k \leq e$$

in  $S \otimes T$  is a broad relation in  $(S - v_A) \otimes (T - v_B)$  if and only if

$$e^{\lambda, (S-v_A) \otimes (T-v_B)} \leq \underline{f}.$$

**Proof** Only the “if” direction needs proof. We argue by  $\leq_d$  induction on  $e = (s, t)$ . The base case, that of  $s$  a leaf of  $S - v_A$  and  $T$  a leaf of  $T - v_B$ , is obvious (we note that the proof will follow even when this case is vacuous).

Otherwise, either  $\underline{f} \leq e^{\uparrow S} \leq e$  or  $\underline{f} \leq e^{\uparrow T} \leq e$  and our assumption ensures, respectively, that  $s \notin A$  or  $t \notin B$ . Writing  $e^{\uparrow*}$  to denote either  $e^{\uparrow S}$  or  $e^{\uparrow T}$  as appropriate, this last observation guarantees that the relation  $e^{\uparrow*} \leq e$  is in  $(S - v_A) \otimes (T - v_B)$ . Further, writing  $\underline{f} = \underline{f}_1 \cdots \underline{f}_k$  and  $e^{\uparrow*} = e_1 \cdots e_k$  so that  $\underline{f}_i \leq e_i$ , the induction hypothesis shows that these last relations are also in  $(S - v_A) \otimes (T - v_B)$ .  $\square$

Recalling Proposition 5.30 hence yields the following.

**Corollary 7.23** A collection of broad relations of the form  $\underline{g}_i \leq f_i$  and  $f_1 \cdots f_k \leq e$  are all in  $(S - v_A) \otimes (T - v_B)$  if and only if the composite relation  $\underline{g}_1 \cdots \underline{g}_k \leq e$  is.

**Corollary 7.24** A subtree  $U \hookrightarrow S \otimes T$  misses  $v_A$  and  $v_B$  if and only if, for  $r$  the root of  $U$ , one has

$$r^{\lambda, (S-v_A) \otimes (T-v_B)} \leq r^{\lambda, U}.$$

**Proof** This follows from Corollary 7.23 since any generating relation in  $U$  is a factor of  $r^{\lambda, U} \leq r$ .  $\square$

The following was first stated in [5, Proposition 1.9 in errata] and proven via a careful combinatorial analysis in [8]. We include here a short broad poset proof. Recall that  $\text{dSet} \times \text{dSet} \xrightarrow{\otimes} \text{dSet}$  is defined by setting  $\Omega[S] \otimes \Omega[T] = \text{Hom}(-, S \otimes T)$  together with the requirement that  $\otimes$  commutes with colimits in each variable.

**Proposition 7.25** Let  $S, T \in \Omega$  be trees which satisfy either (i) both are open or (ii)  $S = [n]$  is linear. Then the square

$$(7.26) \quad \begin{array}{ccc} \partial\Omega[S] \otimes \partial\Omega[T] & \hookrightarrow & \Omega[S] \otimes \partial\Omega[T] \\ \downarrow & & \downarrow \\ \partial\Omega[S] \otimes \Omega[T] & \hookrightarrow & \Omega[S] \otimes \Omega[T] \end{array}$$

consists of normal monomorphisms. Further,

$$(7.27) \quad \partial\Omega[S] \otimes \Omega[T] \coprod_{\partial\Omega[S] \otimes \partial\Omega[T]} \Omega[S] \otimes \partial\Omega[T] \hookrightarrow \Omega[S] \otimes \Omega[T]$$

is also a normal monomorphism.

In particular, (7.26) is a pullback square.

**Proof** Note first that since  $\otimes$  commutes with colimits in each variable,

$$(7.28) \quad \partial\Omega[S] \otimes \Omega[T] = \operatorname{colim}_{F \in \operatorname{Faces}(S) - \{S\}} \Omega[F] \otimes \Omega[T].$$

In the open case, since all faces are full, Lemma 7.19 implies that if  $U \subset F_S \otimes T$  and  $U \subset S \otimes F_T$ , then  $U \subset F_S \otimes F_T$  will hold, showing that (7.27) is a monomorphism whenever (7.26) consists of monomorphisms. By skeletal induction, it thus suffices to check that the right and bottom maps in (7.26) are monomorphisms, and this will follow if for any  $U \subset S \otimes T$  there exists a minimal face  $F \subset S$  such that  $U \subset F \otimes T$ . Clearly  $F = \{s \mid (s, t) \in U \text{ for some } t\}$  will work once we show that this is indeed a face. This is clear when  $U$  is elementary (in which case each vertex of  $U$  either adds a vertex to  $F$  or nothing at all) and holds in general since any  $U$  is an inner face of an elementary subtree.

In the  $S = [n]$  linear case the fact that  $S$  is open still suffices to conclude that (7.27) being a monomorphism will follow once we show that (7.26) consists of monomorphisms. And, similarly, that  $\partial\Omega[n] \otimes \Omega[T] \rightarrow \Omega[n] \otimes \Omega[T]$  is a monomorphism follows by the same argument, building  $F$  in the same way. It remains to show that  $\Omega[n] \otimes \partial\Omega[T] \rightarrow \Omega[n] \otimes \Omega[T]$  is a monomorphism. We note that the projection  $\pi: [n] \otimes T \rightarrow T$  given by  $\pi(k, e) = e$  is a map of broad posets. Given  $U \hookrightarrow [n] \otimes T$  we claim that  $\pi(U)$  is the minimal face such that  $U \hookrightarrow [n] \otimes \pi(U)$ , noting that this is implied by the more general claim that  $U \hookrightarrow [n] \otimes F$  if and only if  $\pi(U) \hookrightarrow F$ . It now suffices to check this when  $F$  is a maximal face, with the case of  $F$  full being obvious from Lemma 7.19 and the stump outer face case following from Corollary 7.24.

The claim that (7.26) is a pullback square is elementary (compare with the proof of [5, Proposition 1.9]). □

### 7.3 Pushout product filtrations

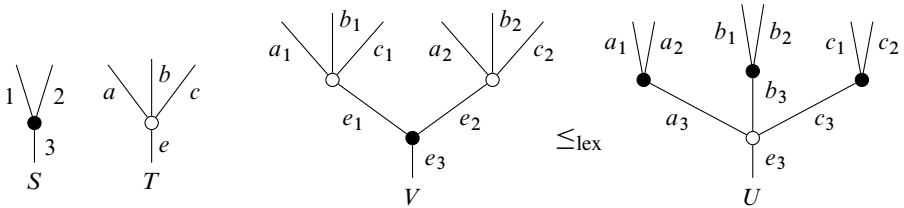
This section features our main technical proofs, namely the proof of Theorem 7.1 and the related but simpler proof of Theorem 7.2.

The majority of the ideas in this section are adapted from the (rather long) proof of [17, Proposition 9.2], but here we will need to significantly repackage those ideas. To explain why, we note that the filtrations in the proof of [17, Proposition 9.2] are actually divided into three nested tiers: an outermost tier described immediately following [17, Corollary 9.3], an intermediate tier described in the proof of [17, Lemma 9.9] and an innermost tier described in the proof of [17, Lemma 9.7]. However, in the equivariant case  $G$  acts transversely to these tiers; ie one can not attach dendrices at an inner tier stage without also attaching dendrices in a different outer tier stage.

Our solution will be to encode the top two filtration tiers as a poset  $\text{IE}_{G\xi}(S \otimes T)$  on which  $G$  acts (to handle the lower tier). To improve readability, however, we first describe our repackaged proof in the nonequivariant case, then indicate the (by then minor) necessary equivariant modifications.

We will make use of an order relation on elementary subtrees (see Definition 7.16) of  $S \otimes T$ .

**Definition 7.29** Write  $V \leq_{\text{lex}} U$  whenever  $U$  is obtained from  $V$  by replacing the intermediate edges in a string of broad relations  $e^{\uparrow S, T} \leq e^{\uparrow S} \leq e$  occurring in  $V$  with the intermediate edges in  $e^{\uparrow S, T} \leq e^{\uparrow T} \leq e$  occurring in  $U$ . An illustrative diagram follows:



**Remark 7.30** The  $\leq_{\text{lex}}$  relation is compatible with the grafting procedure described in Remark 5.16. In particular, we will throughout assume that a relation  $V \leq_{\text{lex}} U$  can be built by first ungrafting  $U$  (Proposition 5.19), then applying  $\leq_{\text{lex}}$  relations to each piece, and finally regrafting the pieces to obtain  $V$ .

In what follows we refer to generating relations of the form  $e^{\uparrow S} \leq e$  (resp.  $e^{\uparrow T} \leq e$ ) in an elementary subtree as  $S$ -vertices (resp.  $T$ -vertices). Also, given vertices  $v = (e^{\uparrow * } \leq e)$  and  $w = (f^{\uparrow * } \leq f)$  we write  $v \leq_d w$  if  $e \leq_d f$ .

**Proposition 7.31** Suppose  $T$  is open (ie has no stumps). Then  $\leq_{\text{lex}}$  induces a partial order on the set of elementary subtrees of  $S \otimes T$ .

Further,  $\leq_{\text{lex}}$  together with the inclusion  $\hookrightarrow$  assemble into a partial order as well, and we denote this latter order simply by  $\leq$ .

