

Minimal *k*-rankings for prism graphs

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We determine rank numbers for the prism graph $P_2 \times C_n$ (P_2 being the connected two-node graph and C_n a cycle of length n) and for the square of an even cycle.

1. Introduction

A *k*-ranking of a graph is a vertex labeling using integers between 1 and *k* inclusive such that any path between two vertices of the same rank contains a vertex of strictly larger rank. When the value of *k* is unimportant, we will refer to a *k*-ranking simply as a ranking. A ranking *f* is minimal if the reduction of any label violates the ranking property [Ghoshal et al. 1996]. Another definition of a minimal ranking is obtained by replacing the reduction of a label by the reduction of labels for any nonempty set of vertices. It was shown in [Jamison 2003] and [Isaak et al. 2009] that these two definitions of minimal rankings are equivalent. The *rank number* of a graph *G*, denoted $\chi_r(G)$ is the smallest *k* such that *G* has a minimal *k*-ranking.

Recall that a vertex coloring of a graph is a vertex labeling in which no two adjacent vertices have the same label. Hence a *k*-ranking is a restricted vertex coloring. Then the rank number is similar to the chromatic number. The *arank* number of a graph G, denoted $\psi_r(G)$, is the largest k such that G has a minimal k-ranking.

The study of the rank number was motivated by applications including the design of very large scale integration (VLSI) layout and Cholesky factorizations associated with parallel processing [de la Torre et al. 1992; Ghoshal et al. 1996; 1999;

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Leiserson 1980; Laskar and Pillone 2001; 2000; Sen et al. 1992]. Numerous related papers have since followed [Bodlaender et al. 1998; Hsieh 2002; Jamison 2003; Dereniowski 2006; 2004; Dereniowski and Nadolski 2006; Kostyuk and Narayan ≥ 2010 ; Kostyuk et al. 2006; Isaak et al. 2009; Novotny et al. 2009a]. Ghoshal, Laskar, and Pillone were the first to investigate minimal *k*-rankings [Ghoshal et al. 1999; 1996; Laskar and Pillone 2001; 2000]. The determination of the rank number and the arank number was shown to be NP-complete [Laskar and Pillone 2000]. The rank number was explored in [Bodlaender et al. 1998] where the authors showed that $\chi_r(P_n) = \lfloor \log_2 n \rfloor + 1$. Rank numbers are known for a few other graph families such as cycles, wheels, complete bipartite graphs, and split graphs [Ghoshal et al. 1996; Dereniowski 2004]. The rank number for ladder graphs $P_2 \times P_n$ and the square of a path P_n^2 were determined in [Novotny et al. 2009b].

Throughout the paper P_n will denote the path on *n* vertices. We use $G \times H$ to denote the *Cartesian product* of *G* and *H*. The *k*-th power of a path, P_n^k , has vertices v_1, v_2, \ldots, v_n and edges (v_i, v_j) for all *i*, *j* satisfying $|i - j| \le k$. The *k*-th power of a cycle, C_n^k , is defined similarly.

In this paper we determine rank numbers for the prism graph $P_2 \times C_n$ and the square of an even cycle.

We begin by restating two elementary results from [Ghoshal et al. 1996].

Lemma 1. In any minimal ranking of a connected graph G the highest label must be unique.

Proof. Suppose there exist two vertices u and v that both have the highest label k. Then any path between u and v will not contain a vertex with a higher label. This is a contradiction.

The following lemma gives a monotonicity result involving the rank number.

Lemma 2. Let *H* be a subgraph of a graph *G*. Then $\chi_r(H) \leq \chi_r(G)$.

Proof. The proof is straightforward. Suppose $\chi_r(H) > \chi_r(G)$. Then we could relabel the vertices of *H* using the corresponding labels used in the ranking of *G*. This produces a ranking with fewer labels, and hence a contradiction.

1.1. The ladder graph L_n . We next describe a family of graphs built using the *Cartesian product*.

Definition 3. The *Cartesian product* of *G* and *H* written $G \times H$ is the graph with vertex set $V(G) \times V(H)$ specified by putting $\{u, v\}$ adjacent to (u', v') if and only if u = u' and $(v, v') \in E(H)$ or v = v' and $(u, u') \in E(G)$.

An example is the ladder graph $L_n = P_2 \times P_n$, shown in Figure 1.