**Proof** One needs only check antisymmetry. Let  $g(U)$  count pairs  $(v_S, v_T)$  of an  $S$ -vertex and a  $T$ -vertex in  $U$  such that  $v_S \leq_d v_T$ . Since the generating relations of  $\leq_{\text{lex}}$  strictly increase  $g$ , we have that  $\leq_{\text{lex}}$  is a partial order. Similarly, letting  $h(U) = \#\{\text{stumps of } U\} + \sum_{l \in \{\text{leaves and stumps of } U\}} d(l, r)$  (where  $d(-, r)$  denotes “distance to the root”, measured in generating  $\leq_d$  relations),  $\hookrightarrow$  increases  $h$  and  $\leq_{\text{lex}}$  either (i) preserves  $h$  if  $e$  (as in [Definition 7.29](#)) is a fork; (ii) is an instance of  $\hookrightarrow$  if  $e$  is a null node (since  $T$  is assumed open). Thus  $\leq$  is a partial order.  $\square$

**Remark 7.32** If  $T$  is not open, then by [Remark 7.18](#) it is possible for the combination of  $\leq_{\text{lex}}$  and  $\hookrightarrow$  to fail antisymmetry.

Henceforth we will let  $\xi$  denote a *fixed* inner edge of  $T$ .

**Definition 7.33** An initial elementary subtree  $U \hookrightarrow S \otimes T$  (see [Definition 7.16](#)) is called  $\xi$ -*internal* if it contains an edge of the form  $(s, \xi)$ , abbreviated as  $\xi_s$ , and the  $T$ -vertex  $\xi_s^{\uparrow T} \leq \xi_s$ .

For  $T$  open, we will denote the subposet of such trees by  $(\text{IE}_{\xi}(S \otimes T), \leq)$ .

Further, when  $S$  is open one can modify the order in  $\text{IE}_{\xi}(S \otimes T)$  by reversing the  $\leq_{\text{lex}}$  order (but not the  $\hookrightarrow$  order). The resulting poset will be denoted by  $\text{IE}_{\xi}^{\text{oplex}}(S \otimes T)$ .

**Lemma 7.34** Suppose  $T$  is open. Let  $U \hookrightarrow S \otimes T$  be an elementary subtree with root vertex a  $T$ -vertex  $e^{\uparrow T} \leq e$  and suppose that  $e^{\lambda, U} \leq e^{\uparrow S}$  (or, by [Proposition 7.10](#), that none of the leaves in  $e^{\lambda, U}$  have the same  $S$  coordinate as  $e$ ).

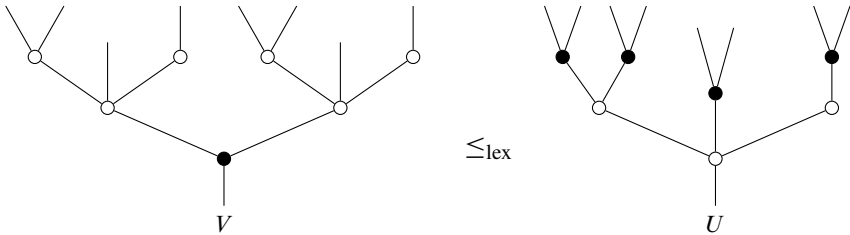
Then there exists an elementary subtree  $V$  such that  $V \leq_{\text{lex}} U$  and  $V$  contains the relations  $e^{\uparrow S, T} \leq e^{\uparrow S} \leq e$ .

**Proof** We argue by induction on the sum of the distances (see proof of [Proposition 7.31](#)) between the leaves and stumps of  $U$  and its root  $e$ . The base case, that of  $U$  the elementary tree generated by  $e^{\uparrow S, T} \leq e^{\uparrow T} \leq e$ , is obvious.

Otherwise, writing  $e = (s, t)$ , we have for each  $t_i \in t^{\uparrow}$  either  $(s, t_i)^{\uparrow S} \leq^U (s, t_i)$  or  $(s, t_i)^{\uparrow T} \leq^U (s, t_i)$ . Applying the induction hypothesis to each of the subtrees  $U^{\leq(s, t_i)}$  (see [Proposition 5.19](#)) in the latter case yields trees  $W_i \leq_{\text{lex}} U^{\leq(s, t_i)}$ , which after

grafted yield a tree  $W \leq_{\text{lex}} U$  such that  $W$  contains all relations  $(s, t_i)^{\uparrow S} \leq (s, t_i)$ . But now  $W$  contains the relations  $e^{\uparrow S, T} \leq e^{\uparrow T} \leq e$  and hence a final generating  $\leq_{\text{lex}}$  relation yields the desired  $V \leq_{\text{lex}} W \leq_{\text{lex}} U$ .  $\square$

**Example 7.35** A typical illustration of the previous result follows:



**Lemma 7.36** Suppose  $U \hookrightarrow S \otimes T$  and  $V \hookrightarrow S \otimes T$  are subtrees with common leaves and root. Then  $F = U \cap V$  defines a full face of both  $U$  and  $V$ .

**Proof** As a set,  $F$  could alternatively be defined as the underlying set of the composite inner face of  $U$  that removes all inner edges of  $U$  not in  $V$ , or vice versa. Thus, the real claim is that both constructions yield the same broad relations. Noting that the  $\leq_d$  order relations on  $U$  and  $V$  are induced from  $S \otimes T$  (as argued in the proof of Proposition 7.15), this follows from Lemma 5.12.  $\square$

**Lemma 7.37** Suppose  $T$  is open. If  $F$  is a common face (resp. inner face) of two elementary subtrees  $U$  and  $V$ , then  $F$  is also a face (resp. inner face) of an elementary subtree  $W$  such that  $W \leq U$  and  $W \leq V$  (resp.  $W \leq_{\text{lex}} U$  and  $W \leq_{\text{lex}} V$ ). In fact, in the inner face case the  $\leq_{\text{lex}}$  inequalities factor through generating  $\leq_{\text{lex}}$  inequalities involving only trees having  $F$  as an inner face.

**Proof** Letting  $r, \underline{l}$  denote the root and leaves of  $F$ , by Corollary 5.23 one can replace  $U$  and  $V$  with  $U \stackrel{\leq r}{\neq \underline{l}}$  and  $V \stackrel{\leq r}{\neq \underline{l}}$ , reducing to the case where  $F, U, V$  have exactly the same leaves and root. Thus, by Lemma 7.36 we are free to assume  $F = U \cap V$ .

If the root vertices  $r^{\uparrow U} \leq r$  and  $r^{\uparrow V} \leq r$  coincide, the result follows by induction on  $\leq$ . Otherwise, we can assume that the root vertex of  $U$  is  $r^{\uparrow S} \leq r$  and that of  $V$  is  $r^{\uparrow T} \leq r$ . Lemma 7.34 now applies to  $V \not\leftarrow r^{\uparrow F}$ , and one can hence build  $W \leq_{\text{lex}} V$  with a strictly larger intersection with  $U$ , finishing the proof.  $\square$

Recall that a subset  $B$  of a poset  $\mathcal{P}$  is called *convex* if  $\bar{b} \leq b$  and  $b \in B$  implies  $\bar{b} \in B$ .

**Proposition 7.38** *Let  $S$  and  $T$  be trees and  $\xi \in T$  an inner edge. Further, assume that either both  $S$  and  $T$  are open or that one of them is linear. Set*

$$A = \Omega[S] \otimes \Lambda^\xi[T] \coprod_{\partial\Omega[S] \otimes \Lambda^\xi[T]} \partial\Omega[S] \otimes \Omega[T]$$

and regard  $A$  and the  $\Omega[V]$  below as subpresheaves of  $\Omega[S] \otimes \Omega[T]$ .

Then, for any convex subsets  $B \subset B'$  of the poset  $\text{IE}_\xi(S \otimes T)$  (or, in the special case of  $S$  a linear tree and  $T$  not open, of the poset  $\text{IE}_\xi^{\text{oplex}}(S \otimes T)$ ), one has that

$$A \cup \bigcup_{V \in B} \Omega[V] \rightarrow A \cup \bigcup_{V \in B'} \Omega[V]$$

is an inner anodyne extension.

To simplify notation we will throughout the proof suppress  $\Omega$  from the notation, eg  $A \cup \bigcup_{V \in B} \Omega[V]$  will be denoted simply by  $A \cup \bigcup_{V \in B} V$ .

The key to the proof is given by the following couple of lemmas.

**Lemma 7.39** *Suppose that either  $S$  and  $T$  are both open or  $T$  is linear.*

Then, for  $U \in \text{IE}_\xi(S \otimes T)$  and  $B$  convex such that  $\{V \mid V < U\} \subset B$ , any edge  $\xi_s$  of  $U$  with vertex  $\xi_s^{\uparrow T} \leq \xi_s$  is (in the sense of [17, Lemma 9.7]) a characteristic edge; ie for each inner face  $F$  of  $U$  containing the edge  $\xi_s$ , we have that  $F$  is in  $A \cup \bigcup_{V \in B} V$  if and only if  $F - \xi_s$  is.

**Proof** Suppose first that  $F - \xi_s$  is in  $A$  but that  $F$  is not.

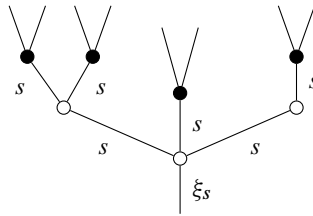
Either  $F - \xi_s \hookrightarrow S' \otimes T$  or  $F - \xi_s \hookrightarrow S \otimes T'$  for  $S' \hookrightarrow S$ ,  $T' \hookrightarrow T$  some maximal subface, where in the latter case  $T' \neq T - \xi$ . Considering the cases in Notation 5.41, the stump cases are excluded by Corollary 7.24 and in the full cases Lemma 7.19 implies that the only possibility is for  $F - \xi_s$  to have no edge with  $S$ -coordinate  $s$  while  $F$  does. Further, if  $s$  were to be a root or leaf of  $S$ , then  $F - \xi_s$  would still contain a root or leaf with  $S$ -coordinate  $s$  (this latter case uses the fact that  $T$  is open). Thus, the only possibility is  $F - \xi_s \hookrightarrow (S - s) \otimes T$  for  $s$  an inner edge of  $S$ .

Now let  $\xi_s^{\uparrow < s} = e_1 \cdots e_k$  consist of the  $\leq_d$ -maximal  $e_i = (s_i, t_i)$  such that both  $e_i <_d \xi_s$  and  $s_i <_d s$ , and consider the subtree  $U \stackrel{\leq \xi_s}{\not\leftarrow \xi_s^{\uparrow < s}}$ . Then (i) this tree has no leaf with  $S$  coordinate  $s$ , or else that would be a leaf of  $U$  (Remark 5.22), and thus also of  $F$ , so that  $F - \xi_s \hookrightarrow (S - s) \otimes T$  could not hold; (ii) the leaf tuple of this tree is

hence  $\xi_s^{\uparrow < s}$ ; (iii) by definition of  $\xi_s^{\uparrow < s}$ , all inner edges of this tree have  $S$  coordinate  $s$ . But the condition  $F - \xi_s \hookrightarrow (S - s) \otimes T$  now implies that  $F$  contains none of the inner edges of  $U \stackrel{\leq \xi_s}{\not\leftarrow \xi_s^{\uparrow < s}}$ , so that Lemma 7.34 implies that  $F$  is a subface of some  $V <_{\text{lex}} U$ , hence contained in  $A \cup \bigcup_{V \in B} V$ .

Suppose now that  $F - \xi_s$  is a subface of some  $V \in B$ . By Lemma 7.37 (and its proof) we can assume that in fact  $F - \xi_s$  is an inner subface of  $V$  and  $V <_{\text{lex}} U$ . Further, by the “in fact” part of Lemma 7.37 one can also assume that this is a generating  $\leq_{\text{lex}}$  relation. But then  $V$  necessarily contains  $\xi_s$ , since generating  $\leq_{\text{lex}}$  relations do not add edges whose vertex is a  $T$ -vertex. Thus  $F - \xi_s \hookrightarrow V$  implies  $F \hookrightarrow V$ .  $\square$

**Example 7.40** The following is a typical tree illustration of  $U \stackrel{\leq \xi_s}{\not\leftarrow \xi_s^{\uparrow < s}}$ :



In words, this subtree always has (i)  $S$ -vertices (in black), all of the same arity, below its leaves; (ii) a  $T$ -vertex above its root  $\xi_s$  (in white); (iii) all its remaining vertices  $T$ -vertices (so that the edges marked  $s$  have  $S$ -coordinate  $s$ ).

**Remark 7.41** For the previous proof to work it is crucial for the tree  $U \stackrel{\leq \xi_s}{\not\leftarrow \xi_s^{\uparrow < s}}$  to in fact have inner edges, as is ensured by the fact that  $\xi$  is not a stump of  $T$ . In this latter case we will instead need to use the following alternative lemma.

**Lemma 7.42** Suppose that  $S$  is a linear tree (ie  $S \simeq [n]$  for  $[n] \in \Delta$ ).

Then for  $U \in \text{IE}_\xi^{\text{oplex}}(S \otimes T)$  the  $\leq_d$ -maximal edge of  $U$  of the form  $\xi_s$  is a characteristic edge (in the sense of Lemma 7.39).

**Proof** Since  $S$  is linear we will for simplicity label its edges by  $0 \leq 1 \leq \dots \leq n$ .

Suppose first that  $F - \xi_s$  is in  $A$ , so that repeating the argument in the proof of Lemma 7.39 we conclude it must be  $F - \xi_s \hookrightarrow (S - s) \otimes T$  for  $s < n$  (note that this makes sense even if  $s = 0$  is the leaf of  $S$ , which must be considered if  $T$  is not open).

Since  $s \neq n$ , one can choose a  $\leq_d$ -minimal edge  $a_{s+1}$  of  $U$  such that  $\xi_s \leq_d a_{s+1}$ . Then (i) the characterization of  $\xi_s$  implies that  $a \neq \xi$ ; (ii) the characterization of  $a_{s+1}$



implies that  $U$  contains the  $S$ -vertex  $a_s = a_{s+1}^{\uparrow S} \leq a_{s+1}$ ; (iii) since  $U$  contains both  $\xi_s$  and  $a_s$ , it contains the  $T$ -vertex  $a_s^{\uparrow T} \leq a_s$ , which can be rewritten as  $a_{s+1}^{\uparrow S, T} \leq a_{s+1}^S$ . Therefore,  $U$  contains the relations  $a_{s+1}^{\uparrow S, T} \leq a_{s+1}^{\uparrow S} \leq a_{s+1}$ ; since  $F$  must collapse  $a_s = a_{s+1}^{\uparrow S}$ , this yields that  $F$  is a subspace of the tree  $V \leq_{\text{oplex}} U$  that is obtained by replacing  $a_s = a_{s+1}^{\uparrow S}$  with  $a_{s+1}^{\uparrow T}$ .

The case of  $F - \xi_s$  a subspace of some  $V \in B$  follows by an argument identical to that in the proof of Lemma 7.39, except now noting that generating  $\leq_{\text{oplex}}$  relations do not add edges whose vertex is a  $S$ -vertex.  $\square$

**Proof of Proposition 7.38** Without loss of generality we can assume that  $B'$  is obtained from  $B$  by adding a single  $\xi$ -internal initial elementary tree  $U$  with  $\xi_s$  its corresponding edge.

We first note that the outer faces of  $U$  are in  $A \cup \bigcup_{V \in B} V$ . Since a maximal outer face  $\bar{U} \hookrightarrow U$  is always still elementary,  $\bar{U}$  will be  $\xi$ -internal initial elementary unless (i)  $\bar{U} = U^{\leq(e, rT)}$  (resp.  $\bar{U} = U^{\leq(rs, e)}$ ) is a root vertex face, in which case  $\bar{U} \hookrightarrow (S^{\leq e}) \otimes T$  (resp.  $\bar{U} \hookrightarrow S \otimes (T^{\leq e})$ ) and is hence in  $A$ ; (ii)  $\bar{U} = U_{\neq \xi_s}$  and  $\bar{U}$  is no longer  $\xi$ -internal since it contains no  $T$ -vertices of the form  $\xi_s^{\uparrow T} \leq \xi_s$ . But then  $\bar{U} \hookrightarrow S \otimes T_{\neq \xi}$  would hold (by either Lemma 7.19 or Corollary 7.24) and thus  $\bar{U}$  is in  $A$ .

Finally, we let  $\text{Inn}_{\xi_s}(U)$  denote the poset of inner faces of  $U$  removing only edges other than  $\xi_s$ . We claim that for any convex subsets  $C \subset C' \subset \text{Inn}_{\xi_s}(U)$  the map

$$(7.43) \quad A \cup \bigcup_{V \in B} V \cup \bigcup_{W \in C} W \rightarrow A \cup \bigcup_{V \in B} V \cup \bigcup_{W \in C'} W$$

is inner anodyne. We argue by induction on  $C$  and again we can assume that  $C'$  is obtained from  $C$  by adding a single  $X \in \text{Inn}_{\xi_s}(U)$  not yet in the domain of (7.43). The concavity of  $C$  and  $C'$  and the characteristic edge condition in Lemmas 7.39 and 7.42 then imply that the only faces of  $X$  not in the source of (7.43) are precisely  $X$  and  $X - \xi_s$ , showing that (7.43) is a pushout of  $\Lambda^{\xi_s}[X] \rightarrow \Omega[X]$ , finishing the proof.  $\square$

In the  $G$ -equivariant case, given an inner edge orbit  $G\xi$ , we write  $\text{IE}_{G\xi}(S \otimes T)$  for the poset of initial elementary trees containing at least one  $T$ -vertex of the form  $(g\xi)_s^{\uparrow T} \leq (g\xi)_s$  (alternatively, one has  $\text{IE}_{G\xi}(S \otimes T) = \bigcup_{g \in G} \text{IE}_{g\xi}(S \otimes T)$ ). Note that in this case the group  $G$  acts on the poset  $\text{IE}_{G\xi}(S \otimes T)$  as well. The following is the equivariant version of Proposition 7.38.

**Proposition 7.44** *Let  $S, T \in \Omega_G$  be  $G$ -trees and  $\xi \in T$  an inner edge. Further, assume that either both  $S$  and  $T$  are open or that one of them is linear (ie of the form  $G/H \cdot [n]$ ). Set*

$$A = \Omega[S] \otimes \Lambda^{G\xi}[T] \coprod_{\partial\Omega[S] \otimes \Lambda^{G\xi}[T]} \partial\Omega[S] \otimes \Omega[T],$$

and regard  $A$  and the  $\Omega[V]$  below as subpresheaves of  $\Omega[S] \otimes \Omega[T]$ .