In this paper we investigate the family of prism graphs $P_2 \times C_n$. We will start with a ladder $P_2 \times P_n$ with *n* even, and insert either a $P_2 \times P_1$ or $P_2 \times P_2$ and



Figure 1. The ladder graph $L_n = P_2 \times P_n$.

"wrap" the ends to form a prism graph $P_2 \times C_{n+1}$ or $P_2 \times C_{n+2}$. In order for this construction to work, it is essential that in the labeling of the vertices labeled 1 of the ladder satisfies an "alternating 1's property": for each vertex v, either vis labeled 1 or all of its neighbors are labeled 1 (Figure 2). That is, the vertices labeled 1 form a particular dominating set of the graph. It was shown in [Novotny et al. 2009b] that in a minimal ranking of a ladder the 1's can be made to alternate.



Figure 2. A graph with the alternating 1s property.

We can insert in $P_2 \times P_n$ either a 1-bridge (Figure 3, left) or a 2-bridge (Figure 3, right). In general, the bridges will contain the labels k and k + 1 where k - 1 is the rank of the original ladder. Our example shows the extension where k = 6.

In each case we insert four edges to connect the bridge to each end of the ladder. When n is even the wrapping of the ladder L_n creates a prism graph where the 1's alternate. When n is odd the 1's alternate except in one place where there are two vertices labeled 1 that are distance 3 apart (Figure 4).

Novotny et al. [2009b] determined the rank number of a ladder graph. This result is stated in our next lemma.

Lemma 4. $\chi_r(L_n) = \lfloor \log_2(n+1) \rfloor + \lfloor \log_2(n+1-2^{\lfloor \log_2 n \rfloor -1}) \rfloor + 1$ for $n \ge 1$.

Applying our construction immediately gives an upper bound for the rank number of the prism graph $P_2 \times C_n$, as stated in our next theorem.

Theorem 5. For $k \ge 2$, both $\chi_r(P_2 \times C_{2k-1})$ and $\chi_r(P_2 \times C_{2k})$ are bounded from above by r(2k-2)+2.

We will show later that this bound is tight.



Figure 3. A 1-bridge (left) and 2-bridge (right).



Figure 4. Prism graphs for *n* even (left) and *n* odd (even).

2. Main results

Theorem 6. Let $l = \chi_r(P_2 \times C_n)$ where $n \ge 3$. If f is a minimal l-ranking of $P_2 \times C_n$, then $l \ge 5$ and the largest four labels of f appear exactly once.

Proof. In the minimal ranking $f : V(P_2 \times C_n) \rightarrow \{1, 2, ..., l\}$ every label appears at least once. Since $G = P_2 \times C_n$ is (vertex) 3-connected, any two distinct vertices of *G* are joined by three internally vertex disjoint paths. Hence each of the largest three labels appears exactly once in *f*.

Assume that l-3 appears at least twice with f(x) = f(y) = l-3, where $x \neq y$. We have $l \geq 5$ because the independence number of G is $2\lfloor n/2 \rfloor$ and $2\lfloor n/2 \rfloor + 3 < 2n = |V(G)|$.

Let *S* be a minimum-sized *x*, *y* vertex separating set. It is clear that |V(S)| = 3. It is well known that every 3-element separating set \tilde{S} is a prism graph *P* is a neighborhood of a single vertex $\tilde{z} \in V(P)$ and the nontrivial component of $P - \tilde{S}$ is induced by $V(P) - (\tilde{S} \cup \{\tilde{z}\})$. Thus, there exists $z \in \{x, y\}$ such that *S* is the neighborhood of *z*. However if *z* has its neighbors labeled l - 2, l - 1, and *l*, then f(x) can be reduced to 1, contradicting the minimality of *f*.

For a positive integer *n* let

$$r(n) = \lfloor \log_2(n+1) \rfloor + \lfloor \log_2(n+1 - (2^{\lfloor \log_2 n \rfloor - 1})) \rfloor + 1.$$
(1)

Then Lemma 4 states that $\chi_r(L_n) = \chi_r(P_2 \times P_n) = r(n)$ for $n \ge 1$.