Then for any  $G$ -equivariant convex subsets  $B \subset B'$  of  $\text{IE}_{G\xi}(S \otimes T)$  (or, in the special case of  $S$  a linear tree and  $T$  not open, of  $\text{IE}_{G\xi}^{\text{oplex}}(S \otimes T)$ ) one has that

$$A \cup \bigcup_{V \in B} \Omega[V] \rightarrow A \cup \bigcup_{V \in B'} \Omega[V]$$

is a  $G$ -inner anodyne extension.

Again we will suppress  $\Omega$  from the notation of the proof.

**Proof** Note we are free to assume that  $S, T \in \Omega^G \subset \Omega_G$ , ie that  $S$  and  $T$  are actual trees with a  $G$ -action rather than  $G$ -indecomposable forests. Indeed, otherwise writing  $S \simeq G \cdot_H S_e$  and  $T \simeq G \cdot_K T_e$  for  $S_e \in \Omega^H$  and  $T_e \in \Omega^K$  yields a decomposition

$$S \otimes T \simeq \coprod_{[g] \in H \backslash G / K} G \cdot_{H \cap gKg^{-1}} S_e \otimes gT_e,$$

where when regarding  $S_e, gT_e \in \Omega^{H \cap gKg^{-1}}$  we omit the forgetful functors.

In analogy with the nonequivariant case we can assume  $B'$  is obtained from  $B$  by adding the  $G$ -orbit of a single  $\xi$ -internal initial elementary tree  $U$  with  $\xi_s$  the corresponding edge. Let  $H \leq G$  denote the  $G$ -isotropy of  $U$  in  $\text{IE}_{G\xi}(S \otimes T)$ .

That the outer faces of any of the conjugates  $gU$  are in  $A \cup \bigcup_{V \in B} V$  follows by the corresponding nonequivariant argument in the proof of [Proposition 7.38](#).

The key is now to prove the equivariant analogues of [Lemmas 7.39](#) and [7.42](#), stating that  $H\xi_s$  is a *characteristic edge orbit* of  $U$ , ie that for each inner face  $F$  of  $U$  containing all edges in  $H\xi_s$ , we have that  $F$  is in  $A \cup \bigcup_{V \in B} V$  if and only if  $F - H\xi_s$  is.

When proving the equivariant analogue of [Lemma 7.39](#), in the case of  $F - H\xi_s$  in  $A$ , the argument in that proof yields that  $F$  itself must already lack all the inner edges of at least one of the  $U_{\substack{\leq h\xi_{hs} \\ \neq h\xi_{hs}^{\uparrow <hs}}}$  subtrees (where  $h \in H$ ) and therefore applying [Lemma 7.34](#)

again shows that  $F$  is a subface of some  $V <_{\text{lex}} U$ . In the case of  $F = H\xi_s$  in some  $V < U$ , repeating the argument in the proof we can again assume  $V <_{\text{lex}} U$  via a generating  $\leq_{\text{lex}}$  relation and such  $V$  must likewise contain all the edges in  $H\xi_s$ .

Proving the equivariant analogue of [Lemma 7.42](#) requires no changes to the proof, since defining  $a_{s+1}$  in the same way one still concludes that  $F$  must lack  $a_s$  (in fact,  $F$  must lack  $Ha_s$ , though that fact is not needed).

Lastly, we equivariantly modify the last two paragraphs in the proof of [Proposition 7.38](#): setting  $\text{Inn}_{\widehat{H\xi_s}}(U)$  to be the  $H$ -poset of inner faces of  $U$  lacking only edges *not* in  $H\xi_s$  (ie containing all edges in  $H\xi_s$ ), we show by induction on  $H$ -equivariant concave subsets  $C \subset C' \subset \text{Inn}_{\widehat{H\xi_s}}(U)$  that the map

$$(7.45) \quad A \cup \bigcup_{V \in B} V \cup \left( \bigcup_{g \in G, W \in C} gW \right) \rightarrow A \cup \bigcup_{V \in B} V \cup \left( \bigcup_{g \in G, W \in C'} gW \right)$$

is inner  $G$ -anodyne. Again we can assume that  $C'$  is obtained from  $C$  by adding the  $H$ -orbit of a single  $X$  containing all edges in  $H\xi_s$  and not in the domain of (7.45). Write  $\bar{H} \leq H$  for the  $H$ -isotropy of  $X$ . At this point an extra equivariant argument is needed: one needs to know that for all subsets  $E \subset H\xi_s$  the  $G$ -isotropy of  $X - E$  coincides with its  $H$ -isotropy (which is necessarily a subgroup of  $\bar{H}$ ). To see this, note that if it were otherwise then  $X - E$  would be contained in both  $U$  and a distinct conjugate  $gU$ , and [Lemma 7.37](#) would imply that  $X$  is already in the domain of (7.45). Finally, repeating the “characteristic edge (orbit)” argument we see that the map (7.45) is a pushout of (note that this is only a generating horn if  $\bar{H} = H$ , so that [Proposition 6.17](#) is needed otherwise)

$$G \cdot \bar{H} (\Lambda^{H\xi_s}[X] \rightarrow \Omega[X]),$$

finishing the proof. □

We now adapt the previous proof to deduce the easier [Theorem 7.2](#).

**Proof of Theorem 7.2** The argument follows by attempting to follow the proof of [Proposition 7.44](#) when  $T = [1]$  and  $\xi = 1$  (note that [Definition 7.33](#) still makes sense, although the edge  $1_s$  of  $U$  may now possibly be the root, hence not internal). The only case where (the equivariant analogue) of [Lemma 7.39](#) does not provide a characteristic edge orbit is when the root vertex of  $U$  is  $1_r \uparrow^T \leq 1_r$ , in which case  $U$  will be in  $A$  unless (using the notation in the proof of [Proposition 7.25](#))  $\pi(U) = S$ , so that in fact

$U = S \otimes \{0\} \cup \{1_r\}$ . Denoting this latest  $U$  as  $S \star \eta$  and noting that it is the maximum of the poset  $IE_{\xi}(S \otimes [1])$  one concludes that letting  $A \rightarrow B$  denote

$$A \rightarrow A \cup \bigcup_{V \in IE_{\xi}(S \otimes [1]), V \neq S \star \eta} V,$$

this is indeed inner  $G$ -anodyne. The pushout (7.3) follows by noting that the only face of  $S \star \eta$  not in  $B$  is  $(S \star \eta) - \eta$  (in fact, the only other face  $F$  such that  $\pi(F) = S$  is  $(S \star \eta) - 0_r$ , which is a common face of the elementary tree  $V$  obtained by applying a  $\leq_{\text{lex}}$  relation to the root of  $S \star \eta$ ).  $\square$

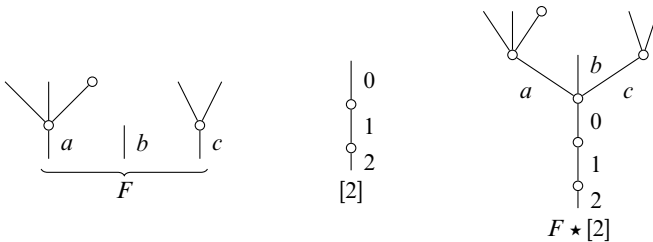
### 7.4 Dendroidal join

We now turn to the equivariant version of the dendroidal join  $\star$  discussed in [5, Section 4], which will be needed to understand the last piece of the filtration in Theorem 7.2. We recall that several categories of forests were discussed in Section 5.2.

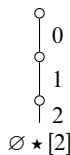
**Definition 7.46** Given an object  $F \in \Phi$  and  $[n] \in \Delta$  we define  $F \star [n] \in \Omega$  as the broad poset having underlying set  $F \amalg [n]$  and relations

- $e_1 \cdots e_n \leq e$  if  $e_i, e \in F$  and  $e_1 \cdots e_n \leq^F e$ ;
- $i \leq j$  if  $i, j \in [n]$  and  $i \leq^{[n]} j$ ;
- $e_1 \cdots e_n \leq i$  if  $e_j \in F, i \in [n]$  and  $e_1 \cdots e_n \leq^F \ulcorner_F$ .

**Example 7.47** As explained in [5, Section 4.3], one can readily visualize  $\star$  when using tree diagrams, such as in the following example:



Further, note that when  $F = \emptyset$  is the empty forest, we have that  $\ulcorner_F = \epsilon$ , and since  $\epsilon \leq \epsilon$ , we also have that  $\emptyset \star [n]$  adds a stump at the top of  $[n]$ :



We now discuss the functoriality of  $\star$ . As implicit in the discussion in [5, Section 4.5],  $\star$  is only functorial with respect to some maps of forests. Indeed, it is clear from the third condition in Definition 7.46 that  $\star$  will be functorial in  $F$  precisely with respect to the maps in  $\Phi_w$ .

Moreover, the canonical inclusions  $[n] \rightarrow F \star [n]$  and  $F \rightarrow F \star [n]$  can be thus encoded: letting  $\Delta_+$ ,  $\Phi_{w+}$  and  $\Phi_{i+}$  denote the categories  $\Delta$ ,  $\Phi_w$  and  $\Phi_i$  together with an additional initial object  $+$ , one has the following.

**Proposition 7.48** *The join  $\star$  defines a bifunctor*

$$\Phi_{w+} \times \Delta_+ \xrightarrow{-\star-} \Phi_{i+}$$

such that  $+ \star [n] = [n]$  and  $F \star + = F$ .

Note that  $\star$  usually lands in  $\Omega_+$ , the only exceptions occurring when the second input is the additional initial object.

We now extend the join operation to presheaves by defining

$$\text{fSet}_w \times \text{sSet} \xrightarrow{-\star-} \text{dSet}$$

to be the composite (writing  $u: \Xi \hookrightarrow \Xi_+$  for the inclusion and  $u^*: \text{Set}^{\Xi_+ \text{op}} \rightleftarrows \text{Set}^{\Xi \text{op}} : u_*$  for the standard adjunction)

$$(7.49) \quad \text{fSet}_w \times \text{sSet} \xrightarrow{u_* \times u^*} \text{Set}^{\Phi_{w+} \text{op}} \times \text{Set}^{\Delta_+ \text{op}} \xrightarrow{\times} \text{Set}^{\Phi_{w+} \times \Delta_+ \text{op}} \xrightarrow{\text{Lan}\star} \text{Set}^{\Phi_{i+} \text{op}} \xrightarrow{u^*} \text{dSet}.$$

**Remark 7.50** Unpacking (7.49) one can write (see [13, Definition 1.2.8.1])

$$(7.51) \quad (X \star Y)(T) = X(T) \amalg Y(T) \amalg \coprod_{F \rightarrow T, [n] \rightarrow T, T \simeq F \star [n]} X(F) \times Y([n]),$$

where  $Y(T) = \emptyset$  when  $T$  is not linear.

**Remark 7.52** Because of the passage through the  $(-)_+$  categories in (7.49), the join  $\text{fSet}_w \times \text{sSet} \xrightarrow{-\star-} \text{dSet}$  does not preserve colimits in each variable. Rather, the functors  $F \star (-)$  and  $(-) \star [n]$  preserve colimits when mapping into the under categories  $\text{dSet}_{F/}$  and  $\text{dSet}_{[n]/}$ . Therefore,  $\star$  does nonetheless preserve *connected colimits* in each variable.

The following is an equivariant generalization of a key technical lemma [5, Lemma 4.10] combined with key arguments in the proof of [5, Theorem 4.2].

**Proposition 7.53** Let  $A \xrightarrow{f} B$  be a normal monomorphism in  $\mathbf{fSet}_w^G$  (defined as in Proposition 6.5(iii)) and  $C \xrightarrow{g} D$  be a left anodyne map in  $\mathbf{sSet}$ . Then

$$f \square^* g: A \star D \coprod_{A \star C} B \star C \rightarrow C \star D$$

is  $G$ -inner anodyne.

**Proof** That  $f \square^* g$  is indeed a monomorphism whenever  $f$  and  $g$  are monomorphisms follows directly from (7.51).

In light of Remark 7.52 concerning connected colimits, it suffices to consider the case where  $f$  and  $g$  have the forms  $\partial\Phi_w[F] \rightarrow \Phi_w[F]$  and  $\Lambda^i[n] \rightarrow \Delta[n]$  for  $0 \leq i < n$ . But in that case  $f \square^* g$  is simply the inner horn inclusion

$$\Lambda^i(F \star [n]) \rightarrow \Omega(F \star [n]). \quad \square$$

**Proof of Theorem 7.4** First note that the conditions on  $S$  are equivalent to saying that  $S \simeq F \star [1]$  for some  $F \in \Phi_w^G$ . One can thus rewrite the left vertical map in Theorem 7.4 as

$$(\Lambda^1[1] \rightarrow \Delta[1]) \square^* (\emptyset \rightarrow \Phi_w[F]),$$

and denoting by  $A \setminus (-)$  the right adjoint to  $\mathbf{sSet} \xrightarrow{A \star (-)} \mathbf{dSet}_{A \star +}^G$ , standard adjunction arguments allows us to convert (7.5) into the equivalent lifting problem

$$(7.54) \quad \begin{array}{ccc} \Lambda^1[1] & \longrightarrow & \Phi_w[F] \setminus X \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta[1] & \longrightarrow & \Phi_w[F] \setminus Y \times_{\emptyset \setminus Y} \emptyset \setminus X \end{array} \quad \begin{array}{ccc} \{0\} & \longrightarrow & \Phi_w[F] \setminus X \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta[1] & \longrightarrow & \Phi_w[F] \setminus Y \times_{i^*(Y^G)} i^*(X^G) \end{array}$$

where the right-hand diagram merely simplifies the notation on the left:  $\emptyset \setminus Z \simeq i^*(Z^G)$ , where we caution that  $\emptyset \neq \Phi_w[\emptyset]$ , the former being the *empty presheaf* and the latter the representable presheaf on the *empty forest*.

Standard repeated applications of Proposition 7.53 (setting  $A = \emptyset$  or  $Y = *$  as needed) yield that (i)  $\Phi_w[F] \setminus X \rightarrow i^*(X^G)$  and  $\Phi_w[F] \setminus Y \rightarrow i^*(Y^G)$  are left fibrations; (ii)  $\Phi_w[F] \setminus X$  and  $\Phi_w[F] \setminus Y$  are left fibrant and thus  $\infty$ -categories; (iii) the rightmost map  $\Phi_w[F] \setminus X \rightarrow \Phi_w[F] \setminus Y \times_{i^*(Y^G)} i^*(X^G)$  in (7.54) is a left fibration. Therefore, the map  $\Phi_w[F] \setminus Y \times_{i^*(Y^G)} i^*(X^G) \rightarrow i^*(X^G)$  is itself a left fibration (it is a pull-back of  $\Phi_w[F] \setminus Y \rightarrow i^*(Y^G)$ ) and since left fibrations conserve equivalences [12, Proposition 2.7] the image of the lower map in (7.54) is a equivalence. The result now

follows since equivalences can be lifted over left fibrations between  $\infty$ -categories [12, Propositions 2.4 and 2.7].  $\square$

**Remark 7.55** The interested reader may note that Proposition 7.53 and Theorem 7.4 have notably shorter proofs than the analogous results [5, Lemma 4.10 and Theorem 4.2]. In fact, some of our arguments closely resemble the proof for the simplicial case as found in [13, Proposition 1.2.4.3]. The brevity of these arguments is the reason for our introduction of the forest category  $\Phi_w$ : from our perspective, many of the arguments in [5] are replaced with the task of showing that the boundaries  $\partial\Phi_w[F]$  satisfy the usual formal properties.

## 8 Model structure on equivariant dendroidal sets

We now adapt the treatment in [5] to equip the category  $dSet^G$  with a model structure where the cofibrations are the normal monomorphisms (see Proposition 6.5) and the fibrant objects are the  $G$ - $\infty$ -operads (see Definition 6.12). Since we have already established the equivariant analogues of the main technical results needed (Theorems 7.1, 7.2 and 7.4), the proofs in [5] will now carry over to our context with only minor changes needed.

### 8.1 Existence of the model structure

Our goal in this first section is to establish Proposition 8.8, which abstractly builds the desired model structure on  $dSet^G$  by generalizing [5, Proposition 3.12].

As it turns out, adapting the treatment in [5, Section 3] requires only minimal modifications (essentially “adding  $G$  to the statements therein”). As such, we will be brief in our discussion, and refer the reader to [5, Section 3] for extra details.

Recall the notation  $J = N(0 \rightrightarrows 1)$  for the nerve of the groupoid generated by a single isomorphism between two distinct objects. As in [5, Section 3.2] we write  $J_d = i_!(J)$  when regarding  $J$  as a dendroidal set, and we will further regard  $J_d$  as a  $G$ -dendroidal set by equipping it with the *trivial*  $G$ -action.

Following [5, Section 3.2], we define  $An$ , the class of  $J$ -anodyne extensions, to be the saturation of the  $G$ -inner horn inclusions together with the maps

$$(8.1) \quad \{i\} \otimes \Omega[T] \coprod_{\{i\} \otimes \partial\Omega[T]} J_d \otimes \partial\Omega[T] \rightarrow J_d \otimes \Omega[T], \quad i = 0, 1, \quad T \in \Omega_G.$$

A  $G$ -dendroidal set  $X$  (resp. map  $X \rightarrow Y$ ) is then called  $J$ -fibrant (resp.  $J$ -fibration) if it has the right lifting property with respect to the maps in  $An$ .

The following generalizes the necessary parts of [5, Proposition 3.3] (note that that result is slightly corrected in the errata to [5]).

**Proposition 8.2** *Let  $A \rightarrow B$  be a normal monomorphism in  $dSet^G$ . Then:*

- (i) *For  $i = 0, 1$ , the map  $\{i\} \otimes B \amalg_{\{i\} \otimes A} J_d \otimes A \rightarrow J_d \otimes B$  is in  $An$ .*
- (ii)  *$\{0, 1\} \otimes B \amalg_{\{0,1\} \otimes A} J_d \otimes A \rightarrow J_d \otimes B$  is a normal monomorphism, which is in  $An$  if  $A \rightarrow B$  is.*

**Proof** Item (i) follows since normal monomorphisms are built by attaching boundary inclusions  $\partial\Omega[T] \rightarrow \Omega[T]$  for  $T \in \Omega_G$ . Similarly, (ii) follows from the “ $S$  linear” cases of Proposition 7.25 and Theorem 7.1 (recall that  $J_d$  is a simplicial set). □

As in [5, Section 3.4], for a  $G$ -dendroidal set  $B$  we now write  $dSet^G/B$  for the category of  $G$ -dendroidal sets over  $B$  and let  $An_B$  denote the class of maps in  $dSet^G/B$  whose image in  $dSet^G$  is in  $An$ .