Theorem 7. For $k \ge 2$, we have

$$\chi_r(P_2 \times C_{2k-1}) = \chi_r(P_2 \times C_{2k}) = \chi_r(P_2 \times P_{2k-2}) + 2 = r(2k-2) + 2.$$

Proof. By Theorem 5, both $\chi_r(P_2 \times C_{2k-1})$ and $\chi_r(P_2 \times C_{2k})$ are bounded from above by r(2k-2)+2. In other words, if m = 2k-1 or 2k, then

$$\chi_r(P_2 \times C_m) \le \chi_r(P_2 \times P_{2\lceil m/2 \rceil - 2}) + 2 = r(2\lceil m/2 \rceil - 2) + 2.$$

To prove the theorem we will show that this last inequality is in fact equality. If k = 2 and m = 2k - 1 or 2k, then $r(2\lceil m/2 \rceil - 2) + 2 = 5$. So by Theorem 6, $\chi_r(P_2 \times C_m) = 5$.

Now assume that m = 2k - 1 or $2k, k \ge 3$, and

$$\chi_r(P_2 \times C_m) = l \le r\left(2\left\lceil \frac{m}{2} \right\rceil - 2\right) + 1.$$
⁽²⁾

Let *f* be an *l*-minimal ranking of $G = P_2 \times C_m$. If k = 3, then $5 \le l = r(4) + 1 \le 5$, l = 5, and by Theorem 6, the label 1 appears 2m - 4 times in *f*. However the independence number of *G* equals $2\lfloor m/2 \rfloor \le m < 2m - 4$, which is a contradiction.

Let $k \ge 4$. This implies $m \ge 7$. Let *i* be the maximum label used at least twice. Since $r(2\lceil m/2 \rceil - 2) + 1 \le r(m-1) + 1 < 2m = |V(G)|$, such a label does exist, and $i \leq l-4$ by Theorem 6. Consider vertices $x_1, x_2 \in V(G)$ with $f(x_1) = i = f(x_2)$, and let y_i be the neighbor of x_i that is not on the "ring" containing x_i . We will refer to this vertex as the special neighbor of x_j for j = 1, 2. There are two distinct subgraphs G_1, G_2 of G that are ladders with corners x_1, x_2, y_1, y_2 . The restriction $f|_{V(G_i)}$ is a ranking of G_i ; hence there is a minimal separating set $S_i \subseteq V(G_i)$ such that min $f(S_i) > i$ and x_1, x_2 are in distinct components of $G_i - S_i$, j = 1, 2. It is easy to see that any minimal separating set that separates two "distant" corners of a ladder on at least six vertices has two vertices and is of one of the two types shown in Figure 3 (consisting of the vertices labeled k and k+1). As all labels in $\{i+1, \ldots, l\}$ are used by f exactly once, any permutation of those labels yields a ranking of G. Therefore, we may suppose without loss of generality that $f(S_1) \cup f(S_2) = \{l-3, l-2, l-1, l\}$. Further, let \overline{S}_i be the set consisting of the vertices of S_i together with their special neighbors (so that $|\bar{S}_i|$ is 2 or 4). The graph $G - (\bar{S}_1 \cup \bar{S}_2)$ is a union of two vertex disjoint ladders H_1 and H₂. Clearly if $|V(H_1)| \ge |V(H_2)|$, then $H_1 = P_2 \times P_q$, where $q \ge \lceil (m-4)/2 \rceil$. Now $f|_{V(H_1)}$ uses only labels from the set $\{1, \ldots, l-4\}$; hence, by (2),

$$\chi_r(H_1) \le l - 4 \le r\left(2\left\lceil \frac{m}{2} \right\rceil - 2\right) - 3. \tag{3}$$

On the other hand if s, t are positive integers with $s \le t$, then $P_2 \times P_s$ is a subgraph of $P_2 \times P_t$. Then by Lemma 2 we have $r(s) = \chi_r(P_2 \times P_s) \le \chi_r(P_2 \times P_t) = r(t)$. Consequently,

$$\chi_r(H_1) = \chi_r(P_2 \times P_q) = r(q) \ge r\left(\left\lceil \frac{m-4}{2} \right\rceil\right). \tag{4}$$

If m is even, then it follows from Equations (3) and (4) that

$$r(m-2) = r\left(2 \cdot \frac{m-4}{2} + 2\right) \ge r\left(\frac{m-4}{2}\right) + 3.$$

If *m* is odd we have

$$r(m-1) = r\left(2 \cdot \frac{m-3}{2} + 2\right) \ge r\left(\frac{m-3}{2}\right) + 3.$$

However both cases lead to a contradiction. From (1) it is easy to see that

$$r(2n+2) - r(n) = 2$$

for any positive integer *n*.