**Proposition 8.3** *Write  $\partial^0: \{0\} \rightarrow J_d$ ,  $\partial^1: \{1\} \rightarrow J_d$  and  $\sigma: J_d \rightarrow \{*\}$  for the standard maps<sup>3</sup> and use the abbreviation  $\mathcal{J} = (J_d \otimes -, \partial^0 \otimes -, \partial^1 \otimes -, \sigma \otimes -)$ .*

*Then whenever  $B$  is normal the pair  $(\mathcal{J}, An_B)$  is a homotopical structure on  $dSet^G/B$  as defined in [4, Definition 1.3.14].*

**Proof** There are two parts: showing that  $\mathcal{J}$  is an elementary homotopical datum as in [4, Definition 1.3.6] and that  $An_B$  is a class of anodyne extensions with respect to  $\mathcal{J}$  as in [4, Definition 1.3.10]. We will refer to the axiom names used therein.

For the first claim, both axioms DH1 and DH2 follow from Proposition 7.25 since when  $B$  is normal all monomorphisms over  $B$  are normal monomorphisms.

For the second claim, axiom An0 follows by [4, Lemma 1.3.52] while An1 and An2 follow from Proposition 8.2. □

Combining the results [4, Theorem 1.3.22, Proposition 1.3.31, Proposition 1.3.36 and Lemma 1.3.52] now yields the following (which generalizes [5, Proposition 3.5 and Remark 3.6]).

<sup>3</sup>Note that  $\{*\} = i_!(*)$  denotes the terminal simplicial set, not the terminal dendroidal set.



**Proposition 8.4** For any normal  $G$ -dendroidal set  $B$ , the category  $\text{dSet}^G/B$  has a left proper cofibrantly generated model structure such that

- the cofibrations are the monomorphisms;
- $J$ -anodyne extensions over  $B$  are trivial cofibrations;
- the fibrant objects are the  $X \xrightarrow{p} B$  such that  $p$  is a  $J$ -fibration in  $\text{dSet}^G$ ;
- a map  $X \xrightarrow{f} X'$  between fibrant  $X \rightarrow B$  and  $X' \rightarrow B$  is a fibration in  $\text{dSet}^G/B$  if and only if  $f$  is a  $J$ -fibration in  $\text{dSet}^G$ .

The following generalizes [5, Lemma 3.7], again following from [4, Corollary 1.3.35].

**Lemma 8.5** Let  $X \rightarrow Y$  be a trivial fibration between normal  $G$ -dendroidal sets. Then any section  $s: Y \rightarrow X$  is in  $\text{An}$ .

Fix once and for all a normalization  $E_\infty$  of the terminal  $G$ -dendroidal set  $*$ , ie a trivial fibration  $E_\infty \rightarrow *$  with  $E_\infty$  a normal  $G$ -dendroidal set.

The following generalizes [5, Lemma 3.9].

**Lemma 8.6** For any normal  $G$ -dendroidal set  $X$  and map  $a: X \rightarrow E_\infty$ , the map  $(a, \text{id}): X \rightarrow E_\infty \times X$  is in  $\text{An}$ .

**Proof** Since  $(a, \text{id})$  is a section of the projection  $X \times E_\infty \rightarrow X$ , this follows from Lemma 8.5. □

The following generalizes [5, Lemma 3.10].

**Lemma 8.7** Let  $i: A \rightarrow B$  be a map of normal  $G$ -dendroidal sets and  $p: X \rightarrow Y$  any map of  $G$ -dendroidal sets. Then  $p$  has the right lifting property with respect to  $i$  if and only if, for any map  $B \rightarrow E_\infty$ , the map  $E_\infty \times X \xrightarrow{\text{id} \times p} E_\infty \times Y$  has the right lifting property with respect to  $i$  in  $\text{dSet}^G/E_\infty$ .

**Proof** Given a lifting problem as on the left below, one obtains a lifting problem as on the right by arbitrarily choosing a map  $B \rightarrow E_\infty$  (such a map always exists since  $B$  is assumed normal). It is clear that the lifting problems are equivalent:

$$\begin{array}{ccc}
 A & \longrightarrow & X \\
 i \downarrow & \nearrow & \downarrow p \\
 B & \longrightarrow & Y
 \end{array}
 \qquad
 \begin{array}{ccc}
 A & \longrightarrow & E_\infty \times X \\
 i \downarrow & \nearrow & \downarrow \text{id} \times p \\
 B & \longrightarrow & E_\infty \times Y
 \end{array}$$

□

We finally obtain the model structure on  $\text{dSet}^G$ , generalizing [5, Proposition 3.12].

**Proposition 8.8**  $\text{dSet}^G$  is equipped with a left proper cofibrantly generated model structure such that

- the cofibrations are the normal monomorphisms;
- inner  $G$ -anodyne extensions are trivial cofibrations;
- the fibrant objects are the  $J$ -fibrant objects;
- the fibrations between fibrant objects are the  $J$ -fibrations.

**Proof** As in [5, Proposition 3.12], the model structure is built via the adjunction

$$p_! : \text{dSet}^G / E_\infty \rightleftarrows \text{dSet}^G : p^*$$

where  $p_!(X \rightarrow B) = X$  and  $p^*(X) = (E_\infty \times X \rightarrow E_\infty)$ . Noting that both  $p_!$  and  $p^*$  preserve all colimits it follows that condition (iii) in [4, Proposition 1.4.23] reduces to verifying that  $p^*p_!(j)$  is a trivial cofibration in  $\text{dSet}^G / E_\infty$  whenever  $j$  is. Since this follows from Lemma 8.6, the transfer model structure exists.

The claim that cofibrations are normal monomorphisms follows since all monomorphisms over  $E_\infty$  are normal. The converse follows since all boundary inclusions  $\partial\Omega[T] \rightarrow \Omega[T]$  are in the image of  $p_!$ , and the claim concerning anodyne extensions follows by the same argument. The fibrancy claims follow from Lemma 8.7. Lastly, left properness follows from that in Proposition 8.4 together with the observation that  $p^*$  preserves both cofibrations and colimits and detects weak equivalences.  $\square$

**Remark 8.9** As in the case of the model structure on  $\text{dSet}$  or of the Joyal model structure on  $\text{sSet}$ , the trivial cofibrations (resp. fibrations) in Proposition 8.8 do not coincide with the inner  $G$ -anodyne extensions (resp.  $J$ -fibrations), but merely contain (resp. are contained in) them.

## 8.2 Characterization of fibrant objects

Much as in [5], the bulk of the work is now that of characterizing the fibrant objects as indeed being the  $G$ - $\infty$ -operads.

We will need to make use of the adjunction

$$u^* : \text{dSet}_G \rightleftarrows \text{dSet}^G : u_*$$

discussed in Proposition 5.53.

**Definition 8.10** The class of *normal monomorphisms* of  $\text{dSet}_G$  is the saturation of the maps of the form  $u_*(A \rightarrow B)$  for  $A \rightarrow B$  a normal monomorphism in  $\text{dSet}^G$ . Further, a map  $X \rightarrow Y$  in  $\text{dSet}_G$  is called a *trivial fibration* if it has the right lifting property with respect to normal monomorphisms.

We now extend some notation from [5, Section 6.1].

As usual, for  $X$  an  $\infty$ -category,  $k(X)$  will denote the maximal Kan complex inside  $X$  and  $\tau(X) \in \text{Cat}$  will denote its homotopy category.

**Notation 8.11** For a  $G$ - $\infty$ -operad  $X$  and simplicial set  $K$  (thought of as having a trivial  $G$ -action), we define  $X^{(K)} \in \text{dSet}_G$  to have  $T$ -dendrices (recall  $T \in \Omega_G$ ) the maps

$$i_!(K) \otimes \Omega[T] \xrightarrow{a} X$$

such that for each edge orbit  $G/H \cdot \eta \xrightarrow{G/H \cdot e} T$ , the induced map

$$K \xrightarrow{a_e} i^*(X^H)$$

factors through  $k(i^*(X^H)) = k((i^*(X))^H)$ .

**Remark 8.12** For  $X$  a  $G$ - $\infty$ -category one always has  $k(X^G) \subset k(X)^G$ , but this inclusion is rarely an equality.

**Notation 8.13** For a normal  $G$ -dendroidal set  $A \in \text{dSet}^G$  and a  $G$ - $\infty$ -operad  $X$  we define  $k(A, X) \in \text{sSet}$  to have  $n$ -simplices the maps

$$i_!(\Delta[n]) \otimes A \xrightarrow{b} X$$

such that for all element orbits  $G/H \cdot \eta \xrightarrow{G/H \cdot a} A$ , the induced map

$$\Delta[n] \xrightarrow{b_a} i^*(X^H)$$

factors through  $k(i^*(X^H)) = k((i^*(X))^H)$ .

Note that there are canonical isomorphisms

$$(8.14) \quad \text{Hom}_{\text{sSet}}(K, k(A, X)) \simeq \text{Hom}_{\text{dSet}_G}(u_*(A), X^{(K)}).$$

The following is the analogue of [5, Theorem 6.4]. Recall that a map  $\mathcal{C} \rightarrow \mathcal{D}$  in  $\text{Cat}$  is called a *categorical fibration* if it has the right lifting property against the inclusion  $1 \rightarrow \widetilde{[1]}$ , where  $\widetilde{[1]} = (0 \rightrightarrows 1)$  is the contractible groupoid with two objects.

**Theorem 8.15** *Let  $p: X \rightarrow Y$  be a  $G$ -inner fibration between  $G$ - $\infty$ -operads. Then*

$$(8.16) \quad \text{ev}_1: X^{(\Delta[1])} \rightarrow Y^{(\Delta[1])} \times_{u_*(Y)} u_*(X)$$

*has the right lifting property with respect to inclusions  $u_*(\partial\Omega[S] \rightarrow \Omega[S])$  for any  $G$ -tree  $S \in \Omega_G$  with at least one vertex.*

*Consequently, the map (8.16) is a trivial fibration in  $\text{dSet}_G$  if and only if all maps  $\tau i^*(X^H \rightarrow Y^H)$  for  $H \leq G$  are categorical fibrations.*

**Proof** Noting that  $S \simeq G \cdot_H S_e$  for some  $S_e \in \Omega^H$  and  $H \leq G$ , it suffices to deal with the case  $S \in \Omega^G$ .

The proof of the main claim now follows exactly as in the proof of [5, Theorem 6.4] by replacing uses of [5, Theorem 5.2] and [5, Theorem 4.2] with the equivariant analogues Theorems 7.2 and 7.4.

For the ‘‘consequently’’ part, one needs only note that in the equivariant context there are now multiple  $G$ -trees with no vertices, namely the  $G$ -trees of the form  $G/H \cdot \eta$ .  $\square$

The following is the analogue of [5, Proposition 6.7].

**Proposition 8.17** *Let  $p: X \rightarrow Y$  be a  $G$ -inner fibration between  $G$ - $\infty$ -operads. If  $\tau i^*(X^H \rightarrow Y^H)$  is a categorical fibration for all  $H \leq G$  then, for any monomorphism between normal dendroidal sets  $A \rightarrow B$ , the map*

$$(8.18) \quad k(B, X) \rightarrow k(B, Y) \times_{k(A, Y)} k(A, X)$$

*is a Kan fibration between Kan complexes.*

**Proof** We will mainly refer to the proof of [5, Proposition 6.7], while indicating the main changes. First, note that it follows from Theorem 7.1(ii) that the map  $\underline{\text{Hom}}(B, X) \rightarrow \underline{\text{Hom}}(B, Y) \times_{\underline{\text{Hom}}(A, Y)} \underline{\text{Hom}}(A, X)$  of simplicial mapping spaces is a  $G$ -inner fibration between  $G$ - $\infty$ -categories, and therefore so is (8.18). As in [5, Proposition 6.7], it now suffices to check that (8.18) has the right lifting property against the ‘‘left pushout products’’

$$\partial\Delta[n] \times \Delta[1] \cup_{\partial\Delta[n] \times \{1\}} \Delta[n] \times \{1\} \rightarrow \Delta[n] \times \Delta[1]$$

(ie thanks to [5, Lemma 6.5] one needs only consider the case  $i = 1$  of (8.1)). A lifting

problem

$$(8.19) \quad \begin{array}{ccc} \partial\Delta[n] \times \Delta[1] \cup \Delta[n] \times \{1\} & \xrightarrow{\quad} & k(B, X) \\ \downarrow & \nearrow h & \downarrow \\ \Delta[n] \times \Delta[1] & \xrightarrow{\quad} & k(B, Y) \times_{k(A, Y)} k(A, X) \end{array}$$

induces an (a priori nonequivalent) lifting problem:

$$(8.20) \quad \begin{array}{ccc} u_*(\partial\Omega[n] \otimes B \cup \Omega[n] \otimes A) & \xrightarrow{\quad} & X^{(\Delta[1])} \\ \downarrow & \nearrow \bar{h} & \downarrow \\ u_*(\Omega[n] \otimes B) & \xrightarrow{\quad} & Y^{(\Delta[1])} \times_{u_*(Y)} u_*(X) \end{array}$$

That the lift  $\bar{h}$  in (8.20) exists follows from [Theorem 8.15](#) and it hence remains to check that the adjoint map  $i_!(\Delta[n] \times \Delta[1]) \otimes B \rightarrow X$  indeed provides the map  $h$  in (8.19). That is, one must check that for any element orbit  $G/H \cdot b$  of  $B$  the induced map  $\Delta[n] \times \Delta[1] \rightarrow i^*(X^H)$  factors through  $k(i^*(X^H))$  (note that the existence of  $\bar{h}$  only guarantees such a factorization for the restriction along  $\{0, 1, \dots, n\} \times \Delta[1] \subset \Delta[n] \times \Delta[1]$ ). The  $n = 0$  case is immediate and the  $n > 0$  case follows by arguing using the “2-out-of-3 property”, just as in the penultimate paragraph of the proof of [\[5, Proposition 6.7\]](#).

Standard arguments (setting  $A = \emptyset$  or  $Y = *$  just as in [\[5, Proposition 6.7\]](#) or as at the end of the proof of [Theorem 7.4](#)) finish the proof.  $\square$

Combining the previous result with (8.14) now yields the following.

**Corollary 8.21** *Let  $p: X \rightarrow Y$  be a  $G$ -inner fibration between  $G$ - $\infty$ -operads such that  $\tau i^*(X^H \rightarrow Y^H)$  is a categorical fibration for all  $H \leq G$ . Then, for any anodyne extension of simplicial sets  $K \rightarrow L$ ,*

$$X^{(L)} \rightarrow Y^{(L)} \times_{Y^{(K)}} X^{(K)}$$

*is a trivial fibration in  $\text{dSet}_G$ .*

We now obtain our sought generalization of [\[5, Theorem 6.10\]](#).

**Theorem 8.22** *A  $G$ -dendroidal set  $X$  is  $J$ -fibrant if and only if it is a  $G$ - $\infty$ -operad. Further, a  $G$ -inner fibration  $p: X \rightarrow Y$  between  $G$ - $\infty$ -operads is a  $J$ -fibration if and only if  $\tau i^*(X^H \rightarrow Y^H)$  is a categorical fibration for all  $H \leq G$ .*

**Proof** It suffices to prove the “further” claim. Moreover, the “only if” direction is a direct consequence of [12, Corollary 1.6]. Unwinding definitions and adjunction properties it thus remains to show that

$$X^{J_d} \rightarrow Y^{J_d} \times_{Y^{\{0\}}} X^{\{0\}}$$

is a trivial fibration in  $\text{dSet}^G$  if  $\tau_i^*(X^H \rightarrow Y^H)$  is a categorical fibration for all  $H \leq G$ . We now note that since any map  $J_d \rightarrow Z$  necessarily factors through  $k(Z)$ , the map

$$u_*(X^{J_d} \rightarrow Y^{J_d} \times_{Y^{\{0\}}} X^{\{0\}})$$

coincides with the map

$$X^{(J_d)} \rightarrow Y^{(J_d)} \times_{Y^{\{0\}}} X^{\{0\}},$$

which is a trivial fibration in  $\text{dSet}_G$  by Corollary 8.21. The result now follows since  $\text{dSet}^G$  is a reflexive subcategory of  $\text{dSet}_G$ , so that  $u_*(f)$  is a trivial fibration if and only if  $f$  is. □

The following follows exactly as in [5, Corollary 6.11].

**Corollary 8.23** *The class of weak equivalences in  $\text{dSet}^G$  is the smallest class that contains the inner  $G$ -anodyne extensions and the trivial fibrations and is closed under “2-out-of-3”.*

## 9 Indexing system analogue results

In this section we follow the lead of [2] and build variant model structures on  $\text{dSet}^G$  associated to indexing systems, a notion originally introduced in [2, Definition 3.22], which we repackage (and slightly extend) in Definition 9.5.

**Definition 9.1** A  $G$ -graph subgroup of  $G \times \Sigma_n$  is a subgroup  $K \leq G \times \Sigma_n$  such that  $K \cap \Sigma_n = *$ .

**Remark 9.2**  $G$ -graph subgroups are graphs of homomorphisms  $G \geq H \rightarrow \Sigma_n$ .

**Definition 9.3** A  $G$ -vertex family is a collection

$$\mathcal{F} = \coprod_{n \geq 0} \mathcal{F}_n,$$

where each  $\mathcal{F}_n$  is a family of  $G$ -graph subgroups of  $G \times \Sigma_n$  closed under subgroups and conjugation.

Further, an  $H$ -set  $X$  for a subgroup  $H \leq G$  is called an  $\mathcal{F}$ -set if for some (and hence any) choice of isomorphism  $X \simeq \{1, \dots, n\}$  the graph subgroup of  $G \times \Sigma_n$  encoding the  $H$ -action on  $\{1, \dots, n\}$  is in  $\mathcal{F}$ .