Since
$$r(2k-3) = r(2k-2)$$
 for $n \ge 3$, we obtain from Theorem 7:

Theorem 8. $\chi_r(P_2 \times C_n) = \chi_r(L_{n-2}) + 2$ for $n \ge 4$.

3. Rankings for other classes of graphs

We now show that the rank number of a prism graph can be used to give the rank number of the square of an even cycle. We recall some earlier facts:

Definition 9 [Ghoshal et al. 1996]. For a graph G and a set $S \subseteq V(G)$ the *reduction* of G, denoted by G_S^{\flat} , is a subgraph of G induced by V - S with an edge uv in $E(G_S^{\flat})$ if and only if there exists a u - v path in G with all internal vertices belonging to S.

Lemma 10 [Ghoshal et al. 1996]. Let G be a graph and let f be a minimal k-ranking of G. If

$$S_1 = \{x \in V(G) : f(x) = 1\}$$
 and $f^{\flat} : V(G_{S_1}^{\flat}) \to \{1, \dots, k-1\}$

is defined by $f^{\flat}(x) = f(x) - 1$, then f^{\flat} is a minimal (k-1)-ranking of $G_{S_1}^{\flat}$.

3.1. The square of a cycle. Next we reduce even prism graphs to squares of cycles.

Theorem 11. $\chi_r(C_n^2) = \chi_r(P_2 \times C_n)$ for even $n \ge 4$.

Proof. (Illustrated in Figure 5.) If n = 2, the result follows from Theorem 7 which states that $\chi_r(P_2 \times C_4) = 5$ and from the fact that $\chi_r(C_4^2) = \chi_r(K_4) = 4$.

Henceforth suppose that $n \ge 3$. Let $k = \chi_r(P_2 \times C_{2n})$ and let $l = \chi_r(C_{2n}^2)$. Let f be a *k*-ranking of $P_2 \times C_{2n}$ in which the 1's alternate. It is straightforward to see that then $(P_2 \times C_{2n})_{S_1}^{\flat}$ is isomorphic to C_{2n}^2 . Therefore, by Lemma 10 $\chi_r(C_{2n}^2) \le k - 1$.

Now let g be an *l*-ranking of C_{2n}^2 . One can easily see that C_{2n}^2 is isomorphic to an *n*-sided antiprism A_n . Pick a new vertex inside each of the 2n triangles of A_n , join it to all three vertices of "its" triangle and delete all edges of A_n . The result is a graph G_n that is isomorphic to $P_2 \times C_{2n}$. Consider the mapping $\tilde{g}: V(G_n) \rightarrow \{1, \ldots, l+1\}$ defined as follows: $\tilde{g}(x) = g(x) + 1$ if $x \in V(C_{2n}^2)$ and $\tilde{g}(x) = 1$ if $x \in V(G_n) - V(C_{2n}^2)$. Since \tilde{g} is a ranking of G_n (a simple exercise left to the reader), we have $\chi_r(P_2 \times C_{2n}) \leq l+1$.

 \square



Figure 5. A minimal 7-ranking of $P_2 \times C_8$ (left) and a minimal 6-ranking of A_4 (right).

Thus $l = \chi_r(C_{2n}^2) \le k - 1 = \chi_r(P_2 \times C_{2n}) - 1 \le (l+1) - 1 = l$, and since both inequalities turn into equalities, we are done.

Combining Theorems 8 and 11 gives:

Corollary 12. Let $n \ge 4$ be even. Then

 $\chi_r(C_n^2) = \chi_r(P_2 \times C_n) - 1 = \lfloor \log_2(n-1) \rfloor + \lfloor \log_2(n-1) - (2^{\lfloor \log_2(n-2) \rfloor - 1}) \rfloor + 2.$

4. Conclusion

We conclude by posing some problems for future research. In this paper we determined the rank number of $P_2 \times C_n$ using known results for the rank number of $P_2 \times P_n$. It would be interesting to determine the rank numbers for grid graphs $P_m \times P_n$ and cylinders $P_m \times C_n$. We found out recently that [Alpert ≥ 2010] gives rank numbers for $P_3 \times P_n$, among other results including an alternate proof of our Theorem 7.

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