**Definition 9.4** Let  $\mathcal{F}$  be a  $G$ -vertex family. A  $G$ -tree  $T$  is called an  $\mathcal{F}$ -tree if for all edges  $e \in T$  with isotropy  $H$  one has that the  $H$ -set  $e^\uparrow$  is an  $\mathcal{F}$ -set.

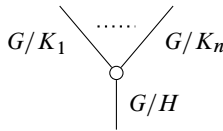
It is clear that whenever  $T \rightarrow S$  is either an outer face or a quotient,  $S$  being an  $\mathcal{F}$ -tree implies that  $T$  is also. However, the same is typically not true for inner faces and degeneracies.

**Definition 9.5** A  $G$ -vertex family  $\mathcal{F}$  is called a *weak indexing system* if  $\mathcal{F}$ -trees form a sieve of  $\Omega_G$ , ie if for any map  $T \rightarrow S$  with  $S$  an  $\mathcal{F}$ -tree,  $T$  is also an  $\mathcal{F}$ -tree. In this case we denote the sieve of  $\mathcal{F}$ -trees by  $\Omega_{\mathcal{F}} \subset \Omega_G$ .

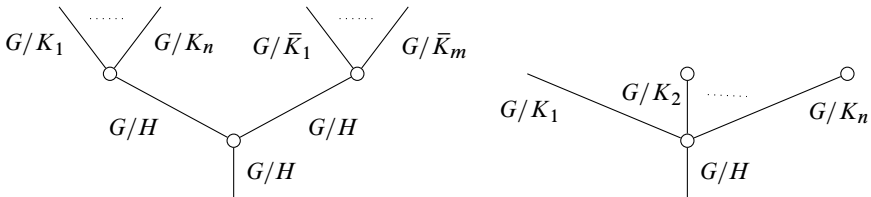
Additionally,  $\mathcal{F}$  is called an *indexing system* if every  $\mathcal{F}_n$  contains all subgroups  $H \times * \leq G \times \Sigma_n$  for  $H \leq G$  and  $n \geq 0$ .

**Remark 9.6** Closure under degeneracies is simply the statement that  $\mathcal{F}_1$  contains all subgroups  $H = H \times \Sigma_1 \leq G \times \Sigma_1$  for  $H \leq G$ .

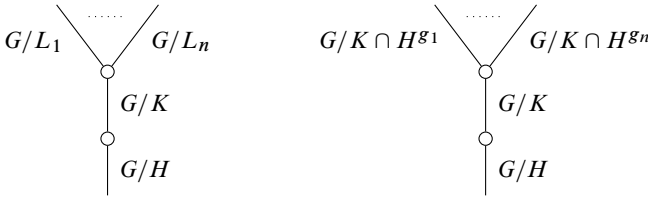
**Remark 9.7** Since [Definition 9.5](#) may at first seem to be quite different from the original [[2](#), Definition 3.22], we now address the equivalence between the two. To an  $H$ -set with orbital decomposition  $H/K_1 \amalg \dots \amalg H/K_n$  one can associate the  $G$ -corolla with orbital representation as follows:



Note that for any of its roots  $r$  one has that  $r^\uparrow$  is a  $G$ -conjugate of the  $H$ -set  $H/K_1 \amalg \dots \amalg H/K_n$ . The conditions (see [[2](#), Definition 3.22]) that indexing systems are closed under disjoint unions [[2](#), Definition 3.19] and subobjects [[2](#), Definition 3.21] of  $\mathcal{F}$ -sets are then encoded by taking inner faces of  $\mathcal{F}$ -trees of the form



while the closure under self-induction [2, Definition 3.20] is similarly encoded by  $\mathcal{F}$ -trees as on the left below:



Closure under cartesian products [2, Definition 3.22] is in fact redundant, as the double coset formula  $H/K \times H/L \simeq \coprod_{[g] \in L \backslash H/K} H/K \cap L^g$  allows such  $H$ -sets to be built using self-inductions as displayed by the rightmost tree above (the case of products of sets with multiple orbits being then obtained via disjoint units).

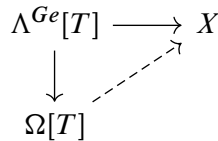
Definitions 6.3, 6.11 and 6.12 admit weak indexing system analogues.

**Definition 9.8** Let  $\mathcal{F}$  be a weak indexing system.

An  $\mathcal{F}$ -boundary inclusion (resp.  $\mathcal{F}$ -inner horn inclusion) is a boundary inclusion  $\partial\Omega[T] \rightarrow \Omega[T]$  (resp. inner horn inclusion  $\Lambda^{Ge}[T] \rightarrow \Omega[T]$ ) with  $T \in \Omega_{\mathcal{F}}$ .

A monomorphism is called  $\mathcal{F}$ -normal (resp.  $\mathcal{F}$ -anodyne) if it is in the saturation of  $\mathcal{F}$ -boundary inclusions (resp.  $\mathcal{F}$ -inner horn inclusions) under pushouts, transfinite compositions and retracts.

Finally, a  $G$ -dendroidal set  $X$  is called an  $\mathcal{F}$ - $\infty$ -operad if it has the right lifting property with respect to all  $\mathcal{F}$ -inner horn inclusions:



We now list the necessary modifications to extend the results in this paper to the indexing system case.

A direct analogue of Proposition 6.5 yields that  $X \in \text{dSet}^G$  is  $\mathcal{F}$ -normal (ie  $\emptyset \rightarrow X$  is an  $\mathcal{F}$ -normal monomorphism) if and only if all dendrices  $x \in X(T)$  have  $\mathcal{F}$ -isotropy, ie if and only if they all have isotropies  $\Gamma \leq G \times \Sigma_T$  that are graph subgroups for partial homomorphisms  $G \geq H \rightarrow \Sigma_T$  such that the induced  $G$ -tree  $G \cdot_H T$  is an  $\mathcal{F}$ -tree.

It then follows that, much like normal dendroidal sets,  $\mathcal{F}$ -normal dendroidal sets form a sieve; ie for any map  $X \rightarrow Y$  with  $Y$  an  $\mathcal{F}$ -normal dendroidal set then so is  $X$ .



Noting that the subtrees of  $S \otimes T$  are  $\mathcal{F}$ -trees whenever  $S$  and  $T$  are  $\mathcal{F}$ -trees (since the generating vertices/broad relations of  $S \otimes T$  are induced from those of  $S$  and  $T$ ), it follows that  $\Omega[S] \otimes \Omega[T]$  is then  $\mathcal{F}$ -normal so that the sieve condition implies that [Proposition 7.25](#) generalizes to the  $\mathcal{F}$ -normal case.

Likewise, the key results [Theorem 7.1](#) and [7.2](#) immediately generalize by replacing the terms “ $G$ -tree” and “ $G$ -anodyne” with “ $\mathcal{F}$ -tree” and “ $\mathcal{F}$ -anodyne”. This is because their proofs, while long, ultimately amount to identifying suitable edge orbits of suitable subtrees of  $S \otimes T$  and then attaching the corresponding equivariant horns.

Likewise, [Proposition 7.53](#) generalizes to the  $\mathcal{F}$  case for the same reason, and hence so does [Theorem 7.4](#), since its proof is an application of [Proposition 7.53](#).

We can now prove [Theorem 2.2](#).

**Proof of Theorem 2.2** The proof of the existence of the model structure follows just as in [Section 8.1](#). The only notable changes are as follows: in defining  $J_{\mathcal{F}}$ -anodyne extensions one uses only  $\mathcal{F}$ -inner horns and those maps in [\(8.1\)](#) for  $T \in \Omega_{\mathcal{F}}$ ; the term “normal” is replaced with “ $\mathcal{F}$ -normal” throughout (note that any monomorphism over an  $\mathcal{F}$ -normal dendroidal set is an  $\mathcal{F}$ -normal monomorphism).

The characterization of the  $J_{\mathcal{F}}$ -fibrant objects as being the  $\mathcal{F}$ - $\infty$ -operads follows by repeating the arguments in [Section 8.2](#), though some care is needed when adapting the definitions preceding [Theorem 8.15](#). Firstly, letting  $\Omega_{\mathcal{F}} \subset \Omega_G$  denote the sieve of  $\mathcal{F}$ -trees, one sets  $\text{dSet}_{\mathcal{F}} = \text{dSet}_{\Omega_{\mathcal{F}}}^{\text{op}}$ , leading to an adjunction

$$u^*: \text{dSet}_{\mathcal{F}} \rightleftarrows \text{dSet}^G : u_*$$

allowing for the  $\mathcal{F}$ -normal monomorphisms of  $\text{dSet}_{\mathcal{F}}$  to be defined from the  $\mathcal{F}$ -normal monomorphisms in  $\text{dSet}^G$  just as in [Definition 8.10](#).

For an  $\mathcal{F}$ - $\infty$ -operad and simplicial set  $K$ , one defines  $X^{(K)} \in \text{dSet}_{\mathcal{F}}$  just as in [Notation 8.11](#) while for  $A \in \text{dSet}^G$  an  $\mathcal{F}$ -normal dendroidal set and an  $\mathcal{F}$ - $\infty$ -operad  $X$  one defines  $k(A, X) \in \text{sSet}$  just as in [Notation 8.13](#).

The proofs of [Theorem 8.15](#), [Proposition 8.17](#) and [Theorem 8.22](#) now extend mutatis mutandis by using the  $\mathcal{F}$  versions of [Theorems 7.2](#) and [7.4](#).  $\square$

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